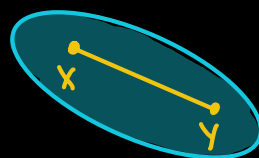


# Math 407 - Linear Optimization

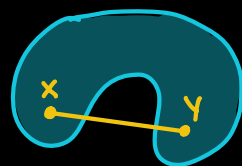
Today: Convexity, Polyhedra, Polytopes

## Convexity Basics

A set  $P \subseteq \mathbb{R}^n$  is convex if for every  $x, y \in P$  and  $\lambda \in [0, 1]$ ,  $\lambda x + (1-\lambda)y \in P$   
(i.e.  $P$  contains the line segment between any two of its points)

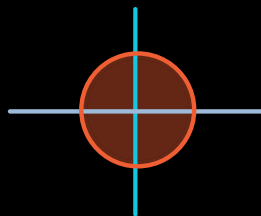


Convex



not convex

Ex:  $\{(x, y) \in \mathbb{R}^2 : x^2 + y^2 \leq 1\}$

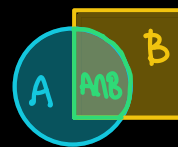


Prop: The intersection of convex sets is convex.

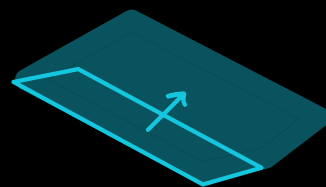
Why?  $A, B$  convex,  $x, y \in A \cap B$ ,  $\lambda \in [0, 1]$

$\Rightarrow \lambda x + (1-\lambda)y$  belongs to  $A$  and to  $B$

(by convexity of  $A, B$ )  $\Rightarrow$  belongs to  $A \cap B$ .

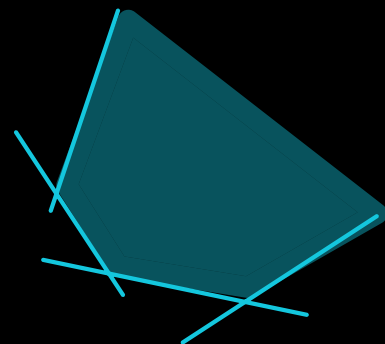


A halfspace in  $\mathbb{R}^n$  is a set of the form  $\{x \in \mathbb{R}^n : a^T x \leq b\}$  where  $a \in \mathbb{R}^n \setminus \{0\}$ ,  $b \in \mathbb{R}$ .



A polyhedron is the intersection of finitely many half-spaces.

(= feasible regions of LPs)



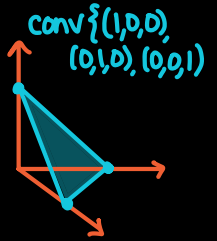
Cov: Any polyhedron is convex.

A convex combination of points  $p_1, \dots, p_s \in \mathbb{R}^n$  has the form  $\lambda_1 p_1 + \lambda_2 p_2 + \dots + \lambda_s p_s$  where  $\lambda_1, \lambda_2, \dots, \lambda_s \geq 0$  and  $\sum_{i=1}^s \lambda_i = 1$ .

The convex hull of a set  $S$  is

$$\text{conv}(S) = \left\{ \lambda_1 p_1 + \dots + \lambda_k p_k : \begin{array}{l} k \in \mathbb{N}, p_1, \dots, p_k \in S \\ \lambda_1, \dots, \lambda_k \geq 0, \sum_{i=1}^k \lambda_i = 1 \end{array} \right\}$$

taking  $k=1$  shows  $S \subseteq \text{conv}(S)$



Prop: For any  $c \in \mathbb{R}^n$  any finite set  $S \subseteq \mathbb{R}^n$

$$\max \{ c^T x \text{ s.t. } x \in S \} = \max \{ c^T x \text{ s.t. } x \in \text{conv}(S) \}$$

(Proof) ( $\leq$ ) follows from  $S \subseteq \text{conv}(S)$

( $\geq$ ) Suppose  $m = \max \{ c^T x : x \in S \}$  and let  $\tilde{x} \in \text{conv}(S)$ .

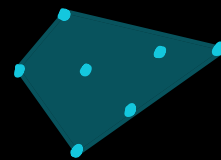
By def,  $\tilde{x} = \lambda_1 p_1 + \dots + \lambda_k p_k$  for some  $p_1, \dots, p_k \in S$ ,

$\lambda_1, \dots, \lambda_k \geq 0$  with  $\sum \lambda_i = 1$ . Then

$$\begin{aligned} c^T \tilde{x} &= c^T (\lambda_1 p_1 + \dots + \lambda_k p_k) = \lambda_1 c^T p_1 + \dots + \lambda_k c^T p_k \\ &\leq \lambda_1 m + \dots + \lambda_k m \\ &= (\lambda_1 + \dots + \lambda_k) m = m \end{aligned}$$

using  $\lambda_i \geq 0$   
and  $c^T x \leq m$   
for  $x \in S$

A polytope is the convex hull of a finite set of points.



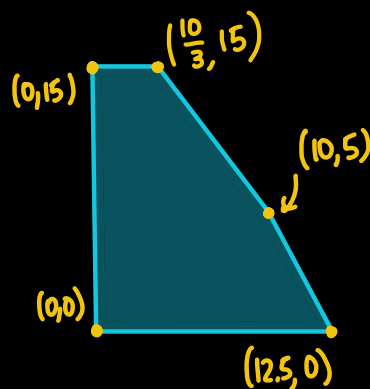
Ex (scarves ; hats from last time)

$$\max 10s + 6h \quad \text{s.t.}$$

$$0 \leq s \leq 15, \quad 0 \leq h \leq 15$$

$$6s + 4h \leq 80,$$

$$200s + 100h \leq 2500$$



Feasible region equals the convex hull of 5 pts.

Given constraints, optimal production will be one of only five possibilities, regardless of what profit per item is!