

## Algorithms reference

### Greatest common factors — Euclidean algorithm

To compute the greatest common factor of  $a$  and  $b$ :

1. Put  $a$  and  $b$  in the first column of a table (bigger number on top).
2. Fill in the next column: if the previous column contains  $x$  and  $y$ , the next column contains  $y$  and  $x \bmod y$ .
3. Repeat step 2 until you get a 0. The number above the 0 is the greatest common factor of  $a$  and  $b$ .

$a$	...	$x$	$y$	...	greatest common factor
$b$		$y$	$x \bmod y$		0

### Multiplicative inverses — extended Euclidean algorithm

To compute the multiplicative inverse of  $a$  modulo  $q$ :

1. Write down the Euclidean algorithm table for  $a$  and  $q$ . A multiplicative inverse exists as long as the greatest common factor of  $a$  and  $q$  is 1.
2. Add another row to the table: if some column contains  $x$  and  $y$ , the third row will contain  $x \div y$  rounded **downwards**. Ignore the column with 0.

$x$	
$y$	
round( $x \div y$ )	

3. Add a fourth row: put a 1 under the first column and a 0 sticking out to the left of the 1, then fill in each entry with (cell two to the left)  $-$  (cell to left)  $\times$  (cell above left).

	$z$	
$x$	$y$	$x - y \times z$

4. When you reach the second-last column, the last number in the fourth row is the multiplicative inverse of  $a$  modulo  $q$ .

### Modular powers

To compute  $a^b \bmod q$ :

1. Set  $T = 1$ .
2. Convert  $b$  to binary, and read left to right. For each digit:
  - If the digit is 0: replace  $T$  with  $T^2 \bmod q$ .
  - If the digit is 1: replace  $T$  with  $T^2 \times a \bmod q$ .
3. The final value of  $T$  is  $a^b \bmod q$ .