

## Textbook problems

### Section 4.1:

8) b) prove  $\forall n \in \mathbb{N} \quad 6 \mid n^3 - n$

Base case: if  $n=1$  then it is true that  $6 \mid 0$

since  $0 = 0 \cdot 6$

Inductive step: assume  $6 \mid k^3 - k$ , that is  $k^3 - k = 6m$

for some  $m \in \mathbb{N}$ ; then  $(k+1)^3 - (k+1) =$

$$= k^3 + 3k^2 + 3k + 1 - k - 1 = k^3 - k + 3(k^2 + k) = 6m + 3(k^2 + k) \quad (*)$$

Claim:  $k^2 + k$  is even; proved at the end.

Using the claim  $k^2 + k = 2t$  for some  $t \in \mathbb{N}$

and  $(k+1)^3 - (k+1) = 6m + 3 \cdot 2t$  (from  $(*)$ )

$$= 6(m+t) \quad \text{so} \quad 6 \mid (k+1)^3 - (k+1).$$

Therefore if  $6 \mid k^3 - k$  then  $6 \mid (k+1)^3 - (k+1)$

Proof of claim:

if  $k$  is even  $k^2$  is even (proved in class) therefore  $k^2 + k$  is even

if  $k$  is odd  $k = 2p + 1$  for some  $p \in \mathbb{N}$

$$\text{and } k^2 + k = (2p+1)^2 + (2p+1) = 4p^2 + 6p + 2$$

$$= 2(2p^2 + 3p + 1) \quad \text{so } k^2 + k \text{ is even}$$

15) c) d) Conjecture  $\forall n \geq 1 \forall a \in \mathbb{R} \frac{d^n}{dx^n} (e^{ax}) = a^n e^{ax}$

Proof

Base case: if  $n=1$  and  $a \in \mathbb{R} \frac{d}{dx} e^{ax} = a e^{ax}$

Inductive step: assume  $\frac{d^k}{dx^k} (e^{ax}) = a^k e^{ax}$  for some  $k \in \mathbb{N}$

$$\text{then } \frac{d^{k+1}}{dx^{k+1}} (e^{ax}) = \frac{d}{dx} \left( \frac{d^k}{dx^k} e^{ax} \right) = \frac{d}{dx} a^k e^{ax} = a^k \cdot a e^{ax} = a^{k+1} e^{ax}$$

18 c)  $P(1)$  is true and the base case is proved correctly

The Inductive step correctly proves  $P(k) \Rightarrow P(k+1)$  if  $k \geq 2$

but the proof does not work for  $P(1) \Rightarrow P(2)$

"Assume all dogs in a set of 1 dog have the same breed

and consider  $D = \{d_1, d_2\}$  to be a set of two dogs

All dogs in  $\{d_2\}$  have the same breed and all dogs in

$\{d_1\}$  have the same breed, but  $\{d_1\}$  and  $\{d_2\}$  have no

common elements so  $d_1$  and  $d_2$  may have different breed.

## Section 4.2

$$1) c) \quad \forall n \geq 3 \quad \left(1 + \frac{1}{n}\right)^n < n$$

proof

$$\text{Base case: if } n=3 \quad \left(1 + \frac{1}{3}\right)^3 = \frac{64}{27} < 3$$

$$\text{Inductive step: assume } \left(1 + \frac{1}{k}\right)^k < k$$

$$\text{then } \left(1 + \frac{1}{k+1}\right)^{k+1} = \left(1 + \frac{1}{k+1}\right) \left(1 + \frac{1}{k+1}\right)^k < \left(1 + \frac{1}{k+1}\right) \cdot \left(1 + \frac{1}{k}\right)^k$$

$$\left(\text{Since } \frac{1}{k+1} < \frac{1}{k}\right)$$

$$< \left(1 + \frac{1}{k+1}\right) k \quad (\text{by induction assumption})$$

$$= k + \frac{k}{k+1} < k+1 \quad \left(\text{since } \frac{k}{k+1} < 1\right)$$

16)  $\forall n \in \mathbb{N} \exists m \text{ odd} \exists k \in \mathbb{N} \text{ s.t. } n = m \cdot 2^k$

Proof by strong induction:

Base case: If  $n=1$  then  $m=1$  and  $k=0$  and  $1=1 \cdot 2^0$

Inductive step: assume the statement is true for  $n=1, 2, \dots, t$

for some  $t \in \mathbb{N}$ , and consider  $t+1$

If  $t+1$  is odd then  $m=t+1$  and  $k=0$  and  $t+1 = \underbrace{(t+1)}_m \cdot 2^0$

If  $t+1$  is even then  $t+1 = 2\ell$  for some  $\ell \in \mathbb{N}$

and we can use the induction assumption on  $\ell$ :

$\ell = 2^k \cdot m$  with  $k \geq 0$   $m$  odd, therefore

$$t+1 = 2\ell = 2^{k+1} \cdot m$$

Now we need to prove  $m, k$  are unique that is

$$n = 2^k m = 2^q p \Rightarrow (k=q) \wedge (m=p) \quad \text{where } k, q \text{ are}$$

non negative integers and  $m$  and  $p$  are odd

Proof: first we will show  $2^k m = 2^q p \Rightarrow k=q$

by contradiction assume  $2^k m = 2^q p$  and  $k \neq q$ .

If  $k \neq q$  then one say  $k$  is bigger than  $q$

so  $2^{k-q} m = p$  but the number on the left

is even and the number on the right is odd

which is impossible, therefore we must have

$$k=q.$$

Now we will show  $(2^k m = 2^q p) \wedge (k=q) \Rightarrow m=p$

$$2^k m = 2^k p \Leftrightarrow m=p \text{ just by algebra}$$

(A proof by strong induction would also work)

## Section 4.3

$$2 \quad b) \quad f_1 = 1$$

$$f_2 = 1$$

$$f_{n+1} = f_n + f_{n-1}$$

$$\forall n \text{ in } \mathbb{N} \quad 5 \text{ div } f_{5n}$$

$$\text{Base case: } f_3 = 2, f_4 = 3, f_5 = 5 \text{ so } 5 \text{ div } f_5$$

$$\text{Inductive step: assume } 5 \text{ div } f_{5k} \text{ for some } k$$

$$f_{5(k+1)} = f_{5k+5} = f_{5k+4} + f_{5k+3} = \underbrace{f_{5k+3} + f_{5k+2}}_{f_{5k+4}} + f_{5k+3} =$$

$$= 2(\underbrace{f_{5k+2} + f_{5k+1}}_{f_{5k+3}}) + f_{5k+2} = 2(\underbrace{f_{5k+1} + f_{5k}}_{f_{5k+2}}) + 2f_{5k+1}$$

$$= \underbrace{5f_{5k+1}}_{5 \text{ div this}} + \underbrace{3f_{5k}}_{\substack{5 \text{ div this} \\ \text{by inductive assumption}}}, \text{ since by inductive}$$

$$\text{assumption } 5 \text{ div } f_{5k}, \text{ then } 5 \text{ div } f_{5(k+1)}$$

$$d) \quad \sum_{l=1}^n f_{2l-1} = f_{2n}$$

$$\text{Base case: if } n=1 \quad \sum_{l=1}^1 f_{2l-1} = f_1 = 1 = f_{2 \cdot 1}$$

Inductive step: Assume  $\sum_{l=1}^k f_{2l-1} = f_{2k}$  for some  $k$

$$\text{in } \mathbb{N}, \text{ then } \sum_{l=1}^{k+1} f_{2l-1} = \sum_{l=1}^k f_{2l-1} + f_{2(k+1)-1} =$$

$$= f_{2k} + f_{2k+1} = f_{2k+2} = f_{2(k+1)}$$

$$11) \quad q_1 = 1$$

$$q_2 = 5$$

$$q_{n+1} = q_n + 2q_{n-1}$$

$$\forall n \text{ in } \mathbb{N} \quad q_n = 2^n + (-1)^n$$

Proof

$$\text{Base cases: if } n=1 \quad q_1 = 1 = 2^1 + (-1)^1$$

$$\text{if } n=2 \quad q_2 = 5 = 2^2 + (-1)^2$$

Inductive step: assume  $q_k = 2^k + (-1)^k$  and

$$q_{k-1} = 2^{k-1} + (-1)^{k-1} \quad \text{for some } k \geq 2$$

$$\text{then } q_{k+1} = q_k + 2q_{k-1} = 2^k + (-1)^k + 2 \cdot 2^{k-1} + 2 \cdot (-1)^{k-1}$$

$$= 2^k + 2^k + (-1)^k + 2(-1)^{k-1} = \begin{cases} 2^{k+1} + 1 & \text{if } k \text{ is odd} \\ 2^{k+1} - 1 & \text{if } k \text{ is even} \end{cases}$$

$$= 2^{k+1} + (-1)^{k+1}$$

Additional problems:

1. Prove that  $\sum_{l=1}^n \frac{1}{l(l+1)} = \frac{n}{n+1}$   
 Proof: Let  $P(n)$  stand for  $\sum_{l=1}^n \frac{1}{l(l+1)} = \frac{n}{n+1}$

Base case: if  $n=1$   $\sum_{l=1}^1 \frac{1}{l(l+1)} = \frac{1}{2}$  and  $\frac{1}{1+1} = \frac{1}{2}$ , so  $P(1)$  is true

Inductive step: Assume  $P(k)$  that is  $\sum_{l=1}^k \frac{1}{l(l+1)} = \frac{k}{k+1}$ , then

$$\sum_{l=1}^{k+1} \frac{1}{l(l+1)} = \sum_{l=1}^k \frac{1}{l(l+1)} + \frac{1}{(k+1)(k+2)} = \frac{k}{k+1} + \frac{1}{(k+1)(k+2)} = \frac{k(k+2)+1}{(k+1)(k+2)} = \frac{(k+1)^2}{(k+1)(k+2)} = \frac{k+1}{k+2}$$

therefore  $P(k+1)$  is true

2. Let  $x$  be a real number, with  $x > -1$ , and let  $P(n)$  stand for  $(1+x)^n \geq 1+nx$ . We want to prove  $\forall n \geq 0 P(n)$  by induction on  $n$ .

Base case.  $P(0)$  is  $(1+x)^0 \geq 1+0x$  which simplifies to  $1 \geq 1$ , obviously true. (Note that we start at  $n=0$ )

Induction step: prove  $P(k) \Rightarrow P(k+1)$ . Assume  $P(k)$ , prove  $P(k+1)$ . In detail this means assume:  $(1+x)^k \geq 1+kx$  (and  $x > -1$ ) and prove

$(1+x)^{k+1} \geq 1+(k+1)x$ . Start from the induction hypothesis  $(1+x)^k \geq (1+kx)$  and multiply both sides of the inequality by  $(1+x)$  (remember that  $(1+x)$  is positive, this is where we use the hypothesis  $x > -1$ ) to get  $(1+x)^k(1+x) \geq (1+kx)(1+x)$  and therefore

(\*)  $(1+x)^{k+1} \geq 1+x+kx+kx^2$  (by multiplying out the right hand side). The right hand side of (\*) is equal to  $1+(k+1)x+kx^2$  and it is greater than or equal to  $1+(k+1)x$  since  $kx^2 \geq 0$ . Therefore  $(1+x)^{k+1} \geq 1+(k+1)x$  and we are done.

3. Guess:  $3^0 + 3^1 + 3^2 + \dots + 3^n = \frac{3^{n+1}-1}{2}$ . Call this statement  $P(n)$ .  
 Proof by induction:

- $P(0)$ :  $3^0 = 1 = \frac{3^1-1}{2}$ , therefore  $P(0)$  is True.
- $P(k) \Rightarrow P(k+1)$ : Assume  $3^0 + 3^1 + 3^2 + \dots + 3^k = \frac{3^{k+1}-1}{2}$  for some  $k \geq 0$ . (this is  $P(k)$ ) then  $3^0 + 3^1 + 3^2 + \dots + 3^k + 3^{k+1} = \frac{3^{k+1}-1}{2} + 3^{k+1}$ .  
 (by the induction hypothesis)  $= \frac{3 \cdot 3^{k+1}-1}{2} = \frac{3^{k+1+1}-1}{2}$  (reading the first and last term of this chain of equalities we have  $P(k+1)$  and we are done).

4. Prove that there are  $u_{n+1}$  (the  $n$ th Fibonacci number) different ways to tile a  $1 \times n$  board using squares (i.e.  $1 \times 1$  tiles) and dominoes (i.e.  $1 \times 2$  tiles).

Proof by induction.

Base cases: if  $n = 1$  a  $1 \times 1$  board can be tiled in only 1 way by using a  $1 \times 1$  square; since  $u_2 = 1$  the statement is true for  $n = 1$ .

if  $n = 2$  a  $2 \times 1$  board can be tiled in 2 ways either by using two  $1 \times 1$  squares, or by using a  $1 \times 2$  tile; since  $u_3 = u_1 + u_2 = 2$  the statement is true for  $n = 2$ .

Induction step: assume the statement is true for  $n = k$  and  $n = k - 1$ , consider a  $1 \times (k+1)$  board: we can start tiling it either by putting down a  $1 \times 1$  square, in which case we are left with  $k$  squares to cover, that is a  $1 \times k$  board, and by induction assumption we can do it in  $u_{k+1}$  ways, or we can start by putting down a domino tile, in which case we are left with  $k-1$  squares to cover, that is a  $1 \times (k-1)$  board, and by induction assumption we can do it in  $u_k$  ways. In total we have  $u_{k+1} + u_k = u_{k+2} = u_{k+1+1}$  ways.

5. Consider the following sequence  $a_{nm}$  where

$$a_{n1} = 1 \text{ for all } n \in \mathbb{N}$$

$$a_{1m} = 0 \text{ for all } m \geq 2$$

$$a_{n+1,m+1} = a_{nm} + a_{n,m+1} \text{ for all } n, m \in \mathbb{N}$$

$$\begin{array}{ccccccc} a_{1m} & \dots & 1 & 0 & 0 & 0 & \dots \\ a_{2m} & \dots & 1 & 1 & 0 & 0 & \dots \\ a_{3m} & \dots & 1 & 2 & 1 & 0 & \dots \end{array}$$

$$\text{Prove that } \forall n \in \mathbb{N} (x+y)^n = \sum_{i=0}^n a_{n+1,i} x^{n-i} y^i$$

In order to prove this we first need to prove

$$\forall n \quad a_{nn} = 1 \wedge \forall m \quad m > n \Rightarrow a_{nm} = 0$$

Prove by induction

If  $n=1$   $a_{11} = 1$  and  $a_{1m} = 0$  for  $m > 1$  by definition of  $a_{nm}$

Induction step: assume for some  $n$   $a_{nn} = 1$  and  $a_{nm} = 0$  if  $m > n$

then  $a_{n+1,n+1} = a_{nn} + a_{n,n+1} = 1 + 0 = 1$  and if  $m > n+1$  then  $m-1 > n$  so

$$a_{n+1,m} = a_{n,m-1} + a_{n,m} = 0$$



$$\begin{matrix} a_{1m} & \dots & 1 & 0 & 0 & 0 & \dots \\ a_{2m} & \dots & 1 & 1 & 0 & 0 & \dots \\ a_{3m} & \dots & 1 & 2 & 1 & 0 & \dots \end{matrix}$$

Proof of  $(x+y)^n = \sum_{l=0}^n a_{n+1, l+1} x^{n-l} y^l$  by induction

Base case: if  $n=1$   $(x+y) = \sum_{l=0}^1 a_{2, l+1} x^{1-l} y^l = a_{2,1} x + a_{2,2} y = x+y$

Induction step: assume  $(x+y)^k = \sum_{l=0}^k a_{k+1, l+1} x^{k-l} y^l$

then  $(x+y)^{k+1} = (x+y)(x+y)^k = (x+y) \sum_{l=0}^k a_{k+1, l+1} x^{k-l} y^l =$

$$\sum_{l=0}^k a_{k+1, l+1} x^{k+1-l} y^l + \sum_{l=0}^k a_{k+1, l+1} x^{k-l} y^{l+1}; \quad x^{k-l} y^{l+1} = x^{k+1-(l+1)} y^{l+1}$$

$$= \sum_{l=0}^k a_{k+1, l+1} x^{k+1-l} y^l + \sum_{l=1}^{k+1} a_{k+1, l} x^{k+1-l} y^l$$

$$= \underbrace{a_{k+1, 1}}_{a_{k+2, 1}} x^{k+1} + \sum_{l=1}^k (\underbrace{a_{k+1, l+1}}_{a_{k+2, l+1}} + \underbrace{a_{k+1, l}}_1) x^{k+1-l} y^l + \underbrace{a_{k+1, k+1}}_1 y^{k+1} = a_{k+2, k+2} y^{k+1}$$

$$\sum_{l=0}^{k+1} a_{k+2, l+1} x^{k+1-l} y^l$$

(If the algebra above is confusing you can try to write down

all the sums so  $\sum_{l=0}^k a_{k+1, l+1} x^{k+1-l} y^l + \sum_{l=0}^k a_{k+1, l+1} x^{k-l} y^{l+1} =$

$$a_{k+1, 1} x^{k+1} + a_{k+1, 2} x^k y + a_{k+1, 3} x^{k-1} y^2 + \dots + a_{k+1, k+1} x y^k +$$

$$+ a_{k+1, 1} x^k y + a_{k+1, 2} x^{k-1} y^2 + \dots + a_{k+1, k} x y^k + a_{k+1, k+1} y^{k+1}$$