

**Definition 1.** Let  $\{0, 1\}^*$  be the set of all binary strings of finite length.

**Definition 2.** A language over  $\{0, 1\}$  is a subset  $A \subset \{0, 1\}^*$ .

**Example 1.**  $EVEN = \{n \geq 0 \in \mathbb{Z} \text{ written in binary } | n \text{ is even}\}$ , then  $n = \{0, 00, 10, 000, 010, 100, 110, \dots\}$ .

**Example 2.**  $PRIME = \{p \geq 2 \in \mathbb{Z} \text{ written in binary } | p \text{ is prime}\}$ , then  $p = \{10, 11, 101, \dots\}$ .

More generally, can consider any set  $\Sigma$  (which we call alphabet) and set  $\Sigma^*$ , the set of all string of  $\Sigma$  of finite length.

**Definition 3.** A language over  $\Sigma$  is a subset  $A \subset \Sigma^*$ .

**Example 3.**  $\Sigma = \{0, 1, 2\}$ .

**Example 4.**  $\Sigma = \mathbb{Z}/n := \{0, 1, \dots, n-1\}$ .

**Example 5.**  $\Sigma = \mathbb{Z}$ .

**Example 6.**  $\Sigma = \mathbb{R}$ .

**Goal 1.** Given a language  $A \subset \Sigma^*$ , we would like a computational model for determining if a given string  $w \in \Sigma^*$  is in  $A$ .

Turing Machines:

**Definition 4** (Conceptual). A **Turing Machine** is a finite state machine (which has a finite set (which we will call states) and a finite set of rules which govern the action depending on the current state and what the head reads from the tape).

The action is either

- a print a symbol form  $\Sigma$ .
- b moves to left or right.
- c move to a different state.
- d terminate with "accept" or "reject."

**Definition 5** (Precise). A **Turing Machine** is a 7-tuple  $M = (Q, \Sigma, \Gamma, \delta, q_0, q_{\text{accept}}, q_{\text{reject}})$

1.  $Q$  is a finite set (of states).
2.  $\Sigma$  is the input alphabet (not containing the blank symbol).
3.  $\Gamma$  is tape alphabet (containing the blank,  $\Sigma \subset \Gamma$ ).
4.  $\delta$  is  $Q \times \Gamma \rightarrow Q \times \Gamma \times \{L, R\}$
5.  $q_0 \in Q$  is the start state.
6.  $q_{\text{accept}} \in Q$  is the accept state.
7.  $q_{\text{reject}} \in Q$  is the reject state.

**Definition 6.** Let  $A \subset \Sigma^*$ . We say a turing maching  $M$  **recognizes**  $A$  if for any  $w = w_1, w_2, \dots, w_n \in \Sigma^*$ , then  $M$  accepts  $w \Leftrightarrow w \in A$ .

**Definition 7.** A language is **recognizable** if there exists a turing machine which recognizes it.