

MA/AMA 514

Today: Matroid Intersection (§10.4, §10.5)

Recall: Independent set matroid axioms $M=(X, \mathcal{I})$

- (1) \mathcal{I} nonempty collection of subsets of X
- (2) closed under containment ($T \in \mathcal{I}, S \subseteq T \Rightarrow S \in \mathcal{I}$)
- (3) $S, T \in \mathcal{I}, |T| > |S| \Rightarrow \exists t \in T \setminus S$ with $S \cup \{t\} \in \mathcal{I}$

Ex (graphic) $G=(V, E)$ $\mathcal{I} = \{F \subseteq E : (V, F) \text{ acyclic}\}$

Ex (linear) $v_1, \dots, v_n \in \mathbb{R}^d$ $\mathcal{I} = \{S \subseteq [n] : \{v_i : i \in S\} \text{ lin. indep}\}$

Partition Matroids ground set X , partition $X = A_1 \uplus \dots \uplus A_m$

$M=(X, \mathcal{I})$ $\mathcal{I} = \{S \subseteq X \text{ s.t. } |S \cap A_i| \leq 1 \text{ for all } i\}$

= transversal matroid with sets A_1, \dots, A_m $\iff \exists$ bijection $f: S \rightarrow [m]$ with $s \in A_{f(s)} \forall s \in S$

Matroid intersection

Given two matroids on the same groundset,

find the largest common indep. set.

Input: $M_1=(X, \mathcal{I}_1)$ and $M_2=(X, \mathcal{I}_2)$

Goal: $\max \{ |I| : I \in \mathcal{I}_1 \cap \mathcal{I}_2 \}$

Ex (bipartite matching) $G=(U \uplus W, E)$ bipartite graph

ground set: $X = E$

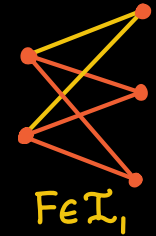
$M_1 =$ transversal matroid with sets $A_u = \delta(u)$ for $u \in U$

$M_2 =$ " " " $A_w = \delta(w)$ for $w \in W$

$F \in \mathcal{I}_1 \Leftrightarrow |F \cap \delta(u)| \leq 1$ for all $u \in U$

$F \in \mathcal{I}_2 \Leftrightarrow |F \cap \delta(w)| \leq 1$ for all $w \in W$

$F \in \mathcal{I}_1 \cap \mathcal{I}_2 \Leftrightarrow F$ matching



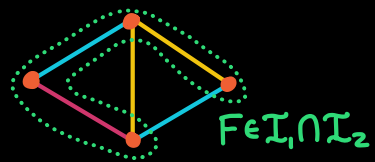
Ex (rainbow forests) $G=(V,E)$ connected graph

Edges E colored with $|V|-1$ colors

$M_1=(E, \mathcal{I}_1)$ graphic matroid of G ($\mathcal{I}_1 = \{\text{forests}\}$)

$M_2=(E, \mathcal{I}_2)$ trans. matroid with $A_i = \{\text{edges with color } i\}$

$F \in \mathcal{I}_2 \Leftrightarrow F$ has at most one edge of each color



$F \in \mathcal{I}_1 \cap \mathcal{I}_2 \Leftrightarrow F$ is a "rainbow" forest

Ex (disjoint spanning trees) $G=(V,E)$ connected graph

$M_1=(E, \mathcal{I}_1)$ graphic matroid of G ($\mathcal{I}_1 = \{\text{forests}\}$)

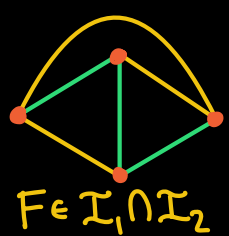
$M_2=(E, \mathcal{I}_2)$ cographic matroid of G ($\mathcal{I}_2 = \{S \subseteq E : S \subseteq E \setminus T \text{ for some span. tree } T\}$)

G has two edge-disjoint spanning trees

$= \{S \subseteq E : (V, E \setminus S) \text{ connected}\}$

$\Leftrightarrow \max\{|I| : I \in \mathcal{I}_1 \cap \mathcal{I}_2\} = |V|-1$

Ex.



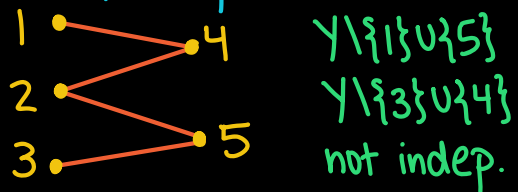
We will give a polynomial time algorithm for matroid intersection. First we need two technical lemmas:

Two exchange lemmas (§10.4)

Let $M=(X, \mathcal{I})$ be a matroid and $Y \in \mathcal{I}$.

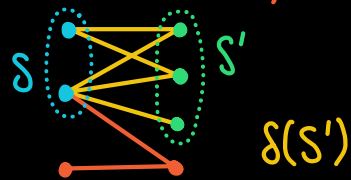
Define $H_M(Y)$ to be the bipartite graph with vertices $Y \sqcup (X \setminus Y)$ and edges $\{y, x\}$ for $y \in Y, x \in X \setminus Y$ with $Y \setminus \{y\} \cup \{x\} \in \mathcal{I}$

Ex: $M = \text{graphic matroid}$
 $Y = \{1, 2, 3\}$



Lemma 10.1 If $Y, Z \in \mathcal{I}$ with $|Y| = |Z|$, then $H_M(Y)$ has a perfect matching on $Y \Delta Z$.

(Proof) Suppose not. By Hall's marriage theorem, $\exists S \subseteq Y \setminus Z$ and $S' \subseteq Z \setminus Y$ with $|S| < |S'|$ and $\delta(S') \subseteq S$. Since $(Y \cap Z) \cup S \in \mathcal{I}$ and $(Y \cap Z) \cup S' \in \mathcal{I}$, both are in \mathcal{I} and $\exists z \in S'$ s.t.



$$T := (Y \cap Z) \cup S \cup \{z\} \in \mathcal{I}.$$

Similarly, we can add elt's of Y to T until we get $\exists U \in \mathcal{I}$ with $T \subseteq U \subseteq T \cup Y$ and $|U| = |Y|$.

\Rightarrow for some $y \in Y \setminus T$, $U = (Y \setminus \{y\}) \cup \{z\} \Rightarrow \{y, z\}$ an edge. \neq
 $(y \notin S, z \in S')$

Lemma 10.2 Let $Y \in \mathcal{I}$ and $Z \subseteq X$ with $|Y| = |Z|$.

If $H_M(Y)$ contains a unique perfect matching on $Y \Delta Z$ then $Z \in \mathcal{I}$.

(Proof) By induction on $k = |Z \setminus Y|$. ($k=0$) $Z = Y \in \mathcal{I}$.

($k \geq 1$) Let N be the unique perfect matching on $Y \Delta Z$.

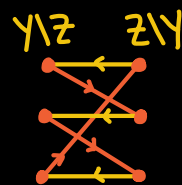
Claim: the restriction of $H_M(Y)$ to $Y \Delta Z$ has a leaf $y \in Y \setminus Z$.

That is $\exists y \in Y \setminus Z$ s.t. $|\{z \in Z \setminus Y : Y \setminus \{y\} \cup \{z\} \in \mathcal{I}\}| = 1$.

(Proof) Define walk on restriction. Start at any $w \in Y \Delta Z$.

At $z \in Z \setminus Y$, walk along edge $\{z, y\} \in N$.

At $y \in Y \setminus Z$, walk along edge $\{y, z\} \notin N$, if it exists.

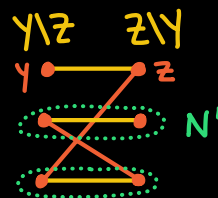


Either walk terminates at a leaf $y \in Y \setminus Z$ or results in an even cycle $C \Rightarrow N \Delta C$ is another perfect matching on $Y \Delta Z$. \square Claim

Let $y \in Y \setminus Z$ be a leaf and $z \in Z \setminus Y$ with $\{y, z\} \in N$.

Take $Z' = (Z \setminus \{z\}) \cup \{y\}$ and $N' = N \setminus \{\{y, z\}\}$.

Then $|Y \Delta Z'| < |Y \Delta Z|$ and N' is unique matching on $Y \Delta Z' \Rightarrow Z' \in \mathcal{I}$ (by induction).



Know: $(Y \setminus \{y\}) \cup \{z\} \in \mathcal{I}$ since $\{y, z\} \in H_M(Y)$.

We can add an elt. $w \in Y \setminus \{y\} \cup \{z\}$ to $Z' \setminus \{y\} \in \mathcal{I}$ s.t.

$S = (Z' \setminus \{y\}) \cup \{w\} \in \mathcal{I}$. If $w = z$, $Z = S \in \mathcal{I}$. \checkmark

Otherwise $w \in Y \setminus \{y\}$. Since $S \in \mathcal{I}$ with $|S| > |Y \setminus \{y\}|$, $\exists z' \in S \setminus Y$ s.t. $(Y \setminus \{y\}) \cup \{z'\} \in \mathcal{I}$, contradicting y a leaf!

Edmond's Alg for matroid intersection

Given $M_1 = (X, \mathcal{I}_1)$, $M_2 = (X, \mathcal{I}_2)$ and $Y \in \mathcal{I}_1 \cap \mathcal{I}_2$ define directed bipartite graph $H_{M_1, M_2}(Y)$ with vertices $X = Y \uplus X \setminus Y$, arcs

(y, x) if $(Y \setminus \{y\}) \cup \{x\} \in \mathcal{I}_1$ ($\Leftrightarrow \{y, x\} \in H_{M_1}(Y)$)

(x, y) if $(Y \setminus \{y\}) \cup \{x\} \in \mathcal{I}_2$ ($\Leftrightarrow \{y, x\} \in H_{M_2}(Y)$)

Idea: look for directed path in $H_{M_1, M_2}(Y)$, use it to "augment" Y to a larger set $Y' \in \mathcal{I}_1 \cap \mathcal{I}_2$.