

# Minimization of convex compositions

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Joint work with  
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- $g: \mathbf{R}^d \rightarrow \overline{\mathbf{R}}$  is closed, convex.
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(Burke '85, Cartis-Gould-Toint '11, Fletcher '82, Lewis-Wright '15, Nesterov '06, Powell '84, Wright '90, Yuan '83, ...)

# Examples

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## Examples:

- Additive composition:

$$\min_x g(x) + c(x)$$

- Least-norm misfit:  $\min_{x \in Q} \|c(x)\|$

- ▶ Robust Phase Retrieval:  $\min_x \sum_{i=1}^m |(a_i^T x)^2 - b_i|$

- ▶ Nonneg. Factorization:  $\min_{X, Y \geq 0} \|XY^T - D\|$

# Prox-linear algorithm

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**Big assumption:**  $x^+$  is computable (for now)

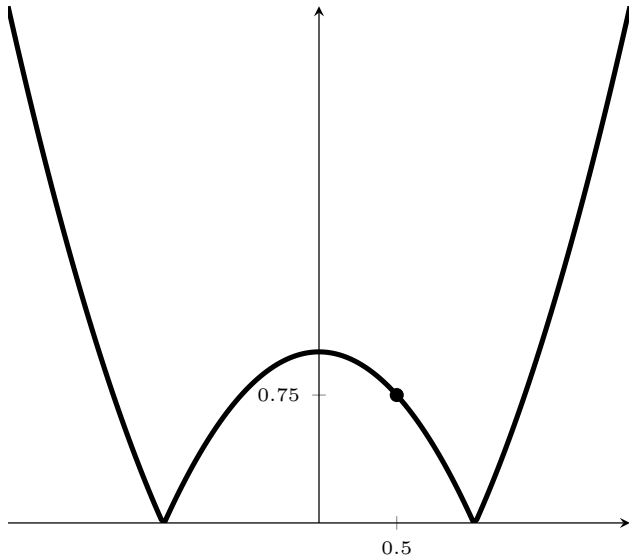


Figure:  $f(x) = |x^2 - 1|$

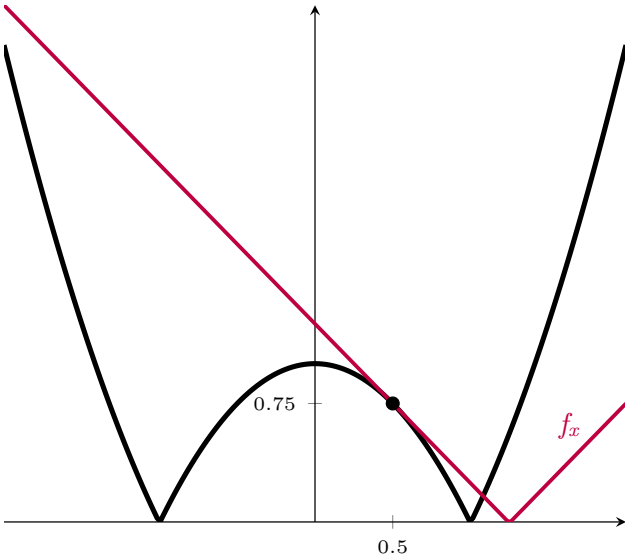


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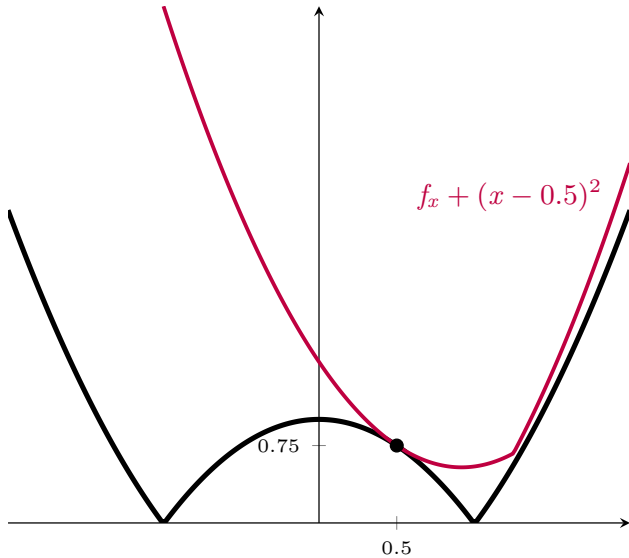


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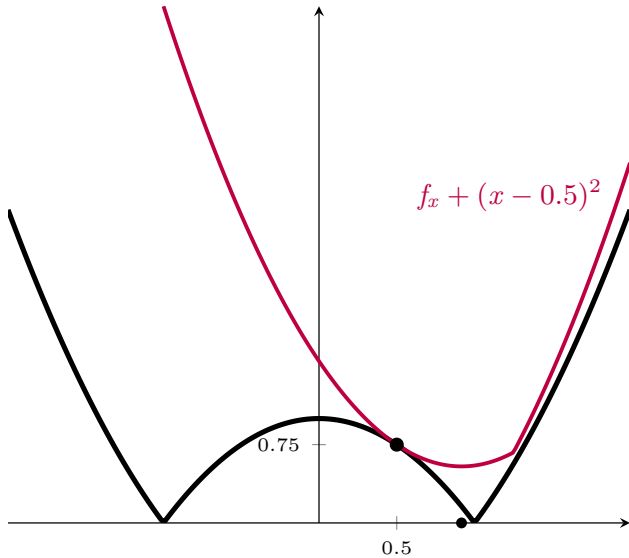


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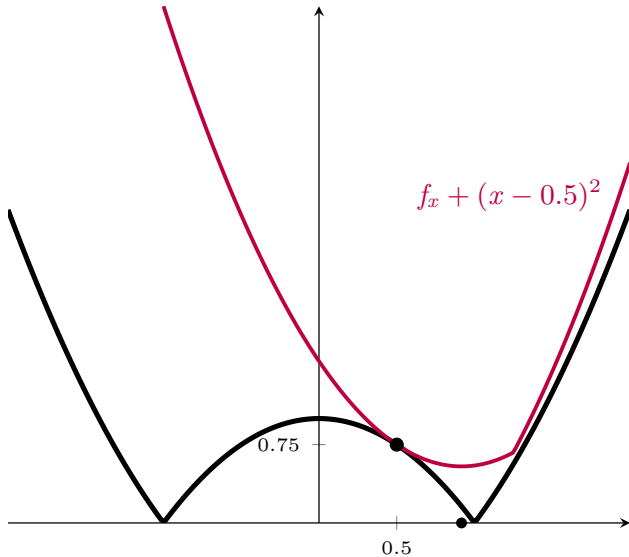


Figure:  $f(x) = |x^2 - 1|$

► No finite termination

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Prox-gradient:

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**Philosophy (Nesterov '13):**  $x^+ = x - \frac{1}{\beta}\mathcal{G}(x)$  and so

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**Sublinear rate:**

$$\|\mathcal{G}(x_k)\| < \epsilon \quad \text{after} \quad \mathcal{O}\left(\frac{\beta}{\epsilon^2} \cdot (F(x_0) - F^*)\right) \quad \text{iterations}$$

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►  $d, m$  are large  $\implies x^+$  inexactly computable

# First-order complexity

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## First-order complexity of a method:

Max # of

- ▶  $c(x), \nabla c(x)v, \nabla c(x)^T v,$
- ▶  $\text{prox}_{th}(x), \text{prox}_{tg}(x)$

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**Thm:** (D-Paquette '16) Define

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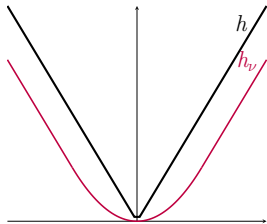
Strategy:

- ▶ smoothing of  $h$ ,
- ▶ inexact prox-linear method,
- ▶ fast-gradient subsolve.

# Smoothing

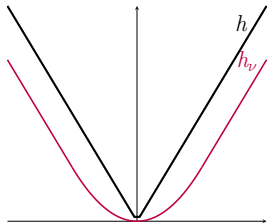
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**Lemma:** For  $\nu > 0$ , consider

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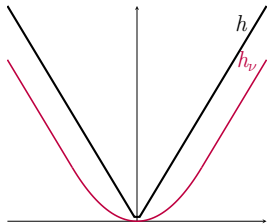
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Therefore set  $\frac{\varepsilon}{2} = \sqrt{\frac{\beta\nu}{2}} \implies \nu = \frac{\varepsilon^2}{2\beta}.$

## Inexact prox-linear

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**Input:** A point  $x_0 \in \text{dom } g$ , a method  $\mathcal{M}$

**Step k:** ( $k \geq 1$ )

Consider: 
$$\min_y \varphi(y) := F_{x_k}(y) + \frac{\beta}{2} \|y - x_k\|^2.$$

Suppose  $\mathcal{M}$  is linearly convergent from  $y_0 = x_k$ :

$$\varphi(y_i) - \varphi(y^*) \leq \gamma(1 - \tau)^i \|y_0 - y^*\|^2 \quad \forall i = 1, 2, \dots$$

Set  $x_{k+1} = y_T$  where  $T := \left\lceil \frac{1}{\tau} \log(4\gamma/\beta) \right\rceil$ .

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[motivated by [Lin-Mairal-Harchaoui '16](#)]

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Lin-Mairal-Harchaui '15, Shalev-Shwartz '15,  
Frostig-Ge-Kakade-Sidford '15, Lan '15, Allen-Zhu '16

What does  $\|\mathcal{G}(x)\| < \epsilon$  mean?

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Additive composite  $F = c + g$  setting:

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$$\implies \text{dist}((x, 0); \text{gph } \partial F) \leq \|\mathcal{G}(x)\|.$$

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1) Error Bound property:

$$\alpha \cdot \text{dist}(x; S) \leq \text{dist}(0; \partial F(x)) \quad \text{for } x \in \mathcal{X}$$

2) Tilt-stability: (Poliquin-Rockafellar '98)

$$v \mapsto \underset{x \in \mathcal{X}}{\text{argmin}} \{F(x) - \langle v, x \rangle\} \quad \text{is } \frac{1}{\alpha}\text{-Lipschitz near } \bar{v} = 0$$

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3) **Sharpness:** (Burke-Ferris '93)

$$\alpha \leq \text{dist}(0; \partial F(x)) \quad \text{for noncritical } x \in \mathcal{X}$$

$$\iff F(x) \geq F^* + \alpha \cdot \text{dist}(x, S_0) \quad \text{for } x \in \mathcal{X}.$$

Regularity	Convergence Guarantee
Error-bound	$\frac{F(x_{k+1})-F^*}{F(x_k)-F^*} \leq 1 - \left(\frac{\alpha}{\beta}\right)^2$
Tilt-stability	$\frac{F(x_{k+1})-F^*}{F(x_k)-F^*} \leq 1 - \frac{\alpha}{\beta}$
Sharpness	$\ x_{k+1} - x^*\  \leq \mathcal{O}(\ x_k - x^*\ ^2)$

(Nesterov '06, D-Lewis '15)



## Recent related work

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“Accelerated” prox-linear method (D.-Paquette '16):

$$\min_{j=1,\dots,k} \|\mathcal{G}(x_j)\|^2 \leq (\beta M)^2 \cdot \mathcal{O}\left(k^{-3} + \rho \cdot k^{-1}\right),$$

assuming  $M = \text{diam}(\text{dom } g)$  finite and  $\rho \in (0, 1)$  satisfies

$$F_x(y) \leq F(y) + \rho \cdot \frac{\beta}{2} \|y - x\|^2 \quad \forall x, y$$

Sampling methods (Duchi-Ruan '16): Almost sure conv. on

$$\min_x g(x) + \int_{\mathcal{S}} h(c(x, s), s) dP(s).$$

Inner-outer subgradient methods w/ rates (Davis-Grimmer '17)

Robust Phase Retrieval (Duchi-Ruan '17):

Quadratic convergence w.h.p on

$$\min_x \frac{1}{m} \sum_{i=1}^m |(a_i^T x)^2 - b_i|$$

[uses Eldar-Mendelson '12]

## Summary

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“Efficiency of minimizing compositions of convex functions and smooth maps”, D-Paquette, 2016, arXiv:1605.00125.

“Error bounds, quadratic growth, and linear convergence of proximal methods”, D-Lewis, Math. Oper. Res. 2017.

- ▶ prox-linear method
- ▶ first-order complexity
- ▶ finite-sum problems
- ▶ termination conditions
- ▶ local linear/quadratic rates

Thank you!