# Avoiding $L_{\infty}$ Discrepancy Optimization

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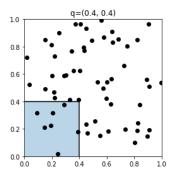


# The $L_{\infty}$ star discrepancy

### $L_{\infty}$ star discrepancy

For P a point set in  $[0;1]^d$ ,

$$d_{\infty}^*(P) = \sup_{q \in [0;1)^d} \left| \frac{\left| P \cap [0,q) \right|}{|P|} - \lambda([0,q)) \right|.$$



#### Local discrepancy:

$$D(q, P) = 0.044$$

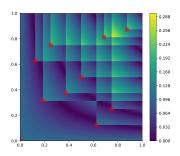
## Very small instances: optimal values

- The minimal star discrepancy,  $d_{\infty}^*(n,d)$ , is the best possible  $L_{\infty}$  star discrepancy value for a point set of size n in dimension d.
- [White, 1977] gave point sets for  $n \le 6$  in dimension 2
- 1-point sets for any d have been solved by [Pillard, Cools and Vandewoestyne, 2006], extended to 2 points by [Larcher and Pillichshammer, 2007]
- For the periodic  $L_2$  discrepancy, [Hinrichs and Oettershagen, 2016] solved the problem for  $n \le 16$

Can we provide point sets matching  $d_{\infty}^*(n,d)$ ?

## Computing the star discrepancy

Calculating the discrepancy is a discrete problem, maximal values can only be reached on a grid defined by the points.



## Computing the star discrepancy

- From the discrete "positions-grid":  $O(n^d)$ ,  $O(n^d/d!)$  if we only count **critical boxes**
- Best known algorithm:  $O(n^{1+d/2})$  by Dobkin, Eppstein and Mitchell (1996)
- Best heuristic in higher dimensions: Threshold Accepting algorithm by Gnewuch, Wahlström and Winzen (2012)

# Optimal constructions<sup>1</sup>

### Optimal $L_{\infty}^*$ star discrepancy set

Given an integer  $n \ge 1$  and a dimension  $d \ge 2$ , find a set P of size n in dimension d of discrepancy  $d_{\infty}^*(n,d)$ .

 Our two non-linear programming formulations rely on the grid structure of the discrepancy calculation.

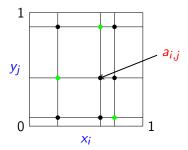
<sup>1</sup>Constructing Optimal  $L_{\infty}$  Star Discrepancy Sets, F.C, C. Doerr, K. Klamroth and L. Paquete, Proceedings of the American Mathematical Society Series B, 2025

## The assignment formulation

- We are trying to determine where the n points should be to minimize the discrepancy value, by placing the underlying grid and deciding which points should generate it.
- The objective z represents the discrepancy of the selected set.
- Variables  $x_i$  correspond to the ordered x coordinates of the points. The  $y_i$  to the ordered y coordinates.
- The binary variables  $a_{i,j}$  correspond to the selected grid-points.

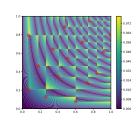
### The assignment formulation

We split the problem in two parts: finding the coordinates and finding an assignment.



### First formulation

s.t. 
$$\frac{1}{n} \sum_{u=1}^{i} \sum_{v=1}^{j} a_{uv} - x_i y_j \le z$$
$$\frac{-1}{n} \sum_{u=1}^{i-1} \sum_{v=1}^{j-1} a_{uv} + x_i y_j \le z$$



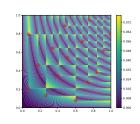
#### For each box, we need:

- the number of points inside:  $\sum_{u=1}^{i} \sum_{v=1}^{j} a_{uv}$
- its volume:  $x_i y_i$

### First formulation

min z

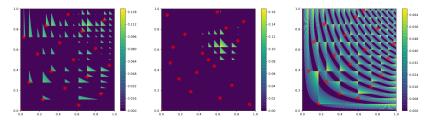
s.t. 
$$\frac{1}{n} \sum_{u=1}^{i} \sum_{v=1}^{j} a_{uv} - x_{i} y_{j} \le z$$
$$\frac{-1}{n} \sum_{u=1}^{i-1} \sum_{v=1}^{j-1} a_{uv} + x_{i} y_{j} \le z$$



For each box, we need:

- the number of points inside:  $\sum_{u=1}^{i} \sum_{v=1}^{j} a_{uv}$
- its volume:  $x_{2i-1}x_{2i}$

# Fibonacci vs Sobol' vs Optimal



Left: Fibonacci 18; Middle: Sobol' 18; Right: Optimal 18

### A computational obstacle

- We can solve these models to optimality for up to around 20 points in 2d.
- Similar models can be found for other discrepancy measures.
- Going further: focus on half of the problem! Either we fix the grid preemptively, or the  $a_{i,j}$ .

# The correct choice: Fixing the permutation <sup>2</sup>

- Fixing the coordinates does not help at all.
- Fixing the assignment (equiv. permutation) makes the problem much easier to solve.
- Can solve for up to 500 points, and the main obstacle is reading the model and presolving.

How to choose the correct permutation?

<sup>&</sup>lt;sup>2</sup>F. C., Carola Doerr, Kathrin Klamroth, Luís Paquete. Searching permutations for constructing uniformly distributed point sets. PNAS 2025

### A natural candidate: shifted Fibonacci sets

For some shift  $j \in \mathbb{N}$ ,  $P = \{(i/n, \{\phi(i+j)\}) : i \in \{1, ..., n\}\}$ 

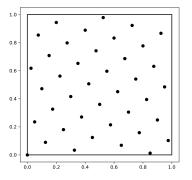


Figure: The unshifted Fibonacci set for 40 points

### A natural candidate: shifted Fibonacci sets

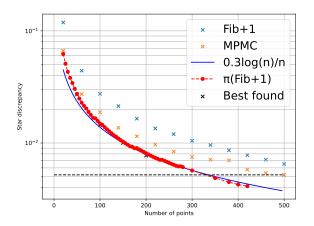


Figure: Best  $L_{\infty}$  star discrepancy values obtained by taking the permutation from the Fibonacci set *offset by 1*.

# Structure of the Fibonacci permutations

- Cycle structure of the permutation is very regular.
- For n = F<sub>k</sub> many points, either there are only cycles of length 2 and fixed points, or cycles of length 4 and a unique fixed points/cycle of length 2.
- Similar observations are possible for quadratic irrationals.
- Is this even relevant for discrepancy?

### Key questions

- How can we determine if a given permutation can lead to a low-discrepancy point set?
- Can we replace the discrepancy optimization by an optimization on the permutation? What kind of structure are we looking for?
- What should be the starting point in higher dimensions? for other discrepancy measures?

## A separate path: the $L_2$ discrepancy

#### L<sub>2</sub> star discrepancy

For P a point set in  $[0;1]^d$ ,

$$d_2^*(P) = \left(\int_{[0,1)^d} D(q,P)^2 dq\right)^{1/2},$$

where D(q, P) is the local discrepancy.

• The main advantage of the  $L_2$  discrepancy is that it is very easy to compute using the Warnock formula [Warnock, 1972].

$$(d_2^*)^2(P) = \frac{1}{3^d} - \frac{n}{2^{d-1}} \sum_{i=1}^n \prod_{k=1}^d (1 - (x_k^{(i)})^2) + \sum_{i,j=1}^n \prod_{k=1}^d (1 - \max(x_k^{(i)}, x_k^{(j)}))$$

# Optimizing the $L_{\infty}$ discrepancy via the $L_2$ discrepancy

- Rusch et al.  $(2024)^3$  use GNN to optimize point placement for the  $L_2$  discrepancy, and they obtain excellent sets also for the  $L_{\infty}$  star discrepancy!
- Preliminary work suggests that gradient descent on a smoothed version of the  $L_2$  discrepancy also leads to similar results as long as the starting set is "good".
- A subset selection approach also leads to low  $L_{\infty}$  sets.

<sup>&</sup>lt;sup>3</sup>T. Konstantin Rusch, N. Kirk, M. M. Bronstein, C. Lemieux and D. Rus, Message-Passing Monte Carlo: Generating low-discrepancy point sets via Graph Neural Networks, 2024

### A good surrogate in low dimensions

- The "base"  $L_2$  discrepancy can be used to optimize for  $L_{\infty}$  only for  $d \le 5$ .
- Strongest results in dimension 2.
- Generalized discrepancy is the better choice in higher dimensions, but not a miracle solution.

What would be the appropriate  $L_2$  surrogate function?

### A greedy approach: the Kritzinger sequence

#### Kritzinger, 2022

Given a starting point  $p_1$ , we define the sequence  $P = (p_i)_{i \in \mathbb{N}}$ , such that

$$p_k := \arg\min_{p \in [0,1)^d} d_2^* (P_{k-1} \cup \{p\}),$$

where  $P_{k=1}$  is the set containing the first k-1 elements of P.

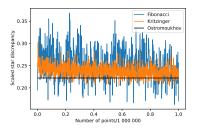
In 1d, this comes down to finding

$$\arg\min_{p\in[0,1)}(n+1)(1-p^2)+(1-p)+2\sum_{i=1}^n(1-\max(x_i,p))$$

### Constructing the sequence

- In one dimension, the next point in the sequence can only come from a "small" set of rationals.
- Linear-time to compute the next point in the sequence.
- Dimensions 2 and 3: exact MILP models, or heuristics.

### A million points



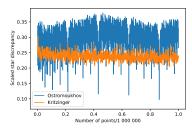


Figure: One million points with the Kritzinger sequence, compared to the Fibonacci sequence and the Ostromoukhov sequence.

# 5 million points!

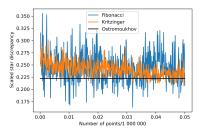
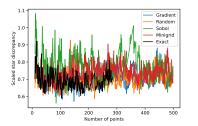


Figure: Five million points for the Kritzinger and Fibonacci sequences.



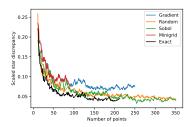


Figure: Performance of the Krizinger sequence in two and three dimensions.

### The Kritzinger sequence

#### The Kritzinger sequence:

- has on average lower discrepancy than the best low-discrepancy sequences.
- is more stable.
- corrects a bad starting set! Even starting with 1 000 badly placed points, the sequence will become low-discrepancy very quickly.

It has not been shown that the sequence is low-discrepancy!

## Key questions

- Can we determine what permutations lead to low-discrepancy sets? Are there specific structures we should look out for?
- Why do near-optimal  $L_2$  discrepancy sets have near-optimal  $L_{\infty}$  sets?
- What function should be used to generalize this to higher dimensions?
- What makes the Kritzinger sequence so good? Can we show it is low-discrepancy?

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Thank you for your attention!