

# A bijection on balanced words reversing both des and maj

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# What are balanced words?

Let  $\mathcal{W}_{n,k}$  be the set of *balanced words*. A word  $w \in \mathcal{W}_{n,k}$  is any arrangement of the letters  $\{1, \dots, k\}$  each of which must occur  $n$  times.

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## Example

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- $|\mathcal{W}_{n,k}| = \binom{nk}{n, \dots, n}$ ,
- $\mathcal{W}_{1,k}$  are permutations of  $[k]$ , and
- $\mathcal{W}_{n,2}$  are balanced binary words or lattice paths in an  $n \times n$  grid.

# Descents and major index

A word in  $\mathcal{W}_{3,4}$  is

|       |   |   |   |   |   |   |   |   |   |    |    |    |
|-------|---|---|---|---|---|---|---|---|---|----|----|----|
|       | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 | 11 | 12 |
| $w =$ | 2 | 3 | 1 | 1 | 2 | 4 | 4 | 1 | 4 | 3  | 2  | 3  |

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|       | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 | 11 | 12 |
| $w =$ | 2 | 3 | 1 | 1 | 2 | 4 | 4 | 1 | 4 | 3  | 2  | 3  |

- The *descent set* of  $w$  is  $\text{Des}(w) = \{2, 7, 9, 10\}$ ,

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|       | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 | 11 | 12 |
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- The *descent set* of  $w$  is  $\text{Des}(w) = \{2, 7, 9, 10\}$ ,
- The number of *descents* in  $w$  is  $\text{des}(w) = 4$ ,

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|       | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 | 11 | 12 |
| $w =$ | 2 | 3 | 1 | 1 | 2 | 4 | 4 | 1 | 4 | 3  | 2  | 3  |

- The *descent set* of  $w$  is  $\text{Des}(w) = \{2, 7, 9, 10\}$ ,
- The number of *descents* in  $w$  is  $\text{des}(w) = 4$ ,
- The *major index* of  $w$  is the sum of the descent set  $\text{maj}(w) = 2 + 7 + 9 + 10 = 28$ .



# Descents and major index

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|       |   |   |   |   |   |   |   |   |   |    |    |    |
|-------|---|---|---|---|---|---|---|---|---|----|----|----|
|       | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 | 11 | 12 |
| $w =$ | 2 | 3 | 1 | 1 | 2 | 4 | 4 | 1 | 4 | 3  | 2  | 3  |

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- The number of *descents* in  $w$  is  $\text{des}(w) = 4$ ,
- The *major index* of  $w$  is the sum of the descent set  $\text{maj}(w) = 2 + 7 + 9 + 10 = 28$ .

## Definition

The bivariate generating polynomial of  $\text{des}$  and  $\text{maj}$  on balanced words is

$$C_{n,k}(x, q) = \sum_{w \in \mathcal{W}_{n,k}} x^{\text{des}(w)} q^{\text{maj}(w)}.$$

For example,

$$C_{2,2}(x, q) = 1 + xq + 2xq^2 + xq^3 + x^2q^4.$$

## Descents and major index

The *unique* word  $w \in \mathcal{W}_{n,k}$  having the minimum value of both  $\text{des}(w) = 0$  and  $\text{maj}(w) = 0$  is

$$1 \ 1 \ \dots \ 1 \ 2 \ 2 \ \dots \ 2 \ \dots \ k \ k \ \dots \ k.$$

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$$1 \ 1 \ \dots \ 1 \ 2 \ 2 \ \dots \ 2 \ \dots \ k \ k \ \dots \ k.$$

The *unique* word having the maximum value of both  $\text{des}(w) = n(k-1)$  and  $\text{maj}(w) = n^2 \binom{k}{2}$  is

$$k \ k-1 \ \dots \ 1 \ k \ k-1 \ \dots \ 1 \ \dots \ k \ k-1 \ \dots \ 1.$$

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The *unique* word  $w \in \mathcal{W}_{n,k}$  having the minimum value of both  $\text{des}(w) = 0$  and  $\text{maj}(w) = 0$  is

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For example, words in  $\mathcal{W}_{2,3}$  can have a maximum  $\text{des}$  of 4 and a maximum  $\text{maj}$  of 12.

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$$k \ k-1 \ \dots \ 1 \ k \ k-1 \ \dots \ 1 \ \dots \ k \ k-1 \ \dots \ 1.$$

For example, words in  $\mathcal{W}_{2,3}$  can have a maximum  $\text{des}$  of 4 and a maximum  $\text{maj}$  of 12.

Words in  $\mathcal{W}_{2,3}$  having  $\text{des} = 1$  and  $\text{maj} = 3$

|   |   |   |   |   |   |
|---|---|---|---|---|---|
| 1 | 1 | 3 | 2 | 2 | 3 |
| 1 | 3 | 3 | 1 | 2 | 2 |
| 1 | 2 | 3 | 1 | 2 | 3 |
| 1 | 2 | 2 | 1 | 3 | 3 |
| 2 | 3 | 3 | 1 | 1 | 2 |
| 2 | 2 | 3 | 1 | 1 | 3 |

## Descents and major index

The *unique* word  $w \in \mathcal{W}_{n,k}$  having the minimum value of both  $\text{des}(w) = 0$  and  $\text{maj}(w) = 0$  is

$$1 \ 1 \ \dots \ 1 \ 2 \ 2 \ \dots \ 2 \ \dots \ k \ k \ \dots \ k.$$

The *unique* word having the maximum value of both  $\text{des}(w) = n(k-1)$  and  $\text{maj}(w) = n^2 \binom{k}{2}$  is

$$k \ k-1 \ \dots \ 1 \ k \ k-1 \ \dots \ 1 \ \dots \ k \ k-1 \ \dots \ 1.$$

For example, words in  $\mathcal{W}_{2,3}$  can have a maximum  $\text{des}$  of 4 and a maximum  $\text{maj}$  of 12.

Words in  $\mathcal{W}_{2,3}$  having  $\text{des} = 1$  and  $\text{maj} = 3$

|   |   |   |   |   |   |
|---|---|---|---|---|---|
| 1 | 1 | 3 | 2 | 2 | 3 |
| 1 | 3 | 3 | 1 | 2 | 2 |
| 1 | 2 | 3 | 1 | 2 | 3 |
| 1 | 2 | 2 | 1 | 3 | 3 |
| 2 | 3 | 3 | 1 | 1 | 2 |
| 2 | 2 | 3 | 1 | 1 | 3 |

Words in  $\mathcal{W}_{2,3}$  having  $\text{des} = 4 - 1 = 3$  and  $\text{maj} = 12 - 3 = 9$

|   |   |   |   |   |   |
|---|---|---|---|---|---|
| 2 | 1 | 3 | 2 | 3 | 1 |
| 3 | 2 | 3 | 1 | 2 | 1 |
| 3 | 2 | 2 | 1 | 3 | 1 |
| 2 | 1 | 3 | 1 | 3 | 2 |
| 3 | 1 | 3 | 2 | 2 | 1 |
| 3 | 1 | 2 | 1 | 3 | 2 |

# The main theorem

## Theorem (Tielker 2023)

*The bivariate generating polynomial of  $\text{des}$  and  $\text{maj}$  on balanced words is palindromic, i.e.*

$$C_{n,k}(x^{-1}, q^{-1}) = x^{-n(k-1)} q^{-n^2 \binom{k}{2}} C_{n,k}(x, q).$$

# The main theorem

## Theorem (Tielker 2023)

*The bivariate generating polynomial of des and maj on balanced words is palindromic, i.e.*

$$C_{n,k}(x^{-1}, q^{-1}) = x^{-n(k-1)} q^{-n^2 \binom{k}{2}} C_{n,k}(x, q).$$

For example, the coefficients of  $C_{2,3}(x, q)$ , written as an array is the same when rotated by  $180^\circ$ .

$$\begin{pmatrix} 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 2 & 5 & 6 & 5 & 2 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & 6 & 10 & 14 & 10 & 6 & 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 2 & 5 & 6 & 5 & 2 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 \end{pmatrix},$$



# The main theorem

We prove the theorem by

- constructing a des and maj reversing bijection,  $f_2$ , on balanced binary words  $\mathcal{W}_{n,2}$ , and
- constructing the general bijection,  $f_k$ , for balanced words  $\mathcal{W}_{n,k}$  (with  $k > 2$ ) via recursion.

# Encoding balanced binary words as pairs of boolean arrays

A *boolean array* is a word in the letters T and F. For example,

$$\mu = \text{TFFTF TT}$$

is one. The *complement* of  $\mu$  is written  $\bar{\mu}$ .

$$\bar{\mu} = \text{FTTFTFF}.$$

Let  $B_n = \{\text{T}, \text{F}\}^n$  be the set of boolean arrays of length  $n$ . Also write  $B_n^k$  to mean boolean arrays of length  $n$  having  $k$  T's.

# Encoding balanced binary words as pairs of boolean arrays

Consider the following word  $w \in \mathcal{W}_{5,2}$ ,

$$w = 1\ 2\ 2\ 2\ 1\ 1\ 2\ 1\ 1\ 2.$$

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$$w = 1\ 2\ 2\ 2\ 1\ 1\ 2\ 1\ 1\ 2.$$

Group each descent into a box like so

$$w = 1\ 2\ 2\ \boxed{2\ 1}\ 1\ \boxed{2\ 1}\ 1\ 2.$$

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Group each descent into a box like so

$$w = 1\ 2\ 2\ \boxed{2\ 1}\ 1\ \boxed{2\ 1}\ 1\ 2.$$

Record the 1's outside the boxes with an F and the boxes themselves with a T,

$$\mu_w = F\ \boxed{T}\ F\ \boxed{T}\ F.$$

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Group each descent into a box like so

$$w = 1\ 2\ 2\ \boxed{2\ 1}\ 1\ \boxed{2\ 1}\ 1\ 2.$$

Record the 1's outside the boxes with an F and the boxes themselves with a T,

$$\mu_w = F\ \boxed{T}\ F\ \boxed{T}\ F.$$

Record the 2's outside the boxes with an F and the boxes themselves with a T,

$$\nu_w = F\ F\ \boxed{T}\ \boxed{T}\ F.$$

# Encoding balanced binary words as pairs of boolean arrays

We can also start with two boolean arrays of 1) the same length and 2) having the same number of T's and construct a binary word. Let,

$$\mu = F \text{ T } F \text{ T } F.$$

$$\nu = F F \text{ T } \text{ T } F.$$

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$$\mu = F \text{ T } F \text{ T } F.$$

$$\nu = F F \text{ T } \text{ T } F.$$

Expand each T into a 21 like so

$$\begin{matrix} 2 & 1 & & 2 & 1 \end{matrix}$$



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$$\mu = F \text{ T } F \text{ T } F.$$

$$\nu = F F \text{ T } \text{ T } F.$$

Expand each T into a 21 like so

$$21 \quad 21$$

Read the number of F's before the first T in  $\mu$  and write

$$1 \quad 21 \quad 21.$$

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Expand each T into a 21 like so

$$21 \quad 21$$

Read the number of F's before the first T in  $\mu$  and write

$$1 \quad 21 \quad 21.$$

Read the number of F's before the first T in  $\nu$  and place it after any 1's

$$1 \quad 2 \quad 2 \quad 21 \quad 21.$$

# Encoding balanced binary words as pairs of boolean arrays

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$$\nu = F F \text{ T } \text{ T } F.$$

Expand each T into a 21 like so

$$\text{2 1} \quad \text{2 1}$$

Read the number of F's before the first T in  $\mu$  and write

$$1 \text{ 2 1} \quad \text{2 1}.$$

Read the number of F's before the first T in  $\nu$  and place it after any 1's

$$1 \text{ 2 2} \text{ 2 1} \quad \text{2 1}.$$

Repeat the process for the second group of F's in array to get

$$1 \text{ 2 2} \text{ 2 1} \text{ 1} \text{ 2 1}.$$

# Encoding balanced binary words as pairs of boolean arrays

We can also start with two boolean arrays of 1) the same length and 2) having the same number of T's and construct a binary word. Let,

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$$\nu = F F \text{ T } \text{ T } F.$$

Expand each T into a 21 like so

$$21 \quad 21$$

Read the number of F's before the first T in  $\mu$  and write

$$1 \quad 21 \quad 21.$$

Read the number of F's before the first T in  $\nu$  and place it after any 1's

$$1 \quad 2 \quad 2 \quad 21 \quad 21.$$

Repeat the process for the second group of F's in array to get

$$1 \quad 2 \quad 2 \quad 21 \quad 1 \quad 21.$$

Finally, repeat the process for the last group to get

$$1 \quad 2 \quad 2 \quad 21 \quad 1 \quad 21 \quad 1 \quad 2.$$

# Encoding balanced binary words as pairs of boolean arrays

The process we saw just now gives us the following invertible map

$$\phi : \mathcal{W}_{n,2} \rightarrow \bigsqcup_{0 \leq k \leq n} B_n^k \times B_n^k$$

by

$$\phi(w) = (\mu_w, \nu_w).$$

# The des and maj reversing bijection on $\mathcal{W}_{n,2}$

Let  $f_2 : \mathcal{W}_{n,2} \rightarrow \mathcal{W}_{n,2}$  be given by

$$f_2(w) = \phi^{-1}(\bar{\mu}_w, \bar{\nu}_w).$$

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$$f_2(w) = \phi^{-1}(\bar{\mu}_w, \bar{\nu}_w).$$

## Theorem

*The map  $f_2$  on  $\mathcal{W}_{n,2}$  is an involution satisfying*

$$\text{des}(w) + \text{des}(f_2(w)) = n,$$

*and*

$$\text{maj}(w) + \text{maj}(f_2(w)) = n^2$$

*for all  $w \in \mathcal{W}_{n,2}$ .*

# The des and maj reversing bijection on $\mathcal{W}_{n,2}$

Consider the same word  $w \in \mathcal{W}_{5,2}$  as earlier,

$$w = 1\ 2\ 2\ \boxed{2\ 1}\ 1\ \boxed{2\ 1}\ 1\ 2.$$



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The map  $\phi$  sends this to

$$\mu_w = F\ \boxed{T}\ F\ \boxed{T}\ F.$$

$$\nu_w = F\ F\ \boxed{T}\ \boxed{T}\ F.$$

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$$\mu_w = F\ \boxed{T}\ F\ \boxed{T}\ F.$$

$$\nu_w = F\ F\ \boxed{T}\ \boxed{T}\ F.$$

The complements of these are

$$\bar{\mu}_w = \boxed{T}\ F\ \boxed{T}\ F\ \boxed{T}.$$

$$\bar{\nu}_w = \boxed{T}\ \boxed{T}\ F\ F\ \boxed{T}.$$

## The des and maj reversing bijection on $\mathcal{W}_{n,2}$

Consider the same word  $w \in \mathcal{W}_{5,2}$  as earlier,

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The map  $\phi$  sends this to

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$$\nu_w = F\ F\ \boxed{T}\ \boxed{T}\ F.$$

The complements of these are

$$\bar{\mu}_w = \boxed{T}\ F\ \boxed{T}\ F\ \boxed{T}.$$

$$\bar{\nu}_w = \boxed{T}\ \boxed{T}\ F\ F\ \boxed{T}.$$

And  $\phi^{-1}(\bar{\mu}_w, \bar{\nu}_w)$  gives

$$f_2(w) = \boxed{2\ 1}\ 1\ \boxed{2\ 1}\ 1\ 2\ 2\ \boxed{2\ 1}.$$

Clearly,  $\text{des}(w) + \text{des}(f_2(w)) = n$ .

# The des and maj reversing bijection on $\mathcal{W}_{n,2}$

The proof that  $\text{maj}(w) + \text{maj}(f_2(w)) = n^2$  would be trivial if  $\text{Des}(w) \cap \text{Des}(f_2(w)) = \emptyset$ ; however, in our example  $\text{Des}(1222112112) = \{4, 7\}$  and  $\text{Des}(2112112221) = \{1, 4, 9\}$ .

|         |   |   |   |   |     |   |     |   |   |
|---------|---|---|---|---|-----|---|-----|---|---|
| $w$     | = | 1 | 2 | 2 | 2 1 | 1 | 2 1 | 1 | 2 |
| $\mu_w$ | = | F |   |   | T   | F | T   | F |   |
| $\nu_w$ | = |   | F | F | T   |   | T   | F |   |

# The des and maj reversing bijection on $\mathcal{W}_{n,2}$

Count 4 with respect to  $\mu_w$  and  $\nu_w$ .

$$\begin{array}{rcccccccc}
 & & & & 4 & & & & \\
 w & = & 1 & 2 & 2 & \boxed{2\ 1} & 1 & \boxed{2\ 1} & 1 & 2 \\
 \mu_w & = & F^1 & & & T^1 & F & T & F & \\
 \nu_w & = & & F^1 & F^1 & T & & T & F & 
 \end{array}$$

# The des and maj reversing bijection on $\mathcal{W}_{n,2}$

Count 4 with respect to  $\mu_w$  and  $\nu_w$ .

$$\begin{array}{rcccccccc}
 & & & & 4 & & & & \\
 w & = & 1 & 2 & 2 & \boxed{2\ 1} & 1 & \boxed{2\ 1} & 1 & 2 \\
 \mu_w & = & F^1 & & & T^1 & F & T & F & \\
 \nu_w & = & & F^1 & F^1 & T & & T & F & 
 \end{array}$$

Next, count 7 with respect to  $\mu_w$  and  $\nu_w$ .

$$\begin{array}{rcccccccc}
 & & & & 4 & & 7 & & \\
 w & = & 1 & 2 & 2 & \boxed{2\ 1} & 1 & \boxed{2\ 1} & 1 & 2 \\
 \mu_w & = & F^2 & & & T^2 & F^1 & T^1 & F & \\
 \nu_w & = & & F^2 & F^2 & T^1 & & T & F & 
 \end{array}$$

# The des and maj reversing bijection on $\mathcal{W}_{n,2}$

$$\begin{array}{rcl}
 w & = & 1 \quad 2 \quad 2 \quad \boxed{2 \ 1} \quad 1 \quad \boxed{2 \ 1} \quad 1 \quad 2 \\
 \mu_w & = & F^2 \quad \quad \quad T^2 \quad F^1 \quad T^1 \quad F \\
 \nu_w & = & \quad F^2 \quad F^2 \quad T^1 \quad \quad T \quad F
 \end{array}$$

# The des and maj reversing bijection on $\mathcal{W}_{n,2}$

$$\begin{array}{rcccccccc}
 w & = & 1 & 2 & 2 & 4 & 1 & 7 & 1 & 2 \\
 & & & & & \boxed{2\ 1} & & \boxed{2\ 1} & & \\
 \mu_w & = & F^2 & & & T^2 & F^1 & T^1 & F & \\
 \nu_w & = & & F^2 & F^2 & T^1 & & T & F & 
 \end{array}$$

Similarly, compute the major index of  $f_2(w)$

$$\begin{array}{rcccccccc}
 w & = & 1 & 4 & 1 & 2 & 2 & 9 \\
 & & \boxed{2\ 1} & & \boxed{2\ 1} & & & \boxed{2\ 1} \\
 \mu_w^c & = & T^3 & F^2 & T^2 & F^1 & & T^1 \\
 \nu_w^c & = & T^2 & & T^1 & & F^1 & F^1 & T
 \end{array}$$



# The des and maj reversing bijection on $\mathcal{W}_{n,2}$

$$\begin{array}{rcccccccc}
 & & & & 4 & & 7 & & \\
 w & = & 1 & 2 & 2 & \boxed{2\ 1} & 1 & \boxed{2\ 1} & 1 & 2 \\
 \mu_w & = & F^2 & & & T^2 & F^1 & T^1 & F & \\
 \nu_w & = & & F^2 & F^2 & T^1 & & T & F & 
 \end{array}$$

Similarly, compute the major index of  $f_2(w)$

$$\begin{array}{rcccccccc}
 & & 1 & & 4 & & & & 9 & \\
 w & = & \boxed{2\ 1} & 1 & \boxed{2\ 1} & 1 & 2 & 2 & \boxed{2\ 1} & \\
 \mu_w^c & = & T^3 & F^2 & T^2 & F^1 & & & T^1 & \\
 \nu_w^c & = & T^2 & & T^1 & & F^1 & F^1 & T & 
 \end{array}$$

Tally the counts for each index of  $\mu_w$  and  $\mu_w^c$ .

$$\begin{array}{rcccccc}
 \mu_w & = & F^2 & T^2 & F^1 & T^1 & F \\
 \mu_w^c & = & T^3 & F^2 & T^2 & F^1 & T^1 \\
 & & 5 & 4 & 3 & 2 & 1
 \end{array}$$

# The des and maj reversing bijection on $\mathcal{W}_{n,2}$

$$\begin{array}{rcccccccc}
 & & & & 4 & & 7 & & \\
 w & = & 1 & 2 & 2 & \boxed{2\ 1} & 1 & \boxed{2\ 1} & 1\ 2 \\
 \mu_w & = & F^2 & & & T^2 & F^1 & T^1 & F \\
 \nu_w & = & & F^2 & F^2 & T^1 & & T & F
 \end{array}$$

Similarly, compute the major index of  $f_2(w)$

$$\begin{array}{rcccccccc}
 & & 1 & & 4 & & & & 9 \\
 w & = & \boxed{2\ 1} & 1 & \boxed{2\ 1} & 1 & 2 & 2 & \boxed{2\ 1} \\
 \mu_w^c & = & T^3 & F^2 & T^2 & F^1 & & & T^1 \\
 \nu_w^c & = & T^2 & & T^1 & & F^1 & F^1 & T
 \end{array}$$

Tally the counts for each index of  $\mu_w$  and  $\mu_w^c$ .

$$\begin{array}{rcccccc}
 \mu_w & = & F^2 & T^2 & F^1 & T^1 & F \\
 \mu_w^c & = & T^3 & F^2 & T^2 & F^1 & T^1 \\
 & & 5 & 4 & 3 & 2 & 1
 \end{array}$$

Tally the counts for each index of  $\nu_w$  and  $\nu_w^c$ .

$$\begin{array}{rcccccc}
 \nu_w & = & F^2 & F^2 & T^1 & T & F \\
 \nu_w^c & = & T^2 & T^1 & F^1 & F^1 & T \\
 & & 4 & 3 & 2 & 1 & 0
 \end{array}$$

$$\text{Hence, } \text{maj}(w) + \text{maj}(f_2(w)) = \sum_{i=1}^n i + \sum_{i=1}^n (i-1) = n^2.$$

## Ascents and comajor index

Clearly, we can modify the map  $\phi$  to encode *ascents* and the *comajor index* instead. Consider the same word  $w \in \mathcal{W}_{5,2}$  as earlier,

$$w = \boxed{1\ 2}\ 2\ 2\ 1\ \boxed{1\ 2}\ 1\ \boxed{1\ 2}.$$

The map  $\underline{\phi}$  sends this to

$$\underline{\mu}_w = \boxed{T}\ F\ \boxed{T}\ F\ \boxed{T}.$$

$$\underline{\nu}_w = \boxed{T}\ F\ F\ \boxed{T}\ \boxed{T}.$$

The complements of these are

$$\bar{\underline{\mu}}_w = F\ \boxed{T}\ F\ \boxed{T}\ F.$$

$$\bar{\underline{\nu}}_w = F\ \boxed{T}\ \boxed{T}\ F\ F.$$

And  $\underline{\phi}^{-1}(\bar{\underline{\mu}}_w, \bar{\underline{\nu}}_w)$  gives

$$\underline{f}_2(w) = 2\ 1\ \boxed{1\ 2}\ 1\ \boxed{1\ 2}\ 2\ 2\ 1.$$

Clearly,  $\text{asc}(w) + \text{asc}(\underline{f}_2(w)) = n$  and  $\text{comaj}(w) + \text{comaj}(\underline{f}_2(w)) = n^2$ .

The bijection  $f_2$  also reverses ascents and comajor index

### Theorem

For  $w \in \mathcal{W}_{n,2}$ ,

$$\text{asc}(w) + \text{asc}(f_2(w)) = n,$$

and

$$\text{comaj}(w) + \text{comaj}(f_2(w)) = n^2.$$

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$$\text{asc}(w) + \text{asc}(f_2(w)) = n,$$

and

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Ascents and descents in binary words are closely related. For any word  $w = w_1 w_2 \dots w_n$ , let  $L(w) = w_2 \dots w_n$  and  $R(w) = w_1 \dots w_{n-1}$ . In our running example, we have

$$\begin{array}{l} \mu_w = F \text{ T } F \text{ T } F, \quad \underline{\mu}_w = \text{ T } F \text{ T } F \text{ T} \\ \nu_w = F F \text{ T } \text{ T } F, \quad \underline{\nu}_w = \text{ T } F F \text{ T } \text{ T} \end{array}$$

Notice that  $\underline{\mu}_w = L(\mu) \text{ T}$  and  $\underline{\nu}_w = \text{ T } R(\nu)$ .

## The bijection $f_2$ also reverses ascents and comajor index

In general, the boolean encoding for ascents is related to the boolean encoding for descents as follows.

|             | $\mu_1 = F$                                                  | $\mu_1 = T$                                                  |
|-------------|--------------------------------------------------------------|--------------------------------------------------------------|
| $\nu_n = F$ | $\underline{\mu} = L(\mu) T$<br>$\underline{\nu} = T R(\nu)$ | $\underline{\mu} = L(\mu) T$<br>$\underline{\nu} = F R(\nu)$ |
| $\nu_n = T$ | $\underline{\mu} = L(\mu) F$<br>$\underline{\nu} = T R(\nu)$ | $\underline{\mu} = L(\mu) F$<br>$\underline{\nu} = F R(\nu)$ |

## The bijection $f_2$ also reverses ascents and comajor index

In general, the boolean encoding for ascents is related to the boolean encoding for descents as follows.

|             | $\mu_1 = F$                                                  | $\mu_1 = T$                                                  |
|-------------|--------------------------------------------------------------|--------------------------------------------------------------|
| $\nu_n = F$ | $\underline{\mu} = L(\mu) T$<br>$\underline{\nu} = T R(\nu)$ | $\underline{\mu} = L(\mu) T$<br>$\underline{\nu} = F R(\nu)$ |
| $\nu_n = T$ | $\underline{\mu} = L(\mu) F$<br>$\underline{\nu} = T R(\nu)$ | $\underline{\mu} = L(\mu) F$<br>$\underline{\nu} = F R(\nu)$ |

The proof that  $f_2$  also reverses ascents and comajor index as well follows from the fact that complementation leaves the table as it is.

# Mirror-symmetric balanced words

We now define a subset of balanced binary words called *mirror-symmetric* words. The following word is mirror-symmetric

$$w = 1\ 2\ 2\ 2\ 2\ 1\ 2\ 2\ 2\ 1\ 2\ 1\ 1\ 1\ 2\ 1\ 1\ 1\ 1\ 2.$$



# Mirror-symmetric balanced words

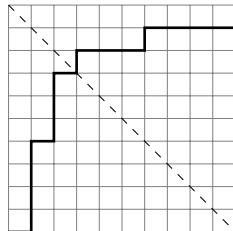
We now define a subset of balanced binary words called *mirror-symmetric* words. The following word is mirror-symmetric

$$w = 1\ 2\ 2\ 2\ 2\ 1\ 2\ 2\ 2\ 1\ 2\ 1\ 1\ 1\ 2\ 1\ 1\ 1\ 2.$$

We can render it as a lattice path by taking an east step for a 1 and a north step for a 2. The path is invariant when reflected across the NW-SE diagonal.

A balanced binary word  $w$  is mirror-symmetric if

$$w_{n+1} \dots w_{2n} = \text{rev}((w_1 \dots w_n)^c).$$



# The bijection $f_2$ preserves mirror-symmetry

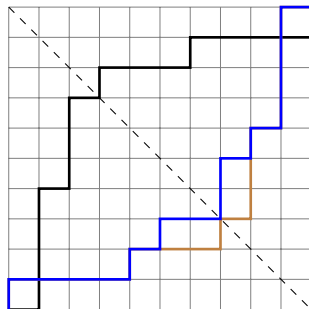
## Theorem

*Let  $w \in \mathcal{W}_{n,2}$  be mirror-symmetric. Then so is  $f_2(w)$ .*

The proof follows from the observation that

$$\mu = \text{rev}(\nu)$$

$$\nu = \text{rev}(\mu).$$



**Figure:** The black path is  $w$ , the blue path is  $f_2(w)$ , and the red path is the complement of  $w$ .

# Overview of the general bijection

For  $w \in \mathcal{W}_{n,k}$  and  $k > 2$ , we construct  $w$  in two steps.

$$w_{-k} \in \mathcal{W}_{n,k-1} \xleftarrow{g(w_{-k}, \mu_w)} \underline{w} \in \underline{\mathcal{W}}_{n,k} \xleftarrow{h(\underline{w}, \nu_w)} w$$

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The word  $w$  is encoded into a triple by

$$\begin{aligned}\Gamma_k(w) &= (w_{-k}, \mu_w, \nu_w) \\ \Gamma^{-1}((w_{-k}, \mu_w, \nu_w)) &= h(g(w_{-k}, \mu_w), \nu_w).\end{aligned}$$

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For example,

$$\begin{aligned}\Gamma_3(123322311321123) &= (1222112112, \text{FFTFFFTFF}, \text{TFTFFFFT}) \\ 123322311321123 &= h(g(1222112112, \text{FFTFFFTFF}), \text{TFTFFFFT}).\end{aligned}$$

# Overview of the general bijection

We then define the bijection as

$$f_k(w) = \Gamma^{-1}(f_{k-1}(w_{-k}), \nu_w, \mu_k).$$

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In other words,

$$f_{k-1}(w_{-k}) \xleftarrow{g(f_{k-1}(w_{-k}), \nu_w)} \underline{f_k(w)} \xleftarrow{h(\underline{f_k(w)}, \mu_w)} f_k(w)$$

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For example,

$$\begin{aligned} f_3(123322311321123) &= \Gamma^{-1}(1222112112, \text{TFTFFFT}, \text{FFTFFFTFF}) \\ &= h(g(f_2(1222112112), \text{TFTFFFT}), \text{FFTFFFTFF}) \\ &= h(g(2112112221, \text{TFTFFFT}), \text{FFTFFFTFF}) \\ &= h(3211321122321, \text{FFTFFFTFF}) \\ &= 323113231122321. \end{aligned}$$



## The subword $w_{-k}$ of $w$

Consider the word

$$w = 1\ 1\ 3\ 3\ 2\ 3\ 3\ 1\ 1\ 2\ 2\ 2\ 3 \in \mathcal{W}_{4,3}.$$

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$$w_{-3} = 1\ 1\ 2\ 1\ 1\ 2\ 2\ 2 \in \mathcal{W}_{4,2}.$$

With this, we refine the descents of  $w \in \mathcal{W}_{n,k}$  having  $\text{des}(w) = d$  as

$$\begin{aligned}\text{des}(w) &= d_1 + d_2, \\ d_1 &= \text{des}(w_{-k}) \\ d_2 &= d - d_1\end{aligned}$$

and say that  $w \in \mathcal{W}_{n,k}^{d_1,d_2}$ .

## The unique subword $\underline{w}$ of $w$

We construct a minimal subword  $\underline{w}$  containing  $w_{-k}$  such that  $\text{des}(\underline{w}) = \text{des}(w)$ . Consider the word

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Now, remove the 3's which do not change the number of descents in  $w$ .

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Now, remove the 3's which do not change the number of descents in  $w$ . We can remove all 3's sandwiched within another descent.

$$1\ 1\ 3\ 3\ \boxed{2}\ \boxed{3\ 3}\ \boxed{1}\ 1\ 2\ 2\ 2\ 3.$$

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We can also remove the 3's that occurs at the end of the word.

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Finally, we can remove all but one 3 from a group of more than one 3's that is sandwiched between a non-descent.

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$$1\ 1\ \boxed{3}\ \boxed{3}\ 2\ 3\ 1\ 1\ 2\ 2\ 2.$$

When all such 3's have been removed, we get

$$\underline{w} = 1\ 1\ 3\ 2\ 1\ 1\ 2\ 2\ 2.$$

## Encoding $\underline{w}$ with respect to $w_k$ with $\mu_w$

Continuing the previous example, we annotate the non-descents in  $w_{-k}$ .

$$\begin{array}{rcccccccccccccccc} w_{-k} & = & \bullet & 1 & \bullet & 1 & \bullet & 2 & 1 & \bullet & 1 & \bullet & 2 & \bullet & 2 & \bullet & 2 \\ \underline{w} & = & \bullet & 1 & \bullet & 1 & 3 & 2 & 1 & \bullet & 1 & \bullet & 2 & \bullet & 2 & \bullet & 2. \end{array}$$

Represent the non-descents with the set  $\text{NDes}(w_{-k}) = \{0, 1, 2, 5, 6, 7, 8\}$ .

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$$\mu_w \in B_{(k-1)n-d_1}^{d_2},$$

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and the map  $g$  has the form

$$g : \mathcal{W}_{n,k-1}^{d_1} \times B_{(k-1)n-d_1}^{d_2} \longrightarrow \underline{\mathcal{W}}_{n,k}^{d_1, d_2}.$$

## Encoding $w$ with respect to $\underline{w}$ with $\nu_w$

Let  $v \in \mathcal{W}_{3,5}^{6,1}$  be the following word where we annotate the descents in  $v$  and the position at the end.

$$v = 3 \bullet 2 \ 4 \ 4 \bullet 1 \ 5 \bullet 3 \bullet 1 \ 2 \bullet 1 \ 4 \bullet 3 \bullet 2 \bullet .$$

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We can construct a balanced word  $w \in \mathcal{W}_{3,5}^{6,1}$  by inserting two 5's into any of the markers.



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We can construct a balanced word  $w \in \mathcal{W}_{3,5}^{6,1}$  by inserting two 5's into any of the markers. Two possibilities are

$$w = 3 \bullet 2 \ 4 \ 4 \ 5 \ 5 \ 1 \ 5 \bullet 3 \bullet 1 \ 2 \bullet 1 \ 4 \bullet 3 \bullet 2 \bullet$$

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We can construct a balanced word  $w \in \mathcal{W}_{3,5}^{6,1}$  by inserting two 5's into any of the markers. Two possibilities are

$$w = 3 \bullet 2 \ 4 \ 4 \ \textcolor{blue}{5} \ \textcolor{blue}{5} \ 1 \ 5 \bullet 3 \bullet 1 \ 2 \bullet 1 \ 4 \bullet 3 \bullet 2 \bullet$$

$$w = 3 \bullet 2 \ 4 \ 4 \bullet 1 \ 5 \bullet 3 \bullet 1 \ 2 \bullet 1 \ 4 \bullet 3 \ \textcolor{blue}{5} \ 2 \textcolor{blue}{5}.$$

As a boolean array, we can represent these arrangements as

$$\nu = F \ \textcolor{blue}{T} \textcolor{blue}{T} \ F \ F \ F \ F \ F \ F \quad \text{and} \quad \nu = F \ F \ F \ F \ F \ F \ \textcolor{blue}{T} \ F \ \textcolor{blue}{T}.$$

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Let  $v \in \mathcal{W}_{3,5}^{6,1}$  be the following word where we annotate the descents in  $v$  and the position at the end.

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We can construct a balanced word  $w \in \mathcal{W}_{3,5}^{6,1}$  by inserting two 5's into any of the markers. Two possibilities are

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In general, given  $w \in \mathcal{W}_{n,k}^{d_1,d_2}$ ,

$$\nu_w \in B_{n+d_1}^{n-d_2},$$

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As a boolean array, we can represent these arrangements as

$$\nu = F \ \textcolor{blue}{T} \ \textcolor{blue}{T} \ F \ F \ F \ F \ F \ F \quad \text{and} \quad \nu = F \ F \ F \ F \ F \ F \ \textcolor{blue}{T} \ \textcolor{blue}{F} \ \textcolor{blue}{T}.$$

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$$\nu_w \in B_{n+d_1}^{n-d_2},$$

and the map  $h$  has the form

$$h : \mathcal{W}_{n,k}^{d_1,d_2} \times B_{n+d_1}^{n-d_2} \longrightarrow \mathcal{W}_{n,k}^{d_1,d_2}.$$

# The map $\Gamma_k$ and $f_k$

## Lemma

*The function  $\Gamma_k$  given by  $\Gamma_k(w) = (w_{-k}, \mu_w, \nu_w)$  is a bijection.*

Formally, define the map  $f_k : \mathcal{W}_{n,k} \rightarrow \mathcal{W}_{n,k}$  for  $k \geq 3$  by

$$f_k(w) = \Gamma_k^{-1}((f_{k-1}(w_{-k}), \nu_w, \mu_w)).$$

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$$f_k(w) = \Gamma_k^{-1}((f_{k-1}(w_{-k}), \nu_w, \mu_w)).$$

## Theorem

*For  $k \geq 3$ ,  $f_k$  satisfies*

$$\text{des}(w) + \text{des}(f_k(w)) = (k-1)n$$

*and*

$$\text{maj}(w) + \text{maj}(f_k(w)) = n^2 \binom{k}{2} \text{ for all } w \in \mathcal{W}_{n,k}.$$

## The map $f_k$

To prove that the major index is reversed by  $f_k$ , we make use of the following lemma.

### Lemma

Let  $w \in \mathcal{W}_{n,k}$  for  $k \geq 3$ . Then,

$$\text{maj}(w) - \text{maj}(w_{-k}) + \text{maj}(f_k(w)) - \text{maj}(f_k(w)_{-k}) = (k-1)n^2.$$

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## Corollary

The bijections satisfy the property  $f_{k-1}(w_{-k}) = f_k(w)_{-k}$  for all  $w \in \mathcal{W}_{n,k}$  for all  $n \geq 1, k \geq 3$ .



# Applying $f_k$ on permutations

Unlike the cases where  $n > 1$ , the descent sets of permutations behave in the following way under  $f_k$ .

## Theorem

Let  $w \in S_k$ . Then  $\text{Des}(f_k(w)) = [k-1] \setminus \text{Des}(w)$ .

## Example

$$\begin{array}{rcl} w & = & 4 \ 6 \ 8 \ 7 \ 2 \ 9 \ 5 \ 1 \ 3 \\ f_8(w) & = & 7 \ 6 \ 1 \ 5 \ 8 \ 3 \ 4 \ 9 \ 2 \end{array}$$

## An example of $f_k$ on $\mathcal{W}_{3,5}$

$w = 325544135121432$

$\underline{w} = 3254413121432$

$\mu_w = \text{FTFFFF}$

$\nu_w = \text{FTFFTFFFF}$

$f_5(w) = 154524225331134$

$w = 321312132$

$\underline{w} = 321312132$

$\mu_w = \text{TTF T}$

$\nu_w = \text{FFFFF}$

$f_3(w) = 122233113$

$w = 324413121432$

$\underline{w} = 3213121432$

$\mu_w = \text{FFFT}$

$\nu_w = \text{FTTFFFFF}$

$f_4(w) = 142422331134$

$w = 211212$

$\mu_w = \text{TFT}$

$\nu_w = \text{TTF}$

$f_2(w) = 122211$

Thank you!