## CONNECTIONS AND DETERMINANTS: A GEOMETRIC FORMULATION

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ABSTRACT. This paper contains a geometric argument for some of the connectiondeterminant relations discussed in section 3.7 of the book *Inverse Problems* for *Electrical Networks* by Edward B. Curtis and James A. Morrow.

In James A. Morrow and Edward B. Curtis' book [1], they prove the following useful property of response matrices  $\Lambda$  of an electrical network: If P and Q are two disjoint sets of boundary vertices with |P| = |Q| = n, such that  $det(\Lambda(P;Q)) \neq 0$ , then there exists an *n*-connection from P to Q. The proof they give is via a rather abstract determinantal formula. It is my hope that a more geometric argument for the same property might be an aid to intuition. In this paper, a more direct geometric proof of the theorem is obtained.

To do this we will make use of the concept of cutsets of a graph:

**Definition 1.** (Cutset, Minimal Cutset): Given a graph  $G = \{V, E\}$  (where V is the set of vertices and E is the set of edges) and two disjoint sets of vertices  $P, Q \subset V$  a cutset of the pair P, Q is a set  $R \subset V$  such that all paths from a vertex  $p \in P$  to a vertex  $q \in Q$  intersect R. A minimal cutset of the pair P, Q is a cutset of P, Q of minimal order among all such cutsets.

This allows us to quote and use a theorem about cutsets, Menger's n-Arc Theorem:

**Theorem 1.** (Menger, 1927 [2]): Let G be a graph with A and B two disjoint n-tuples of graph vertices. Then either G contains n pairwise disjoint AB-paths, or the minimal cutset of A, B in G has order d < n.

Finally, before we begin the proof we introduce the following notation:

**Notation.** Where A is a matrix,  $\vec{u}$  is a vector, and P and Q are two sets of vertices, let A(P;Q) be the submatrix of A with row indices in P and column indices in Q. Let  $\vec{u}(P)$  be the sub-vector of  $\vec{u}$  with indices in P, and let  $\vec{u}[P]$  be the sub-vector of  $\vec{u}$  excluding entries with indices in P.

We are now prepared to give a geometric proof for the connection-determinant theorem.

**Theorem 2.** If P and Q are two disjoint sets of boundary vertices with |P| = |Q| = n, such that  $det(\Lambda(P;Q)) \neq 0$ , then there exists an n-connection from P to Q.

*Proof.* Let  $\Gamma = (G, \gamma)$  be an electrical network on a graph with boundary  $G = \{V, \partial V, E\}$ , and suppose that P and Q are two disjoint sets of boundary vertices with |P| = |Q| = n, such that  $det(\Lambda(P; Q)) \neq 0$ .

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Let  $\hat{G}$  be the graph formed from G by removing all boundary vertices other than the vertices in P or Q, and all edges incident to those vertices. Let  $R = \{r_1, ..., r_d\}$ be a minimal cutset of the pair P, Q, in  $\hat{G}$ . Therefore, in G, any path from a vertex  $p_i \in P$  to a vertex  $q_j \in Q$  must pass through some vertex  $r_k \in R$ , or a boundary vertex  $s \in \partial V - (P \cup Q)$ . Let  $\vec{v} \in \mathbb{R}^n$  be arbitrary, and let  $\vec{u} \in \mathbb{R}^n$  be the vector  $\vec{u} = \Lambda(P; Q)^{-1}\vec{v}$ . Let  $\vec{u}' \in \mathbb{R}^{|\partial V|}$  be the vector with  $\vec{u}'(Q) = \vec{u}$  and  $\vec{u}'[Q] = \vec{0}$ . Consider the vector  $\vec{u}'$  to be a vector of boundary voltages on  $\partial V$ ; this gives rise to a solution of the Dirichlet problem - let  $\vec{w} \in \mathbb{R}^d$  be the vector of voltages on the vertices  $r_1...r_d$  in the Dirichlet solution. Note that the current into the network at the set of nodes P is given by  $(\Lambda \vec{u}')(P) = \Lambda(P, Q)\vec{u} = \vec{v}$ .

Now consider the network  $\Gamma' = (G', \gamma)$  obtained from  $\Gamma = (G, \gamma)$  by declaring all the vertices in R to be boundary vertices, and let  $\Lambda'$  be the response matrix of this new network. The graph G' has the same set of vertices V, but a new set of boundary vertices  $\partial V' = \partial V \cup R$ . Apply a boundary voltage  $\vec{w}'$  to  $\Gamma'$  equal to  $\vec{u}'$ on  $\partial V$  and  $\vec{w}$  on R. (It is possible that  $\partial V \cap R$  will be non-empty, but  $\vec{u}'(i) = \vec{w}(i)$ on any shared vertices i because  $\vec{w}$  is part of the solution to the Dirichlet problem with boundary voltage  $\vec{u}'$ , so this boundary voltage is well-defined.) This boundary voltage gives rise to the same voltages on each vertex in V as the boundary voltage given by  $\vec{u}'$  on the old graph  $\Gamma$ , so the current response is the same - in particular, the currents into the graph at the vertices in P are given by  $\vec{v}$ . The current is given by  $\Lambda'\vec{w}'$ , so  $(\Lambda'\vec{w}')(P) = \vec{v}$ . Since Q + R - P is the set of vertices where  $\vec{w}'$ is nonzero:

$$(\Lambda'\vec{w}')(P) = \Lambda'(P, Q + R - P)\vec{w}'(Q + R - P) = \vec{v}$$

and since  $\vec{v}$  was chosen arbitrarily, we have:

$$Rank(\Lambda'(P, Q + R - P)) = n$$

There is no connection through G' from any vertex  $q \in Q - R$  to any vertex  $p \in P - R$ , so  $\Lambda'(P - R; Q - R) = 0$  (where 0 here is the  $|P - R| \times |Q - R|$  0-matrix). Because  $Rank(\Lambda'(P; Q + R - P)) = n$ , the rows of  $\Lambda'(P; Q + R - P)$  are linearly independent. Ignoring the rows of  $\Lambda'(P; Q + R - P)$  that correspond to vertices in  $P \cap R$ , we have:

$$Rank(\Lambda'(P-R;Q+R-P)) = |P-R|$$

Note that since  $\Lambda'(P-R;Q-R) = 0$ :

$$Rank(\Lambda'(P-R;Q+R-P)) = Rank(\Lambda'(P-R;R-P)) = |P-R|$$

Finally, have:

$$|P| - |P \cap R| = |P - R| = Rank(\Lambda'(P - R; R - P)) \le |R - P| = |R| - |P \cap R|$$

Since  $n = |P| \le |R| = d$ , the degree of the minimal cutset of the pair P, Q is at least n; by Menger's *n*-Arc theorem there is an *n*-connection between P and Q.  $\Box$ 

## References

- Curtis, Edward B. & Morrow, James A., *Inverse Problems for Electrical Networks*, Series on Applied Mathematics - Vol. 13, World Scientifc, New Jersey, 2000.
- [2] Menger, Karl, Kurventheorie, Teubner, Berlin, Germany, 1932.

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