

Lecture 5 Chernoff inequality & random graphs

One downside of Hoeffding's inequality is that it involves σ^2 but not the actual variance. A stronger inequality is possible for Binomial(n, p)

Thm (Chernoff) Let $X_i \sim \text{Ber}(p_i)$ be independent and set $S_n = \sum_{i=1}^n X_i$, $\mu = \mathbb{E}S$. Then

$$P(S_n > t) \leq e^{-\mu} \left(\frac{e\mu}{t}\right)^t \quad (\star)$$

Let's simplify this. Recall

$$\text{Var}(S_n) = \sum_{i=1}^n \text{Var}(X_i) = \sum_{i=1}^n p_i(1-p_i) \approx \mu \quad \text{for } p_i \in (0, \frac{2}{3})$$

So let's set $t = (1+\delta)\mu$ and then

$$\begin{aligned} \log(\star) &= \log\left(e^{-\mu} \left(\frac{e}{1+\delta}\right)^{(1+\delta)\mu}\right) \\ &= -\mu - (1+\delta)\mu \log(1+\delta) \\ &= \mu \left[-1 - (1+\delta) \log(1+\delta) \right] \end{aligned}$$

Define the function $h(s) = (1+s) \log(1+s)$

$$\text{Then } h'(s) = 1 + \log(1+s)$$

$$h''(s) = \frac{1}{1+s} \geq \frac{1}{2} \quad \forall s \in (0,1)$$

So the mean value theorem says that for any $s \geq 0$ $\exists \hat{s} \in [0, s]$ s.t.

$$\begin{aligned} h(s) &= h(0) + h'(0)s + \frac{1}{2} s^2 \underbrace{h''(\hat{s})}_{\geq \frac{1}{2}} \\ &\geq s + \frac{s^2}{4} \end{aligned}$$

So $\log(\star) \leq -\frac{\mu s^2}{4}$, that is

$$P(S_n - \mu \geq \delta \mu) \leq \exp\left(-\frac{\mu \delta^2}{4}\right)$$

for small deviations $s \in (0,1)$.

A similar argument gives the two sided bound

$$P(|S_n - \mu| \geq \delta \mu) \leq 2 \exp\left(-\frac{\mu \delta^2}{4}\right)$$

Pf of Chernoff's inequality:

Proceed as usual

$$P(S_n \geq t) \leq \frac{\prod_i \mathbb{E} \exp(\lambda X_i)}{\exp(\lambda t)}$$
$$= \frac{\prod_i [1 + p_i(\exp(\lambda) - 1)]}{\exp(\lambda t)}$$

The inequality $1 + x \leq e^x$ gives

$$1 + p_i(\exp(\lambda) - 1) \leq \exp(p_i(\exp(\lambda) - 1))$$
$$\leq \exp(\underbrace{\sum_i p_i}_{\mu} (e^\lambda - 1)) \cdot e^{-\lambda t}$$

$$= \exp(\mu(e^\lambda - 1) - \lambda t)$$

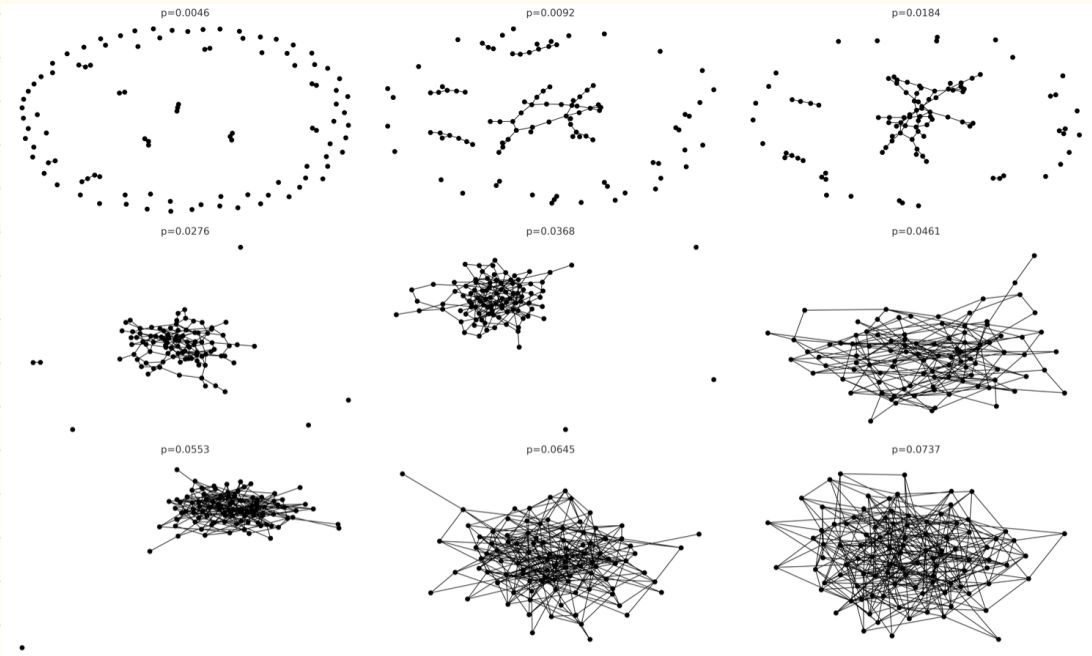
Minimizing in λ gives $\lambda = \ln\left(\frac{t}{\mu}\right)$

$$\leq e^{-\mu \left(\frac{e\mu}{t}\right)^t}$$

as claimed \square

Ex (Regularity of random graphs)

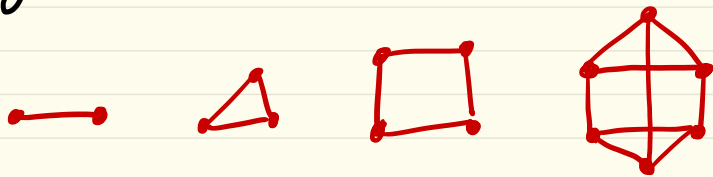
An Erdős-Rényi $G(n, p)$ is a random graph on n vertices obtained by adding each edge independently with probability p .



Some facts: As $n \rightarrow \infty$, the following are true.

- ① $np < 1 \Rightarrow$ every connected component has size $O(\log(n))$
- ② $np = 1 \Rightarrow \exists$ connected component of size $n^{2/3}$
- ③ $np \rightarrow c > 1 \Rightarrow \exists$ one connected component of size $\Omega(n)$ and all other components have size $O(\log(n))$
- ④ $np < (1-\epsilon) \log(n) \rightarrow \exists$ isolated vertices
- ⑤ $np > (1+\epsilon) \log(n) \rightarrow$ graph is connected

Def: A graph is d -regular if $\deg(i) = d$ for all i .



Thm: For any $\delta \in (0, 1)$, if $n(p-1) \geq \frac{8}{\delta^2} \log(2n)$, then $G(n, p)$ is almost d -regular with $d \triangleq (n-1)p$, meaning

$$P[(1-\delta)d \leq \deg(i) \leq (1+\delta)d \quad \forall i] \geq 1 - \frac{1}{2n}$$

pf: For each i , we may write

$$\deg(i) = \sum_{j \neq i} \mathbb{1}_{ij \in E}$$

where E is the set of edges

So $\mathbb{E}[\deg(i)] = n-1$ and Chernoff gives

$$P[|\deg(i) - d| > \delta d] \leq 2 \exp(-\delta^2 d/4)$$

Taking a union bound gives

$$P[|\deg(i) - d| \leq \delta d \quad \forall i] \quad \textcircled{A}$$

$$\geq 1 - 2n \exp(-\delta^2 d/4)$$

$$= 1 - \exp(\log(2n) - \frac{\delta^2 d}{4})$$

$$\geq 1 - \exp(-\log(2n)) = 1 - \frac{1}{2n} \quad \square$$