## Problem Set 4

## **CSE 521 - Design and Analysis of Algorithms**

Fall 2024

## Exercise 1 (10pts)

For  $1 \le p < \infty$  and  $x \in \mathbb{R}^n$ ,  $||x||_p := (\sum_{i=1}^n |x_i|^p)^{1/p}$  is the  $||\cdot||_p$ -norm which generalizes the Euclidean norm  $||\cdot||_2$  and Manhattan norm  $||\cdot||_1$  from the lecture. A probability distribution  $\mathcal{D}_p$  over  $\mathbb{R}$  is said to be *p*-stable if for  $Z, Z_1, \ldots, Z_n$  independently drawn from  $\mathcal{D}_p$  and for any fixed  $x \in \mathbb{R}^n$ , the random variable  $\sum_{i=1}^n x_i Z_i$  is equal in distribution to  $||x||_p \cdot Z$ . Some examples are the standard normal distribution N(0, 1), which is 2-stable. Another less-known example is the Cauchy distribution, which is 1-stable; it has probability density function  $\Phi(x) = 1/(\pi(1+x^2))$ . It is a known theorem that such distributions exist iff  $p \in (0,2]$ . Note that *p*-stable random variables for  $p \neq 2$  cannot have bounded variance, since otherwise the sum of independent copies would have to be gaussian by central limit theorem. In fact, it is known that any *p*-stable distributions have *bounded and continuous* density function and they must have tail bounds  $\Pr(|Z| > \lambda) = O(1/(1+\lambda)^p)$  for all  $\lambda > 0$ . This implies that such distributions cannot exist for p > 2 (since otherwise they would have bounded variance, violating the central limit theorem).

- i) Suppose *Z* is *p*-stable; show that for any  $\alpha > 0$ ,  $\alpha Z$  is also *p*-stable.
- ii) **Optional (0 points).** Let Z be a *p*-stable random variable normalized (by a constant) so that  $Pr[Z \in [-1,1]] = 1/2$  (see previous part). Fix some  $\varepsilon > 0$ . Show that there is a constant c > 0 (as a function of p but not  $\varepsilon$ ) such that

$$\Pr[-1 + \varepsilon < Z < 1 - \varepsilon] \le 1/2 - c\varepsilon,$$
  
$$\Pr[-1 - \varepsilon < Z < 1 + \varepsilon] \ge 1/2 + c\varepsilon.$$

iii) Let *Z* be a *p*-stable random variable with  $\Pr[Z \in [-1,1]] = \frac{1}{2}$ . Let  $P \in \mathbb{R}^{m \times d}$  where  $P_{i,j}$  is an independent sample of *Z*. Let  $x \in \mathbb{R}^d$  arbitrary and y = Px; show that for  $m = O(\log(1/\delta)/\varepsilon^2)$ , with probability at least  $1 - \delta$ , the median of  $|y_1|, \ldots, |y_m|$  is a  $1 \pm \varepsilon$  multiplicative approximation of  $||x||_p$ .

Hint. It might be helpful to review the argument in Theorem 1.37.

## Exercise 2 (10pts)

In this problem we design a *locally sensitive hash function (LSH)* for points in  $\mathbb{R}^d$ , with the  $\ell_1$  distance, i.e.

$$d(p,q) = \sum_{i=1}^{d} |p_i - q_i|.$$

i) Let *a*, *b* be arbitrary real numbers. Fix w > 0 and let  $s \in [0, w)$  chosen uniformly at random. Show that

$$\Pr\left[\left\lfloor\frac{a-s}{w}\right\rfloor = \left\lfloor\frac{b-s}{w}\right\rfloor\right] = \max\left\{0, 1 - \frac{|a-b|}{w}\right\}$$

Recall that for any real number c, |c| is the largest integer which is at most c.

**Hint:** Start with the case where a = 0.

ii) Define a class of hash functions as follows: Fix *w* larger than diameter of the space. Each hash function is defined via a choice of *d* independently selected random real numbers  $s_1, s_2, \ldots, s_d$ , each uniform in [0, w). The hash function associated with this random set of choices is

$$h(x_1,\ldots,x_d) = \left( \left\lfloor \frac{x_1 - s_1}{w} \right\rfloor, \left\lfloor \frac{x_2 - s_2}{w} \right\rfloor, \ldots, \left\lfloor \frac{x_d - s_d}{w} \right\rfloor \right).$$

Let  $\alpha_i = |p_i - q_i|$ . What is the probability that h(p) = h(q) in terms of the  $\alpha_i$  values? Use the approximation  $1 - x \approx e^{-x}$  (just ignore the error term for the sake of simplicity).