

# Math 409: Discrete Optimization

Today: LP for matchings, total unimodularity

## LP for max matchings in $G=(V,E)$

$e \in E \rightarrow$  variable  $x_e$ , constrained  $0 \leq x_e \leq 1$

$$\begin{cases} x_e = 1 \leftrightarrow e \in M \\ x_e = 0 \leftrightarrow e \notin M \end{cases}$$

Constraints:  $\deg_v(M) \leq 1 \rightarrow \sum_{e \ni v} x_e \leq 1 \quad \forall v \in V$

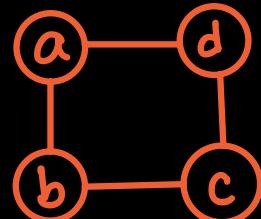
Maximize  $|M|$  (or  $\sum_{e \in M} c_e$  for some  $c: E \rightarrow \mathbb{R}$ )  $\rightarrow \max \sum_{e \in E} c_e x_e$

(IP)  $\max \sum_{e \in E} c_e x_e$  s.t.  $0 \leq x_e \leq 1$  for all  $e \in E$ ,  $x \in \mathbb{Z}^E$   
 $\sum_{e \ni v} x_e \leq 1$  for all  $v \in V$

Claim: This solves max weight matching on  $G$

(Recover  $M$  from opt sol  $x^*$  by  $M = \{e : x_e^* = 1\}$ ) You check!

Ex:



Feas. region of LP relaxation:

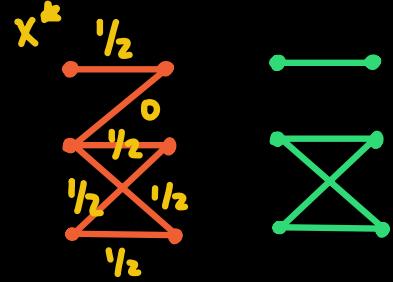
$$\left\{ x \in \mathbb{R}^4 : x_{ab}, x_{bc}, x_{cd}, x_{ad} \in [0, 1] \text{ and } x_{ab} + x_{ad} \leq 1, x_{ab} + x_{bc} \leq 1, x_{bc} + x_{cd} \leq 1, x_{ad} + x_{cd} \leq 1 \right\}$$

Vertices:  $(x_{ab}, x_{bc}, x_{cd}, x_{ad}) = (0, 0, 0, 0), (1, 0, 0, 0), (0, 1, 0, 0), (0, 0, 1, 0), (0, 0, 0, 1)$

All integer!

Thm: For a bipartite graphs, any vertex  $x^*$  of the matching LP is integer. ( $\Rightarrow P = P_I$  for  $P =$  feas. region of matching LP)

(Proof) Suppose  $x^*$  is a vertex. Define  $G_{x^*} = (V, E_{x^*})$  with  $E_{x^*} = \{e \in E \text{ s.t. } 0 < x_e^* < 1\}$ .



(Case 1) Suppose  $G_{x^*}$  has a cycle  $C$

$G_{x^*}$  bipartite  $\Rightarrow C$  has even length:  $C = (e_1, e_2, \dots, e_{2k})$

Define  $w \in \mathbb{R}^E$  by  $w_e = \begin{cases} 1 & \text{if } e = e_{2i} \in C \\ -1 & \text{if } e = e_{2i+1} \in C \\ 0 & \text{o.w.} \end{cases}$



Take  $\varepsilon = \min \{x_e^* : e \in C\} \cup \{1 - x_e^* : e \in C\}$

Then  $x^* + \varepsilon w$  and  $x^* - \varepsilon w$  both feasible  $\Rightarrow x^*$  not a vertex

check  $(x^* \pm \varepsilon w)_e \in [0, 1] \quad \forall e$  and  $\sum_{e \in V} (x^* \pm \varepsilon w)_e = \sum_{e \in V} x_e^*$

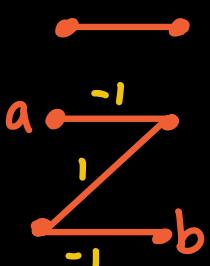
(Case 2)  $G_{x^*}$  is acyclic,  $E_f$  nonempty

$\Rightarrow$  any connected comp. of  $G_{x^*}$  is a tree

$\Rightarrow G_{x^*}$  has two vertices  $a, b$  of  $\deg = 1$ .

Let  $P =$  unique  $a$ - $b$  path in  $G_{x^*}$ ,  $P = (e_1, \dots, e_k)$

Define  $w \in \mathbb{R}^E$  by  $w_e = \begin{cases} 1 & \text{if } e = e_{2i} \in P \\ -1 & \text{if } e = e_{2i+1} \in P \\ 0 & \text{if } e \notin P \end{cases}$



For  $\varepsilon = \min \{x_e : e \in P\} \cup \{1 - x_e : e \in P\}$ ,  $x^* \pm \varepsilon w$  feasible

Check: 1)  $(x^* \pm \varepsilon w)_e$  in  $[0, 1]$   $\Rightarrow x^*$  not a vertex

2)  $\sum_{e \in V} (x^* \pm \varepsilon w)_e = \sum_{e \in V} (x^* \pm \varepsilon w) \quad \forall v \in V \setminus \{a, b\}$

3) For  $v=a,b$ ,  $\exists$  unique edge  $\tilde{e} \ni v$  with  $x_{\tilde{e}}^* \in (0,1)$

$$\Rightarrow \sum_{e \ni v} x_e^* = x_{\tilde{e}}^* + \sum_{\substack{e \ni v \\ e \neq \tilde{e}}} x_e^* \in [0,1] \Rightarrow \sum_{e \ni v} x_e^* = x_{\tilde{e}}^* \in (0,1)$$

must be integer < 1  $\Rightarrow = 0$

$$\Rightarrow \sum_{e \ni v} (x_e^* \pm \varepsilon_w) = (x_{\tilde{e}}^* \pm \varepsilon_w) \in [0,1]$$

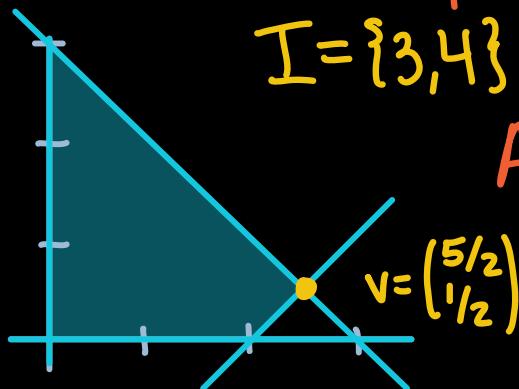
(Case 3)  $E_x = \emptyset \Rightarrow x^* \in \mathbb{Z}^E$

What is special about bipartite graphs?

Recall: A point  $v$  in  $P = \{x \in \mathbb{R}^n : Ax \leq b\}$  is a vertex of  $P \Leftrightarrow \exists I \subseteq \{1, \dots, m\}$  with  $|I|=n$  s.t. the  $n \times n$  submatrix  $A_I$  with rows  $\{a_i^T : i \in I\}$  has full rank  $n$  and  $a_i^T v = b_i$  for all  $i \in I$ .

$$\Rightarrow v = A_I^{-1} b_I \quad (\text{i.e. } v \text{ is the unique sol. to } a_i^T x = b_i \text{ for all } i \in I)$$

Ex: (P)  $\max 2x_1 + x_2$  s.t.  $-x_1 \leq 0, -x_2 \leq 0, x_1 + x_2 \leq 3, x_1 - x_2 \leq 2$



$$A = \begin{pmatrix} -1 & 0 \\ 0 & -1 \\ 1 & 1 \\ 1 & -1 \end{pmatrix}$$

$$A_{34} = \begin{pmatrix} 1 & 1 \\ 1 & -1 \end{pmatrix} \quad v = A_{34}^{-1} b_{34}$$

$$= \frac{1}{2} \begin{pmatrix} 1 & 1 \\ 1 & -1 \end{pmatrix} \begin{pmatrix} 3 \\ 2 \end{pmatrix}$$

$$= \frac{1}{2} \begin{pmatrix} 5 \\ 1 \end{pmatrix} = \begin{pmatrix} 5/2 \\ 1/2 \end{pmatrix}$$

For  $2 \times 2$  matrices  $A = \begin{pmatrix} a & b \\ c & d \end{pmatrix}$ ,  $A^{-1} = \frac{1}{ad-bc} \begin{pmatrix} d & -b \\ -c & a \end{pmatrix}$   
 If  $a, b, c, d \in \mathbb{Z}$  and  $\det(A) = \pm 1$  then entries of  $A^{-1}$  are also integer!

More generally...

Cramers Rule: Let  $B \in \mathbb{R}^{n \times n}$  with  $\text{rank}(B) = n$ .  
Then  $B^{-1}$  has entries

$$(B^{-1})_{ij} = (-1)^{i+j} \left( \frac{1}{\det(B)} \right) \det(M_{ji})$$

where  $M_{ji}$  is the  $(n-1) \times (n-1)$  matrix obtained by removing the  $j^{\text{th}}$  row and  $i^{\text{th}}$  col of  $B$ .

Cor: If  $B \in \mathbb{Z}^{n \times n}$  and  $\det(B) = \pm 1$ , then  $B^{-1} \in \mathbb{Z}^{n \times n}$ .  
 $\Rightarrow$  for all  $b \in \mathbb{Z}^n$ ,  $B^{-1}b \in \mathbb{Z}^n$

Def: A matrix  $A \in \mathbb{R}^{m \times n}$  is totally unimodular (TU) if every square submatrix has determinant 0,  $\pm 1$ .  
( $1 \times 1$  submatrices  $\Rightarrow$  all entries of  $A$  are 0,  $\pm 1$ )

Cor: If  $A$  is TU, for any  $b \in \mathbb{Z}^m$ , all vertices of  $P = \{x \in \mathbb{R}^n : Ax \leq b\}$  are integer.