
Linear Analysis
Lecture 17

The QR for Algorithm for the Schur Factorization

The QR algorithm is used to compute a specific Schur unitary triangularization of a matrix $A \in \mathbb{C}^{n \times n}$.

The algorithm is **iterative**. That is, we generate a sequence

$$A = A_0, A_1, A_2, \dots$$

of matrices that are unitarily similar to A . The goal is to get the subdiagonal elements to converge to zero, since then the eigenvalues will appear on the diagonal.

If A is Hermitian, then so are A_1, A_2, \dots , so if the subdiagonal elements $\rightarrow 0$, also the superdiagonal elements converge to 0, and (in the limit) we have diagonalized A .

Variations of the QR algorithm are the most commonly used methods for computing all the eigenvalues (and eigenvectors if wanted) of a matrix. It behaves well numerically since all the similarity transformations are unitary.

Upper-Hessenberg Form

When used in practice, a matrix is first reduced to **upper-Hessenberg form**

$$\begin{bmatrix} * & \cdots & * & * \\ * & \cdots & * & * \\ 0 & \ddots & & \vdots \\ 0 & \cdots & * & * \end{bmatrix}$$

($h_{ij} = 0$ for $i > j + 1$) using unitary similarity transformations built from Householder reflections (or Givens rotations), quite analogous to computing a QR factorization.

However, since similarity transformations are being performed, we require left and right multiplication by the Householder transformations — leading to an inability to zero out the first subdiagonal ($i = j + 1$) in the process.

If A is Hermitian and upper-Hessenberg, A is tridiagonal.

This initial reduction decreases the computational cost of the QR algorithm. It is successful because the upper-Hessenberg form is preserved by the iterations: if A_k is upper Hessenberg, so is A_{k+1} .

There are many variants of the QR algorithm. We consider the basic algorithm over \mathbb{C} .

The Basic QR Algorithm

Given $A \in \mathbb{C}^{n \times n}$, let $A_0 = A$. For $k = 0, 1, 2, \dots$, starting with A_k , do a QR factorization of A_k :

$$A_k = Q_k R_k.$$

The set

$$A_{k+1} = R_k Q_k .$$

Remark

$$R_k = Q_k^H A_k \quad \text{so} \quad A_{k+1} = Q_k^H A_k Q_k$$

is unitarily similar to A_k . In general,

$$A_{k+1} = Q_k^H Q_{k-1}^H \dots Q_0^H A_0 Q_0 \dots Q_{k-1} Q_k .$$

Thus

$$\|A_k\| = \|A_0\| \quad \forall k = 1, 2, \dots .$$

The algorithm uses the Q of the QR factorization of A_k to perform the next unitary similarity transformation.

Convergence of the QR Algorithm

We now show convergence. Under the hypotheses that $A \in \mathbb{C}^{n \times n}$ has eigenvalues $\lambda_1, \dots, \lambda_n$ with $|\lambda_1| > |\lambda_2| > \dots > |\lambda_n| > 0$.

Lemma. Let $Q_j \in \mathbb{C}^{n \times n}$ $j = 1, 2, \dots$ be unitary matrices, and let $R_j \in \mathbb{C}^{n \times n}$ $j = 1, 2, \dots$ be upper triangular matrices with positive diagonal entries. Suppose

$$Q_j R_j \rightarrow I \quad \text{as } j \rightarrow \infty .$$

Then $Q_j \rightarrow I$ and $R_j \rightarrow I$.

Proof Sketch. Let Q_{j_k} be any subsequence of Q_j . Since the set of unitary matrices in $\mathbb{C}^{n \times n}$ is compact, there exists a sub-subsequence $Q_{j_{k_l}}$ and a unitary Q such that $Q_{j_{k_l}} \rightarrow Q$. So

$$R_{j_{k_l}} = Q_{j_{k_l}}^H Q_{j_{k_l}} R_{j_{k_l}} \rightarrow Q^H \cdot I = Q^H .$$

So Q^H is unitary, upper triangular, with nonnegative diagonal elements, which implies easily that $Q^H = I$. Thus every subsequence of Q_j has in turn a sub-subsequence converging to I . Consequently, $Q_j \rightarrow I$, and thus $R_j = Q_j^H Q_j R_j \rightarrow I \cdot I = I$. \square

Convergence of the QR Algorithm

Theorem: Suppose $A \in \mathbb{C}^{n \times n}$ has eigenvalues

$$\lambda_1, \dots, \lambda_n \quad \text{with} \quad |\lambda_1| > |\lambda_2| > \dots > |\lambda_n| > 0.$$

Choose $X \in \mathbb{C}^{n \times n}$ such that

$$X^{-1}AX = \Lambda \equiv \text{diag}(\lambda_1, \dots, \lambda_n).$$

Suppose X^{-1} has an LU decomposition. Generate the sequence

$$A = A_0, A_1, A_2, \dots$$

using the QR algorithm. Then

the subdiagonal entries of $A_k \rightarrow 0$ as $k \rightarrow \infty$,

and

the j^{th} diagonal entry $\rightarrow \lambda_j \quad 1 \leq j \leq n$.

Define

$$\tilde{Q}_k = Q_0 Q_1 \cdots Q_k \quad \text{and} \quad \tilde{R}_k = R_k \cdots R_0.$$

Then $A_{k+1} = \tilde{Q}_k^H A \tilde{Q}_k$.

Claim: $\tilde{Q}_k \tilde{R}_k = A^{k+1}$

Proof: Proceed by induction. Clearly the Claim holds for $k = 0$.

Suppose $\tilde{Q}_{k-1} \tilde{R}_{k-1} = A^k$. Then

$$R_k = A_{k+1} Q_k^H = \tilde{Q}_k^H A \tilde{Q}_k Q_k^H = \tilde{Q}_k^H A \tilde{Q}_{k-1},$$

so

$$\tilde{R}_k = R_k \tilde{R}_{k-1} = \tilde{Q}_k^H A \tilde{Q}_{k-1} \tilde{R}_{k-1} = \tilde{Q}_k^H A^{k+1},$$

so $\tilde{Q}_k \tilde{R}_k = A^{k+1}$. □

Recall

$$X^{-1}AX = \Lambda \equiv \text{diag}(\lambda_1, \dots, \lambda_n).$$

Set

$$X = QR,$$

$$X^{-1} = LU,$$

Q unitary,

L unit lower triangular,

R nonsingular upper triangular,

U nonsingular upper triangular.

Then

$$\begin{aligned} A^{k+1} &= X\Lambda^{k+1}X^{-1} \\ &= QR\Lambda^{k+1}LU \\ &= QR(\Lambda^{k+1}L\Lambda^{-(k+1)})\Lambda^{k+1}U. \end{aligned}$$

Let

$$E_{k+1} = \Lambda^{k+1}L\Lambda^{-(k+1)} - I \quad \text{and} \quad F_{k+1} = RE_{k+1}R^{-1}.$$

Claim: $E_{k+1} \rightarrow 0$ (and thus $F_{k+1} \rightarrow 0$) as $k \rightarrow \infty$.

Proof: Let ℓ_{ij} denote the elements of L . E_{k+1} is strictly lower triangular, and for $i > j$ its ij element is

$$\left(\frac{\lambda_i}{\lambda_j}\right)^{k+1} \ell_{ij} \rightarrow 0 \quad \text{since } |\lambda_i| < |\lambda_j|. \quad \square$$

Now $A^{k+1} = QR(I + E_{k+1})\Lambda^{k+1}U = Q(I + F_{k+1})R\Lambda^{k+1}U$.

Choose a QR factorization of $I + F_{k+1}$ (which is invertible)

$$I + F_{k+1} = \widehat{Q}_{k+1}\widehat{R}_{k+1} \rightarrow I,$$

where \widehat{R}_{k+1} has positive diagonal entries. By the Lemma,

$$\widehat{Q}_{k+1} \rightarrow I \quad \text{and} \quad \widehat{R}_{k+1} \rightarrow I.$$

Since

$$A^{k+1} = Q(\widehat{Q}_{k+1}\widehat{R}_{k+1})R\Lambda^{k+1}U = (Q\widehat{Q}_{k+1})(\widehat{R}_{k+1}R\Lambda^{k+1}U)$$

and $A^{k+1} = \widetilde{Q}_k\widetilde{R}_k$, the essential uniqueness of QR factorizations of invertible matrices implies there exists a unitary diagonal matrix D_k for which

$$Q\widehat{Q}_{k+1}D_k^H = \widetilde{Q}_k \quad \text{and} \quad D_k\widehat{R}_{k+1}\Lambda^{k+1}U = \widetilde{R}_k.$$

So $\tilde{Q}_k D_k = Q \hat{Q}_{k+1} \rightarrow Q$, and thus

$$D_k^H A_{k+1} D_k = D_k^H \tilde{Q}_k^H A \tilde{Q}_k D_k \rightarrow Q^H A Q.$$

But

$$\begin{aligned} Q^H A Q &= Q^H (X \Lambda X^{-1}) Q \\ &= Q^H (Q R \Lambda X^{-1}) Q R R^{-1} \\ &= R \Lambda R^{-1} \end{aligned}$$

is upper triangular with diagonal entries $\lambda_1, \dots, \lambda_n$ in that order. Since D_k is unitary and diagonal, the diagonal and lower triangular entries of $R \Lambda R^{-1}$ and of $D_k R \Lambda R^{-1} D_k^H$ are the same, namely $\text{diag}[\lambda_1, \dots, \lambda_n]$ and

$$\|A_{k+1} - D_k R \Lambda R^{-1} D_k^H\| = \|D_k^H A_{k+1} D_k - R \Lambda R^{-1}\| \rightarrow 0.$$

The Theorem follows. □

Comments on the Proof

Note that the proof shows that there is a sequence $\{D_k\}$ of unitary diagonal matrices for which $D_k^H A_{k+1} D_k \rightarrow R \Lambda R^{-1}$. So although the superdiagonal ($i < j$) elements of A_{k+1} may not converge, the magnitude of each superdiagonal element converges.

As a partial explanation for why the QR algorithm works, we show how the convergence of the first column of A_k to $[\lambda_1, 0, \dots, 0]$ follows from the power method.

Suppose $A \in \mathbb{C}^{n \times n}$ is diagonalizable and has a unique e-value $\lambda_1 > 0$ of maximum modulus with unit e-vec x_1 . Then if $x \in \mathbb{C}^n$ has $\langle x, x_1 \rangle \neq 0$, then $A^k x / \|A^k x\| \rightarrow x_1 / \|x_1\|$.

If X^{-1} has an LU factorization, the $(1, 1)$ entry of X^{-1} is nonzero. Thus when e_1 is expanded in terms of the eigenvectors

x_1, \dots, x_n (cols. of X), the x_1 -coefficient is nonzero. So

$$A^{k+1} e_1 / \|A^{k+1} e_1\| \rightarrow \alpha x_1 \quad \text{for some } \alpha \in \mathbb{C} \quad \text{with } |\alpha| = 1.$$

Let $(\tilde{q}_k)_1$ denote the first column of \tilde{Q}_k and $(\tilde{r}_k)_{11}$ denote the $(1, 1)$ -entry of \tilde{R}_k . Then

$$A^{k+1} e_1 = \tilde{Q}_k \tilde{R}_k e_1 = (\tilde{r}_k)_{11} \tilde{Q}_k e_1 = (\tilde{r}_k)_{11} (\tilde{q}_k)_1,$$

so $(\tilde{q}_k)_1 \rightarrow \alpha x_1$. Since $A_{k+1} = \tilde{Q}_k^H A \tilde{Q}_k$, the first column of A_{k+1} converges to $[\lambda_1, 0, \dots, 0]$.