

Chapter 8

Sparse Recovery

- High Dimensional Signal Recovery
- Signal recovery based on M^* bound
- exact recovery based on escape
- deterministic conditions for recovery
- Recovery with noise (LASSO/BPDN)

Problem: Fix a vector $x_{\#} \in \mathbb{R}^n$ and we get to see $y \in \mathbb{R}^m$ satisfying

$$y = Ax_{\#} + w$$

where $w \in \mathbb{R}^m$ is an error vector.

Notation

- $x_{\#}$ is called the signal
- y is called the measurement
- A is called the design matrix
- w is called the noise

$$\begin{matrix} m \\ | \\ y \end{matrix} = \begin{matrix} n \\ | \\ A \end{matrix} \begin{matrix} n \\ | \\ x_{\#} \end{matrix} + \begin{matrix} m \\ | \\ w \end{matrix}$$

We'll be interested in setting

under additional prior information $x_{\#} \in T$
 $m < n$

For the time being assume $w=0$.

Summary:

Having access to y , find $x \in \mathbb{R}^n$ s.t.
~~(*)~~ $y = Ax$ and $x \in T$.

Thm: Suppose $A \in \mathbb{R}^{m \times n}$ has independent,
isotropic, σ -subGaussian rows. Then any
point \hat{x} satisfying ~~(*)~~ satisfies

$$\mathbb{E} \|\hat{x} - x_{\#}\|_2 \leq \frac{C\sigma^2 \omega(T)}{\sqrt{m}}$$

pf: Since $\hat{x}, x_{\#} \in T$ satisfy $A\hat{x} = Ax_{\#} = y$,
we know $\hat{x} \in x_{\#} + \ker(A)$. So m^* bound

$$\begin{aligned} \mathbb{E} \|\hat{x} - x_{\#}\|_2 &\leq \mathbb{E} \text{diam}((x_{\#} + \ker(A)) \cap T) \\ &\leq \frac{C\sigma^2 \omega(T)}{\sqrt{m}} \quad \square \end{aligned}$$

Remark: So can achieve

$$\mathbb{E} \|\hat{x} - x^*\|_2 \leq \epsilon \text{diam}(T)$$

using $m \geq C \frac{\sigma^4}{\epsilon^2} d(T)$ measurements.

Often prior information encodes sparsity.

Define $\|x\|_0 = |\{i : x_i \neq 0\}|$

Suppose it's the case that

$$\|x^*\|_0 = s \ll n$$

This corresponds to setting

$$T = \{x \in \mathbb{R}^n : \|x\|_0 \leq s\}$$

Finding a point x satisfying

$$Ax = y \text{ and } x \in T \quad \textcircled{1}$$

can be difficult.

① It can be done efficiently under RTP assumptions by alternating projections

Instead, let's use a convex surrogate
Replace $\|x\|_0$ by $\|x\|_1$.

Intuition

$\lim_{p \rightarrow 0} \|x\|_p = \|x\|_0$ and $\|x\|_1$ is the first true norm.

How big should the ℓ_1 -ball be?

Lemma: $\{x \in \mathbb{R}^n : \|x\|_0 \leq s, \|x\|_2 \leq 1\} \subseteq \sqrt{s} \mathcal{B}^n$

pf: $\sum_{i=1}^n |x_i| \leq \sum_{i \in \text{supp } x} |x_i| \leq \sqrt{s} \|x\|_2 \leq \sqrt{s}$ \square

So new plan

① Find x : $y = Ax$ and $x \in \sqrt{s} \mathcal{B}^n$

Thm: Assume $\|x^*\|_0 \leq s$ and $\|x^*\|_2 \leq 1$

Then any solution of ① satisfies

$$\mathbb{E} \|\hat{x} - x^*\|_2 \leq C \sigma^2 \sqrt{\frac{s \log(n)}{m}}$$

pf: Set $T = \sqrt{s} \mathcal{B}^n$. Then $x^*, \hat{x} \in T$ and therefore

$$\mathbb{E} \|\hat{x} - x^*\|_2 \leq \underbrace{C \sigma^2 \omega(T)}_{\sqrt{m}} = C \sigma^2 \sqrt{\frac{s \log n}{m}} \quad \square$$

So recovery is possible if
 $m \sim s \log(n)$

Remark: dependence can be improved to
 $m \sim s \log(\frac{n}{s})$

Notice that (1) requires to know s and
 the theorem requires $\|x_{\#}\|_2 \leq 1$. We can do better.

Then Suppose $\|x_{\#}\|_0 \leq s$. Then any
 solution \hat{x} of

$$\begin{array}{ll} \min_x & \|x\|_1 \\ \text{s.t.} & Ax = y \end{array}$$

satisfies

$$\mathbb{E} \|\hat{x} - x_{\#}\|_2 \leq C \sqrt{\frac{s \log n}{m}} \|x_{\#}\|_2$$

p.t: We know $A\hat{x} = Ax = y$

$$\text{Also } \|x_{\#}\|_1 \leq \sqrt{s} \|x_{\#}\|_2$$

$$\Rightarrow \|\hat{x}\|_1 \leq \sqrt{s} \|x_{\#}\|_2.$$

$\hat{x}, x_{\#}$ satisfy $Ax' = y$, $x' \in \sqrt{s} \|x_{\#}\|_2 B^n$

$$\Rightarrow \mathbb{E} \|\hat{x} - x_{\#}\|_2 \leq C \sigma^2 \sqrt{\frac{s \log n}{m}} \|x_{\#}\|_2$$

The previous results extend easily to low-rank matrix recovery.

Model: There is $X \in \mathbb{R}_{\#}^{d \times d}$ satisfying

$$y_i := \langle A_i, X_{\#} \rangle \quad \text{for } i=1, \dots, m$$

where $A_i \in \mathbb{R}^{d \times d}$ are independent.

The side information we assume is $\text{rank}(X_{\#}) =: r \ll d$

Let $S(X)$ be the singular values of X .

Then $\text{rank}(X) = \|S(X)\|_2$

$$\sum_{i=1}^d s_i(X) = \|X\|_* = \|S(X)\|,$$

Lemma: $\{X : \text{rank}(X) \leq r, \|X\|_F \leq 1\} \subset \mathcal{F} B_*$

$$\text{pf: } \|X\|_* = \sum_{i=1}^r s_i(X) \leq \sqrt{r} \|S(X)\|_2 = \sqrt{r} \|X\|_F \leq \sqrt{r} \quad \square$$

Lemma: $\omega(B_\infty) \leq C\sqrt{d}$

Pf.: $\mathbb{E} \sup_{\|X\|_\infty \leq 1} \langle G, X \rangle = \mathbb{E} \|G\|_2 \leq C\sqrt{d}$ \square

Thm: Suppose that A_i are independent and $\text{vec}(A_i)$ are σ -subGaussian. Assume $X_\#$ satisfies $\|X_\#\|_F \leq 1$.

Then any \hat{X} satisfying

$y_i = \langle A_i, \hat{X} \rangle \quad \forall i=1, \dots, m, \quad \|\hat{X}\|_\infty \leq \sqrt{r}$
satisfies

$$\mathbb{E} \|\hat{X} - X_\#\|_F \leq C\sigma^2 \sqrt{\frac{rd}{m}}$$

Pf.: Set $T = \sqrt{r} B_\infty$. Then we can

treat X, \hat{X} as long vectors and then

$$\mathbb{E} \|\hat{X} - X_\#\|_F \leq C \frac{\sigma^2 \omega(T)}{\sqrt{m}} = C\sigma^2 \sqrt{\frac{rd}{m}} \quad \square$$

As in the vector case, we can instead solve

$$\begin{array}{ll} \min & \|X\|_* \\ \text{s.t.} & X = \langle A_i, X \rangle \quad \forall i=1,\dots,m. \end{array}$$

Exact Recovery Guarantees:

Consider the sparse recovery problem

$$\begin{array}{ll} \min & \|x\|_1 \\ \text{s.t.} & Ax = b \end{array}$$

We will now show that with h.p. $\hat{x} = x^*$.

Convex Optimization:

$$\begin{array}{ll} \min & f(x) \\ \text{s.t.} & Ax = b \end{array}$$

where f is convex. Then \bar{x} is optimal

$$\text{iff } \left\{ \begin{array}{l} A\bar{x} = b \\ f'(\bar{x}; v) \geq 0 \quad \forall v \in \text{ker } A \end{array} \right\}$$

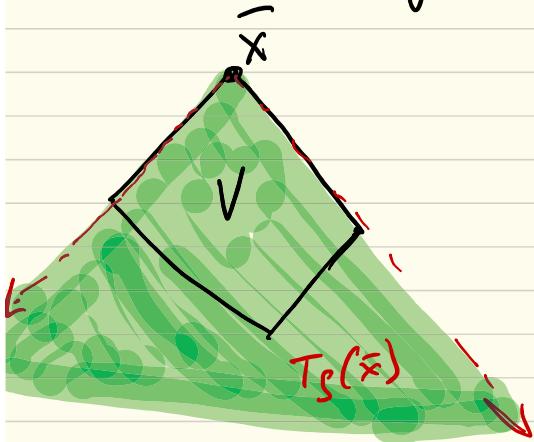
for any convex set $V \subseteq \mathbb{R}^n$ and $\bar{x} \in V$, define the tangent cone $T_V(\bar{x}) = \text{cl} \cap_{t>0} (V - \bar{x})$

Lemma:

As long as \bar{x} is not a minimizer of f ,
have

$$T_V(\bar{x}) = \{v : f'(\bar{x}, v) \leq 0\}$$

$$\text{where } V = \{y : f(y) \leq f(\bar{x})\}$$



So conditions

$$\begin{cases} A\bar{x} = b \\ T_{[f(x) \leq f(\bar{x})]}(\bar{x}) \cap \ker(A) = \{0\} \end{cases}$$

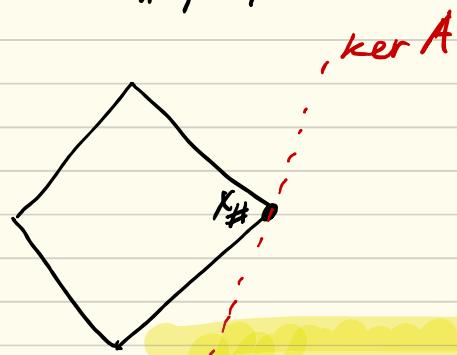
guarantees \bar{x} is a soln, and is in
fact the unique soln.

So to show that $x_{\#}$ is the unique solution of

$$\begin{array}{ll} \min & \|x\|, \\ \text{s.t.} & Ax = b \end{array}$$

we estimate probability

$$T_{\|x_{\#}\|, B} (x_{\#}) \cap \ker A = \{0\}$$



Define $Q := T_{\|x_{\#}\|, B} (x_{\#})$

Recall by the escape theorem, the probability is governed by

$$\omega(Q \cap S^{n-1})$$

Lemma: $\|h\|_1 \leq 2\sqrt{s} \|h\|_2 \quad \forall h \in Q$.

Pf: Let $S = \text{supp}(x_\#)$. Let's verify

$$\|h_{S^c}\|_1 \leq \|h_S\|_1 \quad \forall h \in Q$$

Let $x \in \|x_\#\|_1 B$, and define $h = x - x_\#$.

Then

$$\begin{aligned} \|x_\#\|_1 &\geq \|x\|_1 = \|(x_\# + h)\|_1 = \|((x_\#)_S + h_S)\|_1 + \|((x_\#)_{S^c} + h_{S^c})\|_1 \\ &\geq \|(x_\#)_S\|_1 - \|h_S\|_1 + \|h_{S^c}\|_1 \end{aligned}$$

$$\Rightarrow \|h_S\|_1 \geq \|h_{S^c}\|_1 \quad \text{ok!}$$

$$\text{So } \|h\|_1 = \|h_S\|_1 + \|h_{S^c}\|_1 \leq 2\|h_S\|_1 \leq 2\sqrt{s} \|h\|_2$$

So we learned

$$Q \cap B_2 \subseteq 2\sqrt{s} B_1$$

Thm: Suppose rows of $A \in \mathbb{R}^{m \times n}$ are independent, isotropic, σ -subGaussian. Then as long as

$$m \geq c\sigma^4 \log n$$

then with probability $1 - 2\exp(-\frac{cm}{\sigma^4})$, $x_\#$ is the unique soln of

$$\min_{x \in \mathbb{R}^n} \|x\|,$$

$$\text{s.t. } Ax = b$$

pf: The escape from shows that as long as

$$m \geq c\sigma^4 \omega^2 (Q \cap S^{n-1})$$

then

$$(Q \cap S^{n-1}) \cap \ker A = \emptyset$$

W.p. $1 - 2\exp(-c m / \sigma^4)$.

From the lemma

$$\omega(Q \cap B_2) \leq c \cdot \omega(\sqrt{\Sigma} B_1) \leq C \sqrt{s \log(n)}$$

We next look at deterministic conditions for exact recovery.

Looking back at the previous proof, define

$$\mathcal{C}(S) := \{h \in \mathbb{R}^n : \|h_{S^c}\|_1 \leq \|h_S\|\}$$

Recall we showed:

$$T_{\|x_*\|_1, B_r}(x_*) \subseteq \mathcal{C}(S)$$

where $S = \text{supp}(x_*)$.

Defn: $A \in \mathbb{R}^{m \times n}$ satisfies the restricted nullspace property with respect to

$S \subseteq \{1, \dots, n\}$, if $\mathcal{C}(S) \cap \ker(A) = \{0\}$.

Thm: The following properties are equivalent.

(a) $A \in \mathbb{R}^{m \times n}$ satisfies the restricted nullspace property with respect to S

(b) for any $x_{\#} \in \mathbb{R}^n$ with $\text{supp}(x_{\#}) = S$,
the problem

$$\min_{x \in \mathbb{R}^n} \|x\|,$$

$$\text{s.t. } Ax = b$$

has the unique soln $x_{\#}$

Pf: We already proved (a) \Rightarrow (b)

To see converse, fix $\bar{x} \in \ker A \setminus \{0\}$.

Then $\min_{x \in \mathbb{R}^n} \|x\|,$

$$\text{s.t. } Ax = A \begin{bmatrix} \bar{x}_S \\ 0 \end{bmatrix}$$

has $\begin{bmatrix} \bar{x}_S \\ 0 \end{bmatrix}$ as the unique optimal soln.

Since $A\bar{x} = 0$, it follows $\begin{bmatrix} 0 \\ -\bar{x}_{S^c} \end{bmatrix}$ is
also feasible.

$$\Rightarrow \|\bar{x}_S\|_1 < \|\bar{x}_{S^c}\|_1 \Rightarrow \bar{x} \notin C(S) \quad \square$$

Recall⁺ we also showed

$$\|h\|_1 \leq 2\sqrt{s} \|h\|_2 \quad \forall h \in C(S)$$

So $C(S) \cap B_2 \subset 2\sqrt{s} B_1$,

and escape then show that

if $A \in \mathbb{R}^{m \times n}$ is random (independent, subGaussian, isotropic rows) and

$$m \geq C \sigma^4 \sqrt{s} \log(n)$$

then w.p. $1 - 2\exp(-\frac{cm}{\sigma^4})$ have

$$C(S) \cap \ker A = \{0\}$$

Recall the $C(S)$

$$\|h\|_1 = \|h_s\|_1 + \|h_{S^c}\|_1 \leq 2\|h\|_S \leq 2\sqrt{s} \|h\|_2$$

Q: How to verify the restricted nullspace property?

Here's a simpler condition

Defn: $A \in \mathbb{R}^{m \times n}$ satisfies the restricted isometry property relative to $\{S \subseteq \{1, \dots, n\}\}$ with constant $\alpha, \beta > 0$, if

$$\alpha \|x\|_2 \leq \frac{1}{\sqrt{m}} \|Ax\|_2 \leq \beta \|x\|_2$$

for all x with $\|x\|_0 \leq s$.

Exercise: This is equivalent to requiring
{singular values of $\frac{1}{\sqrt{m}} A_S$ } $\subset [\alpha, \beta]$
where A_S is submatrix of columns S
and $|S| \leq s$.

$$\left[\begin{array}{ccc} \boxed{1} & \boxed{2} & \boxed{3} \\ \end{array} \right]$$

s-many

Thm: If $\frac{2}{3} < \alpha < \beta < \frac{4}{3}$, then A satisfies the restricted nullspace property for any $S \subseteq \{1, \dots, n\}$ with $|S| \leq s$
 [See Wainwright Proposition 7.11]

Thm: Suppose $A \in \mathbb{R}^{m \times n}$ satisfies RIP with parameters $\alpha, \beta, (1+\lambda)s$ where $\lambda > (\frac{\beta}{\alpha})^2$. Then every s -sparse $x \in \mathbb{R}^n$ can be recovered exactly by solving $\min_{\substack{x \\ Ax=b}} \|x\|_1$

Pf: Define $h = \hat{x} - x$ where \hat{x} solves $(*)$. Let $S_0 = \text{supp}(x)$, let S_1 index the largest elements of $h_{S_0^c}$ in absolute value, let S_2 index the next largest in $h_{S_0^c}$, and so on.

Set $S_{o,1} = S_o \cup S_1$.

Since $Ah = Ax - Ax_o = 0$, we have

$$0 = \|Ah\|_2 \geq \|A_{S_{o,1}} h_{S_{o,1}}\|_2 - \|A_{S_o^c} h_{S_o^c}\|_2$$

Since $|S_{o,1}| \leq s + \lambda s$, RIP gives

$$\|A_{S_{o,1}} h_{S_{o,1}}\|_2 \geq \sqrt{\alpha} \cdot \|h_{S_{o,1}}\|_2$$

and

$$\begin{aligned} \|A_{S_o^c} h_{S_o^c}\|_2 &\leq \sum_{i \geq 2} \|A_{S_i} h_{S_i}\|_2 \\ &\leq \beta \sqrt{m} \sum_{i \geq 2} \|h_{S_i}\|_2 \end{aligned}$$

So

$$\boxed{\alpha \|h_{o,1}\|_2 \leq \beta \sum_{i \geq 2} \|h_{S_i}\|_2}$$

Notice

$$\{ \text{each entry of } h_{S_j} \} \leq \frac{\| h_{S_{j-1}} \|}{\lambda^s}$$

for $j \geq 2$. [For simplicity assume]
 n divides λ^s

$$\text{So } \| h_{S_j} \|_2 \leq \frac{1}{\sqrt{\lambda^s}} \| h_{S_{j-1}} \|,$$

$$\Rightarrow \sum_{i \geq 2} \| h_{S_i} \|_2 \leq \frac{1}{\sqrt{\lambda^s}} \sum_{i \geq 1} \| h_i \|,$$

$$= \frac{1}{\sqrt{\lambda^s}} \| h_{S_0} \|,$$

$$\leq \frac{1}{\sqrt{\lambda^s}} \| h_{S_0} \|, \quad [C(S) \text{ contains tangent cone}]$$

$$\leq \frac{1}{\sqrt{\lambda}} \| h_{S_0} \|_2$$

$$\leq \frac{1}{\sqrt{\lambda}} \| h_{S_{0,1}} \|_2$$

$$\Rightarrow \frac{\beta}{\sqrt{\lambda}} \| h_{S_{0,1}} \|_2 \geq \alpha \| h_{S_{0,1}} \|_2 \Rightarrow \boxed{\alpha = 0}$$

Then let $A \in \mathbb{R}^{m \times n}$ have independent, isotropic, σ -subGaussian rows. Assume

$$m \geq \frac{\sigma^4 s}{\delta^2} \log(\epsilon n/s)$$

Then with probability $1 - 2 \exp(-\frac{c\delta^2 m}{\sigma^4})$ have

$$(1-\delta) \|x\|_2 \leq \frac{1}{\sqrt{m}} \|Ax\|_2 \leq (1+\delta) \|x\|_2 \quad \forall x \text{ with } \|x\| \leq s.$$

p.f.: we will control singular values of $A_{S'}$ uniformly over all S' with $|S'| = s$.

For any fixed S' , we know

$$\sqrt{m} - r \leq \gamma_n(A_{S'}) \leq \gamma_1(A_{S'}) \leq \sqrt{m} + r$$

with probability $1 - 2 \exp(-t^2)$ where $r = c\sigma^2(\sqrt{s} + t)$. Ensure

$$\frac{r}{\sqrt{m}} \leq \delta \text{ by setting } t = \frac{\delta \sqrt{m}}{c\sigma^2} - \sqrt{s}$$

Then $1 - \delta \leq \gamma_n(A_{S'}) \leq \gamma_1(A_{S'}) \leq 1 + \delta \quad \forall S' : |S'| = s$

w.p.

$$\begin{aligned} 1 - 2 \left(\frac{n}{s} \right) \exp(-t^2) &\leq 1 - 2 \exp(s \log(\frac{\epsilon n}{s}) - t^2) \\ &\leq 1 - 2 \exp\left(-\frac{c\delta^2 m}{\sigma^4}\right) \quad \square \end{aligned}$$

Sparse Recovery in the noisy setting

Suppose now that $x_{\#}$ satisfies

$$y = Ax_{\#} + w \quad \text{and} \quad \|x_{\#}\|_0 \leq s$$

and the noise vector $w \in \mathbb{R}^m$ is nonzero.

Three Problem Formulations: $S = \text{supp}(x_{\#})$

① LASSO ("Least Absolute Shrinkage and Selection Operator")

$$\hat{x}_a \in \arg \min_x \left\{ \frac{1}{2m} \|y - Ax\|_2^2 + \lambda_m \|x\|_1 \right\}$$

② Basis Pursuit DeNoise (BPDN I)

$$\hat{x}_b \in \arg \min_x \frac{1}{2m} \|y - Ax\|_2^2 \quad \text{subject to} \quad \|x\|_1 \leq R$$

③ Basis Pursuit DeNoise (BPDN II)

$$\hat{x}_c \in \arg \min_x \|x\|_1 \quad \text{subject to} \quad \frac{1}{2m} \|y - Ax\|_2^2 \leq b^2$$

Remark: Solutions are the same for appropriate choices of (λ_m, R, b) but one can not easily set R from b etc.

We will need to assume a modified restricted nullspace property. Define

$$\mathcal{C}_\lambda(S) := \{ h \in \mathbb{R}^d : \|h_{S^c}\|_1 \leq \lambda \|h_S\|_1 \}$$

for some $\lambda \geq 0$. Note $T_{\|x_{\#}\|_1, B_1(x_{\#})} \subset \mathcal{C}_\lambda(\text{supp}(x_{\#}))$

Defn: $A \in \mathbb{R}^{m \times n}$ satisfies the restricted eigenvalue (RE) condition with parameters (λ, γ) if

$$\frac{1}{m} \|Ah\|_2^2 \geq \gamma \|h\|_2^2 \quad \forall h \in \mathcal{C}_\lambda(S)$$

Remark:

$(\lambda, 1)$ is the restricted nullspace property.

Intuition: $D^2 \left(\frac{1}{2m} \|b - Ax\|_2^2 \right) = \frac{1}{m} A^T A$ singular

RE ensures D^2 is positive definite in important directions.

Thm: Suppose $A \in \mathbb{R}^{m \times n}$ has independent, isotropic, δ -subGaussian rows. Then

$$\frac{1}{\sqrt{m}} \|Ah\|_2 \geq \left(1 - \delta - C\delta^2(1+\lambda) \sqrt{\frac{\log n}{m}}\right) \|h\|_2 \quad \forall h \in \mathbb{R}^n$$

w.p. $1 - 2\exp(-\frac{C\delta^2 m}{64})$

Pf: Matrix Deviation with $T = \mathcal{C}_\alpha(\mathcal{S}) \cap \mathcal{S}^{n-1}$,

$$\sup_{x \in T} \left| \|Ax\|_2 - \sqrt{m} \|x\|_2 \right| \leq C\delta^2 (\omega(T) + \kappa) \quad (*)$$

w.p. $1 - 2\exp(-u^2)$.

Observe for $h \in T$ have

$$\|h\|_1 = \|h_S\|_1 + \|h_{S^c}\|_1 \leq (1+\lambda) \|h_S\|_1 \leq (1+\lambda) \sqrt{\frac{n}{m}} \|h_S\|_2 \leq 1$$

$$\Rightarrow \omega(T) \leq C(1+\lambda) \sqrt{\log(n)}$$

$$\text{So } \frac{1}{\sqrt{m}} \|Ax\|_2 \geq 1 - C\delta^2 \left((1+\lambda) \sqrt{\frac{\log n}{m}} + \frac{\kappa}{\sqrt{m}} \right)$$

w.p. $1 - 2\exp(-u^2)$. Set $\kappa = \frac{8\sqrt{m}}{C\delta^2}$ □

Thm: Suppose $\text{supp}(x_{\#}) = S'$ with $\|x_{\#}\|_0 \leq \zeta$,
and A satisfies the RE condition (x3)

(a) If $\lambda_m \geq 2 \left\| \frac{A^T w}{m} \right\|_{\infty}$, then

$$\|\hat{x}_a - x_{\#}\|_2 \leq \frac{3}{\sqrt{\kappa}} \sqrt{\zeta} \lambda_m \approx \frac{6}{\sqrt{\kappa}} \sqrt{\zeta} \left\| \frac{A^T w}{m} \right\|_{\infty}$$

(b) If $R = \|x_{\#}\|_1$, then

$$\|\hat{x}_b - x_{\#}\|_2 \leq \frac{4}{\sqrt{\kappa}} \sqrt{\zeta} \left\| \frac{A^T w}{m} \right\|_{\infty}$$

(c) If $b^2 \geq \frac{\|w\|_2^2}{2m}$, then

$$\|\hat{x}_c - x_{\#}\|_2 \leq \frac{4}{\sqrt{\kappa}} \sqrt{\zeta} \left\| \frac{A^T w}{m} \right\|_{\infty} + \frac{2}{\sqrt{\kappa}} \sqrt{b^2 - \frac{\|w\|_2^2}{2m}}$$

In all cases $\|\hat{x} - x_{\#}\|_1 \leq 4\sqrt{\zeta} \|\hat{x} - x_{\#}\|_2$

Example: Suppose $w \sim N(0, \sigma^2)$ and A is Gaussian

(a) Set $\lambda_m \approx \sigma \left(\sqrt{\frac{\log n}{m}} + \delta \right)$

$$\Rightarrow \|\hat{x}_1 - x_{\#}\|_2 \approx \frac{\sigma \cdot \sqrt{\zeta}}{\sqrt{\kappa}} \left(\sqrt{\frac{\log n}{m}} + \delta \right)$$

W.P. $1 - 2 \exp\left(\frac{-m\delta^2}{2}\right)$

(b), (c) are similar. λ and b are easy to estimate.

Example: Suppose $A_{ij} \sim N(0, 1)$ and $\|w\|_2 \leq b$

Same guarantees for LASSO:

pf: (i) Let $R = \|X_{\#}\|_1$ and let \hat{x} be a minimizer. Then

$$\frac{1}{2m} \|y - A\hat{x}\|_2^2 \leq \frac{1}{2m} \|y - Ax_{\#}\|_2^2$$

Plugging in $y = Ax_{\#} + w$

$$\Rightarrow \frac{\|Ah\|_2^2}{m} \leq \frac{2w^T Ah}{m} \quad \text{where } h = x - x_{\#}.$$

Since $h \in T_{RB_j}(x_{\#})$ we know $h \in C_{\alpha}(S')$
 $\subseteq C_{\alpha}(S)$

$$\Rightarrow \lambda \|h\|_2^2 \leq 2\|h\|_1 \cdot \left\| \frac{A^T w}{m} \right\|_{\infty} \leq 2\sqrt{s} \|h\|_2 \left\| \frac{A^T w}{m} \right\|_{\infty} \quad \square$$

(ii) Observe $\frac{1}{2m} \|y - Ax_{\#}\|_2^2 = \frac{\|w\|_2^2}{2m} \leq b^2$

so $x_{\#}$ is feasible $\Rightarrow \|\hat{x}\|_1 \leq \|x_{\#}\|_1$

$$\Rightarrow h = \hat{x} - x_{\#} \in C_{\alpha}(S).$$

$$\text{Observe } \frac{1}{2m} \|y - A\hat{x}\|_2^2 \leq b^2 = \frac{1}{2m} \|y - Ax_{\#}\|_2^2 + \left(b - \frac{\|w\|_2}{m} \right)^2$$

Rearranging

$$\frac{\|Ah\|_2^2}{m} \leq 2 \frac{w^T Ah}{m} + 2 \left(b^2 - \frac{\|w\|^2}{2m} \right)$$

$$\Rightarrow \lambda_m \|h\|_2^2 \leq 4 \sqrt{5} \|h\|_2 \left\| \frac{Aw}{m} \right\|_\infty + 2 \left(b^2 - \frac{\|w\|^2}{2m} \right)$$

$$\leq \max \left\{ 4 \sqrt{5} \|h\|_2 \left\| \frac{Aw}{m} \right\|_\infty, 2 \left(b^2 - \frac{\|w\|^2}{2m} \right) \right\}$$

(@) Define

$$L(x, \lambda_m) = \frac{1}{2m} \|y - Ax\|_2^2 + \lambda_m \|x\|_1$$

We know

$$y = Ax_{\#} + w$$

$$L(\hat{x}, \lambda_m) \leq L(x_{\#}, \lambda_m) = \frac{1}{2} \|w\|_2^2 + \lambda_m \|x_{\#}\|_1$$

$$\Rightarrow \frac{1}{2m} \|Ah\|_2^2 \leq \frac{w^T Ah}{m} + \lambda_m (\|x_{\#}\|_1 - \|\hat{x}\|_1) \quad (*)$$

where $h = \hat{x} - x_{\#}$. We can write

$$\underbrace{\|x_{\#}\|_1 - \|\hat{x}\|_1}_{\leq \|x_{\#}\|_1 - \|x_{\#} + h_S\|_1 - \|h_S\|_1} \leq \|h_S\|_1 - \|h_{Sc}\|_1$$

So

$$0 \leq \frac{1}{m} \|Ah\|_2^2 \leq 2\|h\|_1 \left\| \frac{Aw}{m} \right\|_\infty + 2\lambda_m (\|h_S\|_1 - \|h_{Sc}\|_1)$$

$$\text{choice } \frac{1}{\lambda_m} \leq \lambda_m (\|h\|_1 + 2\|h_S\|_1 - 2\|h_{Sc}\|_1) \leq \lambda_m (3\|h_S\|_1 - \|h_{Sc}\|_1)$$

So $h \in \mathcal{C}_3(S)$. RE condition gives

$$\underbrace{\kappa \|h\|_2^2 \leq 3\lambda_m \sqrt{s} \|h\|_2}_{\square}$$

In fact we proved bounds on MSE. Namely for \hat{x} :

$$\begin{aligned} \frac{1}{2m} \|A(\hat{x} - x_\#)\|_2^2 &\leq \frac{c_s \sqrt{s}}{\kappa} \|h\|_2 \left\| \frac{A^T \omega}{m} \right\|_\infty \\ &\leq \frac{c_s s}{\kappa^2} \left\| \frac{A^T \omega}{m} \right\|_\infty^2 \end{aligned}$$

So if $\omega \sim N(0, I)$, get

$$\approx \frac{c_s \log(d)}{\kappa^2 m}$$

w.h.p.

$$\frac{1}{2m} \|y - A\hat{x}\|_2^2 + \lambda_m \|\hat{x}\|_1 \leq \frac{1}{2m} \|w\|_2^2 + \lambda_m \|\hat{x}\|_1$$

$\|A\hat{x}\|_2 + w$

$$\frac{1}{2m} \|w - A(\hat{x} - x_{\#})\|_2^2 + \lambda_m \|\hat{x}\|_1$$

$\|h\|$

$$\frac{1}{2m} \|w\|_2^2 - \frac{w^T A h}{m} + \frac{1}{2m} \|Ah\|_2^2 + \lambda_m \|\hat{x}\|_1$$

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