

HW

5/1/2019

PERFECT GP

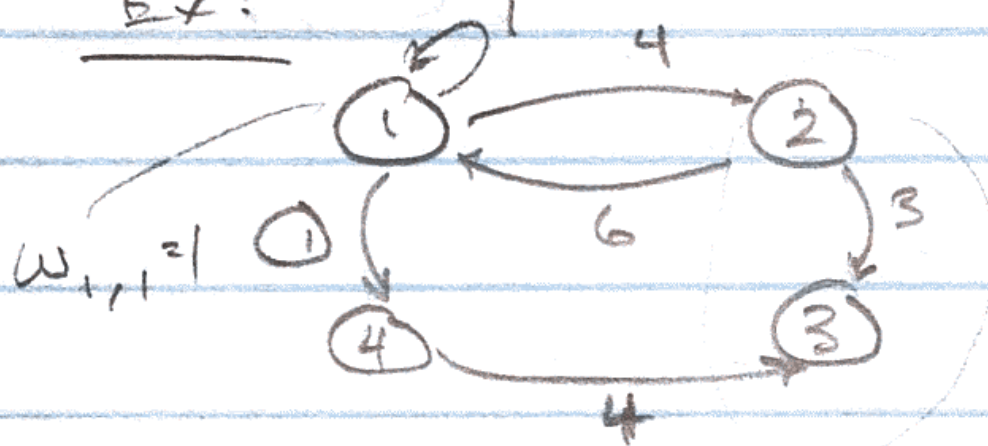
NP-complete is Counting Version

Today: Determinants & permanents.

Approach via graphs

Let Graph G be directed graph w/ vertices $1, 2, \dots, n$.

Ex:



$w_{1,1} = 1$

$w_{2,3} = 3$

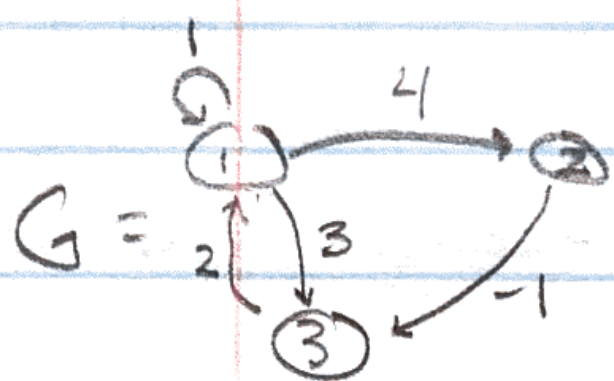
Let $w_{ij} = 0$ if \nexists edge $i \rightarrow j$

Def: Adjacency Matrix

$$A_G = \begin{pmatrix} w_{11} & \dots & w_{1n} \\ \vdots & \ddots & \vdots \\ w_{n1} & \dots & w_{nn} \end{pmatrix}$$

Ex:

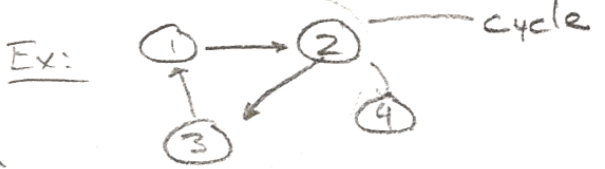
$$A = \begin{pmatrix} 1 & 4 & 3 \\ 0 & 0 & 1 \\ 2 & 0 & 0 \end{pmatrix}$$



Prop: \exists 1:1 bijection $G \rightarrow A_G$ & $G_A \leftarrow A$

Def: Let G be labelled direct graph w/ vertices $1, \dots, n$.

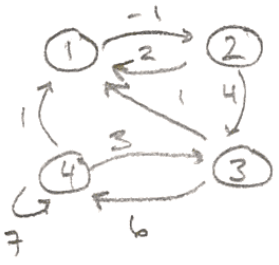
A cycle σ in G is path starting & ending @ same vertex & doesn't pass through the same vertex twice.



Def:

Cycle cover = union of finitely many cycles such that every vertex in the graph is in precisely one cycle.

Ex: Let G be below:



Def: If G is labelled directed graph:

$$\text{perm}(G) = \sum_{\sigma} \text{weight}(\sigma)$$

$$\det(G) = \sum_{\sigma} \text{sgn}(\sigma) \text{weight}(\sigma)$$

Cycle Cover:

- $\sigma_1: 1 \rightarrow 2 \rightarrow 3 \rightarrow 4 \rightarrow 1$
- $\sigma_2: 1 \rightarrow 2 \rightarrow 1 \quad 3 \rightarrow 4 \rightarrow 3$
- $\sigma_3: 1 \rightarrow 2 \rightarrow 3 \rightarrow 1 \quad 4 \rightarrow 4$

Ex:

$$\text{perm}(G) = (-24) + (-36) + (-28) = -88$$

$$\det(G) = (-1) \cdot 24 + 1 \cdot (-36) + 1 \cdot (-28) = -40$$

If A is $n \times n$ matrix associated

with graph G_A , then
 $\text{perm}(A) := \text{perm}(G_A)$
 $\det(A) := \det(G_A)$

Cycle Notation

- $(1 \ 2 \ 3 \ 4)$
- $(1 \ 2) (3 \ 4)$
- $(1 \ 2 \ 3) (4)$

Weight:

$$(-1) \cdot 4 \cdot 6 \cdot 1 = -24$$

$$(-1) \cdot 2 \cdot 3 \cdot 6 = -36$$

$$(-1) \cdot 4 \cdot 1 \cdot 7 = -28$$

Def: Sign of a cycle cover:

$$\sigma = \sigma_1 \cup \sigma_2 \cup \dots \cup \sigma_n$$

is

$$\text{sign}(\sigma) = \prod (-1)^{|\sigma_i| - 1}$$

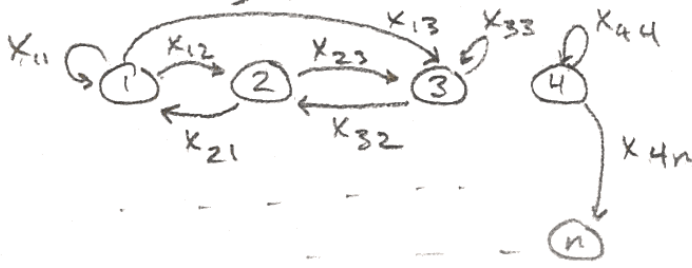
where $|\sigma_i| = \#$ of vertices in cycle σ_i .

Det & Perm'n of poly.'s:

Consider the $n \times n$ matrix:

$$X_n = \begin{pmatrix} x_{11} & x_{12} & \dots & x_{1n} \\ \vdots & \vdots & \ddots & \vdots \\ x_{ni} & \dots & \dots & x_{nn} \end{pmatrix}$$

Corresponding graph is:



Def:

$$\det \det_n := \det(X_n) = \det(G_n)$$

$$\text{perm}_n := \text{perm}(X_n) = \text{perm}(G_n).$$

$$n=3: X_3 = \begin{pmatrix} x_{11} & x_{12} & x_{13} \\ \vdots & \vdots & \vdots \\ x_{31} & \dots & x_{33} \end{pmatrix} \quad * \text{ You can draw graph.}$$

$$\begin{aligned} \text{perm}_3 = & x_{11} x_{22} x_{33} + x_{12} x_{21} x_{33} \\ & + \dots + x_{13} x_{32} x_{21} \end{aligned}$$

Prop: $\{\det_n\} \in VP$

Sketch: $\left. \begin{array}{l} \# \text{ of var.'s} = n^2 \\ \deg(\det_n) = n \end{array} \right\} \text{ poly bdd.}$

Complexity?

FACT: (Lin. dg). $\det(AB) = \det(A) \cdot \det(B)$.

Cor: Elementary row/column operations do not change det of matrix.

HW: $G =$ bipartite graph

perf. matchings = $\text{perm}(G)$.

For