

Linear Least-Squares Problems

1. LINEAR LEAST-SQUARES AS AN OPTIMIZATION PROBLEM

Let $A \in \mathbb{R}^{m \times n}$ and $b \in \mathbb{R}^m$ and assume that $m \gg n$, i.e., m is much greater than n . In this setting it is highly unlikely that there exists a vector $x \in \mathbb{R}^n$ such that $Ax = b$. As an alternative goal, we try to find the x that is as *close* to solving $Ax = b$ as possible. But first we must define a notion of *close*. One way is to try to find the vector x that minimizes the norm of the residual error $\|Ax - b\|_2$. That is, we wish to find a vector \bar{x} such that

$$\|A\bar{x} - b\|_2 \leq \|Ax - b\|_2 \quad \forall x \in \mathbb{R}^n.$$

Equivalently, we wish to solve the optimization problem

$$\mathcal{LLS} \quad \min_{x \in \mathbb{R}^n} \frac{1}{2} \sum_{i=1}^m \left(\sum_{j=1}^n A_{ij}x_j - b_i \right)^2 = \frac{1}{2} \|Ax - b\|_2^2.$$

If we set

$$f(x) = \frac{1}{2} \|Ax - b\|_2^2,$$

then the first-order necessary conditions for a point x to solve \mathcal{LLS} is that $\nabla f(x) = 0$. In order to use this fact, we need to compute an expression for the gradient of f .

For $i = 1, 2, \dots, m$, set $\phi_i(x) = \left(\sum_{j=1}^n A_{ij}x_j - b_i \right)^2$, then $f(x) = \frac{1}{2} \sum_{i=1}^m \phi_i(x)$. Observe that for $i \in \{1, 2, \dots, m\}$ and a given $j_0 \in \{1, 2, \dots, n\}$

$$\frac{\partial}{\partial x_{j_0}} \phi_i(x) = \frac{\partial}{\partial x_{j_0}} \left(\sum_{j=1}^n A_{ij}x_j - b_i \right)^2 = 2A_{ij_0} \left(\sum_{j=1}^n A_{ij}x_j - b_i \right).$$

Consequently,

$$\frac{\partial}{\partial x_{j_0}} f(x) = \sum_{i=1}^m A_{ij_0} \left(\sum_{j=1}^n A_{ij}x_j - b_i \right) = A_{j_0}^T (Ax - b),$$

and so

$$\nabla f(x) = \begin{bmatrix} \frac{\partial}{\partial x_1} f(x) \\ \frac{\partial}{\partial x_2} f(x) \\ \vdots \\ \frac{\partial}{\partial x_n} f(x) \end{bmatrix} = \begin{bmatrix} A_{j_0}^T (Ax - b) \\ A_{j_0}^T (Ax - b) \\ \vdots \\ A_{j_0}^T (Ax - b) \end{bmatrix} = A^T (Ax - b).$$

Therefore, if $x \in \mathbb{R}^n$ solves \mathcal{LLS} , then it must be the case that $0 = \nabla f(x) = A^T (Ax - b)$ or equivalently

$$A^T Ax = A^T b.$$

This system of equations are called the *normal equations* for the linear least-squares problem \mathcal{LLS} .

We must now address the question of whether there exists a solution to the normal equations. For this we make use of the following lemma.

Lemma 1. *For every matrix $A \in \mathbb{R}^{m \times n}$ we have*

$$\text{Null}(A^T A) = \text{Null}(A) \quad \text{and} \quad \text{Ran}(A^T A) = \text{Ran}(A^T).$$

Proof. Note that if $x \in \text{Null}(A)$, then $Ax = 0$ and so $A^T Ax = 0$, that is, $x \in \text{Null}(A^T A)$. Therefore, $\text{Null}(A) \subset \text{Null}(A^T A)$. Conversely, if $x \in \text{Null}(A^T A)$, then

$$A^T Ax = 0 \implies x^T A^T Ax = 0 \implies (Ax)^T (Ax) = 0 \implies \|Ax\|_2^2 = 0 \implies Ax = 0,$$

or equivalently, $x \in \text{Null}(A)$. Therefore, $\text{Null}(A^T A) \subset \text{Null}(A)$, and so $\text{Null}(A^T A) = \text{Null}(A)$.

Since $\text{Null}(A^T A) = \text{Null}(A)$, the Fundamental Theorem of the Alternative tells us that

$$\text{Ran}(A^T A) = \text{Ran}((A^T A)^T) = \text{Null}(A^T A)^\perp = \text{Null}(A)^\perp = \text{Ran}(A^T),$$

which proves the lemma. \square

Let us now examine the existence of solutions to the the normal equations in light of this lemma. The normal equations are $A^T A x = A^T b$. By definition, $A^T b \in \text{Ran}(A^T)$. The lemma tells us that $\text{Ran}(A^T) = \text{Ran}(A^T A)$. Hence, there must exist x such that $A^T A x = A^T b$, that is, the normal equations are always consistent regardless of the choice of matrix $A \in \mathbb{R}^{m \times n}$ and vector $b \in \mathbb{R}^m$.

Theorem 2. *The normal equations are consistent for all $A \in \mathbb{R}^{m \times n}$ and $b \in \mathbb{R}^m$.*

This tells us that the linear least-squares problem \mathcal{LLS} always has a critical point. But it does not tell us when these critical points for \mathcal{LLS} are global solutions to \mathcal{LLS} . In this regard, we have the following surprising result.

Theorem 3. *Let $A \in \mathbb{R}^{m \times n}$ and $b \in \mathbb{R}^m$. Then every solution \bar{x} to $A^T A x = A^T b$ satisfies*

$$(1) \quad \|A\bar{x} - b\|_2 \leq \|Ax - b\|_2 \quad \forall x \in \mathbb{R}^n,$$

that is \bar{x} is a global solution to \mathcal{LLS} .

Proof. Given $u, v \in \mathbb{R}^m$, we have

$$(2) \quad \|u + v\|_2^2 = (u + v)^T(u + v) = u^T u + 2u^T v + v^T v = \|u\|_2^2 + 2u^T v + \|v\|_2^2.$$

Consequently, for every $x \in \mathbb{R}^n$,

$$\begin{aligned} \|Ax - b\|_2^2 &= \|(Ax - A\bar{x}) + (A\bar{x} - b)\|_2^2 \\ &= \|A(x - \bar{x})\|_2^2 + 2(A(x - \bar{x}))^T(A\bar{x} - b) + \|A\bar{x} - b\|_2^2 && \text{(by (2))} \\ &\geq 2(x - \bar{x})^T A^T(A\bar{x} - b) + \|A\bar{x} - b\|_2^2 && \text{(since } \|A(x - \bar{x})\|_2^2 \geq 0) \\ &= \|A\bar{x} - b\|_2^2 && \text{(since } A^T(A\bar{x} - b) = 0), \end{aligned}$$

or equivalently, (1) holds. \square

So far we know that the normal equations are consistent and that every solution to the normal equations solves the linear least-squares problem. That is, a solution to the linear least-squares problem always exists. We now address the question of when the solution is unique. This is equivalent to asking when the normal equations have a unique solution. From our study of nonsingular matrices, we know this occurs precisely when the matrix $A^T A$ is nonsingular or equivalently, invertible, in which case the unique solution is given by $\bar{x} = (A^T A)^{-1} A^T b$. Note that $A^T A$ is invertible if and only if $\text{Null}(A^T A) = \{0\}$. But, by Lemma 1, this is equivalent to $\text{Null}(A) = \{0\}$.

Theorem 4. *The linear least-squares problem \mathcal{LLS} has a unique solution if and only if $\text{Null}(A) = \{0\}$.*

2. ORTHOGONAL PROJECTION ONTO A SUBSPACE

In the previous section we stated the linear least-squares problem as the optimization problem \mathcal{LLS} . We can view this problem in a somewhat different light as a least distance problem to a subspace, or equivalently, as a projection problem for a subspace. Suppose $S \subset \mathbb{R}^m$ is a given subspace and $b \notin S$. The least distance problem for S and b is to find that element of S that is as close to b as possible. That is we wish to solve the problem

$$(3) \quad \min_{z \in S} \frac{1}{2} \|z - b\|_2^2.$$

The solution is the point $\bar{z} \in S$ such that

$$\|\bar{z} - b\|_2 \leq \|z - b\|_2 \quad \forall z \in S.$$

If we now take the subspace to be the range of A , $S = \text{Ran}(A)$, then the problem (3) is closely related to the problem \mathcal{LLS} since

$$(4) \quad \text{if } \bar{z} \text{ solves (3) and } \bar{x} \text{ solves } \mathcal{LLS}, \text{ then } \bar{z} = A\bar{x} \text{ (why?).}$$

Below we discuss the connection between the notion of a projection matrix and solutions to (3). Since the norm $\|\cdot\|_2$ is generated by the dot product, $\|w\|_2 = \sqrt{w \bullet w}$, least norm problems of this type are solved using the notion of *orthogonal projection onto a subspace*.

To understand orthogonal projections, we must first introduce the notion of projection. A matrix $P \in \mathbb{R}^{m \times m}$ is said to be a projection if $P^2 = P$. In this case we say that P is a projection onto the range of P , $S = \text{Ran}(P)$. Note that if $x \in \text{Ran}(P)$, then there is a $w \in \mathbb{R}^m$ such that $x = Pw$, therefore, $Px = P(Pw) = P^2w = Pw = x$. That is, P leaves all elements of $\text{Ran}(P)$ fixed. Also, note that, if P is a projection, then

$$(I - P)^2 = I - P - P + P^2 = I - P,$$

and so $(I - P)$ is also a projection. Since for all $w \in \mathbb{R}^m$,

$$w = Pw + (I - P)w,$$

we have

$$\mathbb{R}^m = \text{Ran}(P) + \text{Ran}((I - P)).$$

In this case we say that the subspaces $\text{Ran}(P)$ and $\text{Ran}((I - P))$ are *complementary* since their sum is the whole space and their intersection is the origin, i.e., $\text{Ran}(P) \cap \text{Ran}((I - P)) = \{0\}$ (why?).

Conversely, given any two subspaces S_1 and S_2 such that $S_1 \cap S_2 = \{0\}$ and $S_1 + S_2 = \mathbb{R}^m$, there is a projection P such that $S_1 = \text{Ran}(P)$ and $S_2 = \text{Ran}((I - P))$. We do not show how to construct these projections here, but note that they can be generated with the aid of bases for S_1 and S_2 .

This relationship between projections and complementary subspaces allows us to define a notion of *orthogonal projection*. For every subspace $S \subset \mathbb{R}^m$, we know that

$$S \cap S^\perp = \{0\} \quad \text{and} \quad S + S^\perp = \mathbb{R}^m \quad (\text{why?}).$$

Therefore, there is a projection P such that $\text{Ran}(P) = S$ and $\text{Ran}((I - P)) = S^\perp$, or equivalently,

$$(5) \quad ((I - P)y)^T(Pw) = 0 \quad \forall y, w \in \mathbb{R}^m.$$

The orthogonal projection plays a very special role among all possible projections onto a subspace. For this reason, we denote the orthogonal projection onto the subspace S by P_S .

We now use the condition (5) to derive a simple test of whether a projection is an orthogonal projection. For brevity, we write $P := P_S$ and set $M = (I - P)^T P$. Then, by (5),

$$M_{ij} = e_i^T M e_j = 0 \quad \forall i, j = 1, \dots, n,$$

i.e., M is the zero matrix. But then,

$$P = P^T P = (P^T P)^T = P^T.$$

Conversely, if $P = P^T$ and $P^2 = P$, then $(I - P)^T P = 0$. Therefore, a matrix P is an orthogonal projection if and only if $P^2 = P$ and $P = P^T$. An orthogonal projection for a given subspace S can be constructed from any orthonormal basis for that subspace. Indeed, if the columns of the matrix Q form an orthonormal basis for S , then the matrix $P = QQ^T$ satisfies

$$P^2 = QQ^T QQ^T \stackrel{\text{why}}{=} Q I_k Q^T = QQ^T = P \quad \text{and} \quad P^T = (QQ^T)^T = QQ^T = P,$$

so that P is an orthogonal projection. Moreover, since we know that $\text{Ran}(QQ^T) = \text{Ran}(Q) = S$, P is necessarily the orthogonal projector onto S .

Let us now return to the problem (5). We show that \bar{z} solves this problem if and only if $\bar{z} = P_S b$ where P_S is the orthogonal projection onto S . To see this, let $P := P_S$ and $z \in S$ so that $z = Pz$. Then

$$\begin{aligned} \|z - b\|_2^2 &= \|Pz - Pb - (I - P)b\|_2^2 \\ &= \|P(z - b) + (I - P)b\|_2^2 \\ &= \|P(z - b)\|_2^2 + 2(z - b)^T P(I - P)b + \|(I - P)b\|_2^2 \\ &= \|P(z - b)\|_2^2 + \|(I - P)b\|_2^2 \\ &\geq \|(P - I)b\|_2^2 \\ &= \|\bar{z} - b\|_2^2, \end{aligned}$$

which shows that $\|\bar{z} - b\|_2 \leq \|z - b\|_2$ for all $z \in S$.

Let us now consider the linear least-squares problem \mathcal{LLS} when $m \gg n$ and $\text{Null}(A) = \{0\}$. In this case, we have shown that $\bar{x} = (A^T A)^{-1} A^T b$ solves \mathcal{LLS} , and $\bar{z} = P_S b$ solves (5) where P_S is the orthogonal projector onto $S = \text{Ran}(A)$. Hence, by (4),

$$P_S b = \bar{z} = A\bar{x} = A(A^T A)^{-1} A^T b.$$

Since this is true for all possible choices of the vector b , we have

$$(6) \quad P_S = A(A^T A)^{-1} A^T !$$

That is, the matrix $A(A^T A)^{-1} A^T$ is the orthogonal projector onto the range of A . One can also check this directly by showing that the matrix $M = A(A^T A)^{-1} A^T$ satisfies $M^2 = M$, $M^T = M$, and $\text{Ran}(M) = \text{Ran}(A)$.

3. MINIMAL NORM SOLUTIONS TO $Ax = b$

Again let $A \in \mathbf{R}^{m \times n}$, but now we suppose that $m \ll n$. In this case A is short and fat so the matrix A most likely has rank m , or equivalently,

$$(7) \quad \text{Ran}(A) = \mathbf{R}^m .$$

In this case, regardless of the choice of the vector $b \in \mathbf{R}^m$, the set of solutions to $Ax = b$ will be infinite since the nullity of A is $n - m$. Indeed, if x^0 is any particular solution to $Ax = b$, then the set of solutions is given by $\{x^0 + z \mid z \in \text{Null}(A)\}$. In this setting, one might prefer the solution to the system having least norm. This solution is found by solving the problem

$$(8) \quad \min_{z \in S} \frac{1}{2} \|z + x^0\|_2^2 ,$$

where S is the null-space of A . This problem is of the form (3). Consequently, the solution is given by $\bar{z} = -P_S x^0$ where P_S is now the orthogonal projection onto $S := \text{Null}(A)$.

In this context, note that $(I - P_S)$ is the orthogonal projector onto $\text{Null}(A)^\perp = \text{Ran}(A^T)$. Recall that the formula (6) shows that if $M \in \mathbf{R}^{k \times s}$ is such that $\text{Null}(M) = \{0\}$, then the orthogonal projector onto $R := \text{Ran}(M)$ is given by

$$(9) \quad P_R = M(M^T M)^{-1} M^T .$$

In our case, $M = A^T$ and $M^T M = AA^T$. Our working assumption (7) implies that

$$\text{Null}(M) = \text{Null}(A^T) = \text{Ran}(A)^\perp = (\mathbf{R}^m)^\perp = \{0\}$$

and consequently, by (9), the orthogonal projector onto $R = \text{Ran}(A^T) = S^\perp$ is given by

$$P_{S^\perp} = A^T (AA^T)^{-1} A .$$

Therefore, the orthogonal projector onto S is

$$P_S = I - P_{S^\perp} = I - A^T(AA^T)^{-1}A .$$

Putting this all together, we find that the solution to (8) is

$$\bar{z} = -P_S x^0 = (A^T(AA^T)^{-1}A - I)x^0 ,$$

and the solution to $Ax = b$ of least norm is

$$\bar{x} = x^0 + \bar{z} = A^T(AA^T)^{-1}Ax^0 ,$$

where x^0 is any particular solution to $Ax = b$, i.e., $Ax^0 = b$. Plugging \bar{x} into $Ax = b$ gives

$$A\bar{x} = AA^T(AA^T)^{-1}Ax^0 = Ax^0 = b .$$