

Math 582G

Today: Hyperbolic Programming

A hyperbolic program is an problem of the form

$$\min \langle c, x \rangle \quad \text{s.t.} \quad x \in \overline{C(f, e)} \quad \text{and} \quad \langle a_i, x \rangle = b_i \quad i=1, \dots, m$$

where $f \in \mathbb{R}[x_1, \dots, x_n]_d$ is hyperbolic w.r.t. $e \in \mathbb{R}^n$,

$$c, a_1, \dots, a_m \in \mathbb{R}^n, \quad b_1, \dots, b_m \in \mathbb{R}.$$



Ex: $\overline{C(f, e)} = \mathbb{R}_{\geq 0}^n \rightarrow$ linear programming (LP)

Ex: $\overline{C(f, e)} = \text{PSD}_d \rightarrow$ semidefinite programming (SDP)

Güler ('97) and Renegar (early '00's) showed that interior pt methods for LPs and SDPs extend to general hyperbolic programs.

Idea: Use $-\log(f)$ as a "barrier" function for $C(f, e)$

Note: For $a \in C(f, e)$, $f(a) > 0 \Rightarrow -\log(f(a))$ defined

$$-\log(f(a)) \rightarrow \infty \quad \text{as} \quad a \rightarrow \partial \overline{C(f, e)} \quad (\text{and} \quad f(a) \rightarrow 0)$$

Thm: If f is hyperbolic w.r.t. e and $\overline{C(f, e)} \cap (-\overline{C(f, e)}) = \{0\}$ then the Hessian of $-\log(f)$ is positive definite at every point $a \in C(f, e)$.

(Proof) Let $a \in C(f, e)$ and $v \in \mathbb{R}^n$. Then

$$v^T (\nabla^2 \log(f(a))) v = \left[\frac{d^2}{dt^2} \log(f(a+tv)) \right]_{t=0}.$$

WTS this is ≤ 0 and < 0 for $v \neq 0$.

Since f is hyp. w.r.t. a , $f(a+tv) = f(a) \prod_{j=1}^d (1 + \lambda_j t)$

where $\lambda_k = k^{\text{th}}$ eig. val of v w.r.t. a .

$$\Rightarrow \frac{d^2}{dt^2} \log(f(a+tv)) = \frac{d^2}{dt^2} \left[\log(f(a)) + \sum_{j=1}^d \log(1 + \lambda_j t) \right]$$

$$= \sum_{j=1}^d \frac{d^2}{dt^2} \log(1 + \lambda_j t) = \sum_{j=1}^d \frac{d}{dt} \left(\frac{\lambda_j}{(1 + \lambda_j t)} \right) = \sum_{j=1}^d \frac{-\lambda_j^2}{(1 + \lambda_j t)^2}$$

$$\Rightarrow v^T (\nabla^2 \log(f(a))) v = \left[\frac{d^2}{dt^2} \log(f(a+tv)) \right]_{t=0} = \sum_{j=1}^d -\lambda_j^2$$

Since $\lambda_1, \dots, \lambda_d \in \mathbb{R}$, this is ≤ 0

and equals zero $\Leftrightarrow \lambda_1 = \dots = \lambda_d = 0 \Leftrightarrow v \in \overline{C(f, e)} \cap -\overline{C(f, e)}$.

Cor: For any $\lambda \in \mathbb{R}_{>0}$, the Hessian of

$\langle c, x \rangle - \lambda \log(f)$ is positive definite at $a \in C(f, e)$.

If $S = \{x \in \overline{C(f, e)} : \langle a_i, x \rangle = b_i, i=1, \dots, m\}$ is compact

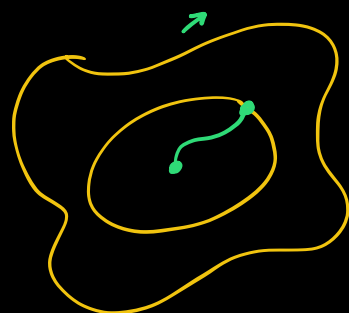
then for each $\lambda > 0$, $\langle c, x \rangle - \lambda \log(f(x))$ is

minimized by a unique point $x^*(\lambda) \in S$.

$x^*(\lambda)$ is continuous and $x^*(\lambda) \rightarrow$ minimizer of $\langle c, x \rangle$ as $\lambda \rightarrow 0$

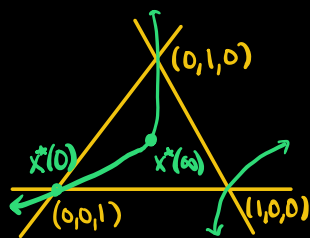
Def: The central path of a hyperbolic program is

$$\{x^*(\lambda) : \lambda \in \mathbb{R}_{>0}\}.$$



Idea of interior pt. methods: find piecewise linear approximation of central path as $\lambda \rightarrow 0$

Ex: $\min x_1 + 2x_2$ s.t. $x \in \mathbb{R}_{\geq 0}^3$, $x_1 + x_2 + x_3 = 1$



$$F_\lambda(x) = x_1 + 2x_2 - \lambda \log(x_1 x_2 x_3)$$

$$\nabla F_\lambda(x) = (1, 2, 0) - \lambda \left(\frac{1}{x_1}, \frac{1}{x_2}, \frac{1}{x_3} \right)$$

$$\nabla F_\lambda(x) = \mu(1, 1, 1) \Rightarrow 0 = \det \begin{pmatrix} \frac{1}{x_1} & \frac{1}{x_2} & \frac{1}{x_3} \\ \frac{1}{x_1} & \frac{1}{x_2} & \frac{1}{x_3} \\ \frac{1}{x_1} & \frac{1}{x_2} & \frac{1}{x_3} \end{pmatrix} = -\frac{2}{x_1} + \frac{1}{x_2} + \frac{1}{x_3}$$

$$\Rightarrow -2x_2 x_3 + x_1 x_3 + x_1 x_2 = 0.$$

Thm (Renegar) All faces of $\overline{C(f, e)}$ are exposed.

Idea: For $x \in \mathbb{R}^n$, define $\text{rank}(x) = \#$ nonzero eig. val. of x .

If f is homog. of degree d , this is $d -$ (mult of 0 as a root of $f(te-x)$).

• If $\text{rank}(a) = d-1$ and $a \in \overline{C(f, e)}$

$\mathbb{R}_{>0} \{a\}$ is exposed by $x \mapsto \langle x, \nabla f(a) \rangle$



- If F is a face of $\overline{C(f,e)}$, then $\text{rank}(a)$ is constant on the relative interior of F
- If $\text{rank}(a) = r$ for $a \in \text{relint}(F)$ then F is a face of $C(D_e^{d-r-1} f, e)$ with mult. 1 \Rightarrow exposed by $x \mapsto \langle x, \nabla D_e^{d-r-1} f(a) \rangle$.



Open Question: Is hyperbolic programming stronger than semidefinite programming?

Related Open Question: Is every hyperbolicity cone $C(f, e)$ linearly isomorphic to $L \cap \text{PSD}_N$ for some N and linear space L ?