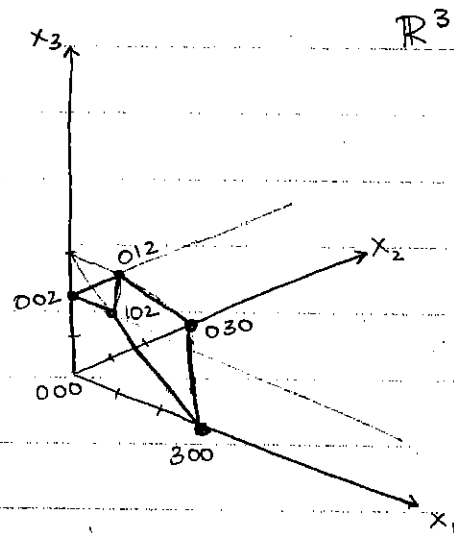


The Geometry of the Simplex Method

$$\begin{aligned} \max \quad & x_1 + 10x_2 + 100x_3 \\ \text{s.t.} \quad & x_1 + x_2 + x_3 \leq 3 \\ & x_3 \leq 2 \\ & x_1, x_2, x_3 \geq 0 \end{aligned}$$



first dictionary

$$x_4 = 3 - x_1 - x_2 - x_3$$

$$x_5 = 2 - x_3$$

$$z = x_1 + 10x_2 + 100x_3$$

hyperplanes corresponding to the constraints:

$$H_4 = H(x_1 + x_2 + x_3 = 3), \quad H(x_3 = 2) = H_5$$

$$H_1 = H(x_1 = 0), \quad H_2 = H(x_2 = 0), \quad H(x_3 = 0) = H_3$$

Prob: $\max z \mid x_1, \dots, x_5 \geq 0$

$$x_0 = (0, 0, 0, 3, 2) \quad z = 0$$

If $x = (x_1, \dots, x_5)$ is a feasible solⁿ,

- basis / nonbasis
- Choose x_1 to enter the basis
 \Rightarrow *allowing distance of solⁿ from $H(x_1=0)$ to increase from 0.

x_1	- distance of x from	$H(x_1=0)$
x_2	- " "	$H(x_2=0)$
x_3	- " "	$H(x_3=0)$
x_4	- " "	$H(x_1+x_2+x_3=3)$
x_5	- " "	$H(x_3=2)$

* Keep $x_2 = x_3 = 0$

- Min ratio test says that x_4 leaves
 \Rightarrow solⁿ has hit the hyperplane $H(x_1 + x_2 + x_3 = 3)$

\rightarrow solⁿ travels along an edge to new basic feas. solⁿ

- new basis $\{x_1, x_5\}$ nonbasis $\{x_2, x_3, x_4\}$
- monotone increase.

2nd dictionary:

$$x_1 = 3 - x_2 - x_3 - x_4$$

$$x_5 = 2 - x_3$$

$$z = 3 + 9x_2 + 99x_3 - x_4$$

Choose x_3 to enter - allowing solⁿ to leave the $H_{(x_3=0)}$

keep $x_2 = x_4 = 0$ \therefore edge is uniquely chosen.

min ratio test - x_5 leaves.

- solⁿ has hit the hyperplane $H_{(x_3=2)}$

new basis $\{x_1, x_3\}$ nonbasis $\{x_2, x_4, x_5\}$

cost has improved.

3rd dictionary

$$x_3 = 2 - x_5$$

$$x_1 = 1 - x_2 - x_4 + x_5$$

$$z = 201 + 9x_2 - 99x_5 - x_4$$

Choose x_2 to enter ... x_1 leaves.

$$x_2 = 1 - x_1 - x_4 + x_5$$

$$x_3 = 2 - x_5$$

$$z = 210 - 9x_1 - 10x_4 - 90x_5$$

← optimal dictionary

Opt solⁿ: $(0, 1, 2, 0, 0)$

← solⁿ lies on the hyperplanes H_1, H_4, H_5

Q: Why are there always 3 elts in the nonbasis?

→ x_1, \dots, x_n measure distance from H_1, \dots, H_n

→ x_{n+1}, \dots, x_{n+m} measure distance to the complicated hyperplane

In general,

$$\max c \cdot x$$

$$Ax \leq b$$

$$x \geq 0$$

$$A = \begin{pmatrix} a_1 \\ \vdots \\ a_m \end{pmatrix} \in \mathbb{R}^{m \times n}$$

$$\max \sum_{j=1}^n c_j x_j$$

$$\text{s.t. } \sum_{j=1}^n a_{ij} x_j \leq b_i \quad i=1, \dots, m$$

$$x_j \geq 0 \quad j=1, \dots, n$$

$P = \{x \in \mathbb{R}^n : Ax \leq b, x \geq 0\}$ - polyhedron / polytope if it is bounded.

→ convex: if $a, b \in P$ then $[a, b] \in P$

→ has flat "faces" or sides

→ dim (usually) $n \Rightarrow$ sides have dim $n-1$

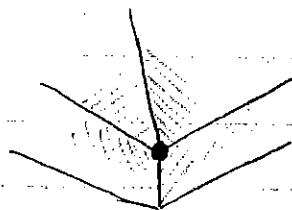
→ has faces of dim $n-1 \rightarrow 0 \leftarrow$ vertices / corners

→ Cost hyperplane gets optimized on a face \rightarrow include a corner

→ Simplex method creates a monotone path (cost never becomes worse) from an initial solⁿ to the opt along edges of the feasible region.

Every basic feasible solⁿ has at least n vars = 0
indeed a corner lies at the intersection of at least n hyperplanes.

→ More vars can be zero:



4 hyperplanes create this corner

Some basic var is also zero.

but nonbasic