

# Large Example of the Dual Simplex Method

UW Math 407, Fall 2022

Below is a large example of the dual simplex method, carried through until an optimal solution is found. Afterwards, is a side-by-side comparison of using the usual simplex method on the dual LP.

## LARGE EXAMPLE

Original LP:  $\max -x_1 - x_2 - x_3$  such that  $A\mathbf{x} = \mathbf{b}$  and  $\mathbf{x} \geq 0$  where

$$A = \begin{pmatrix} 1 & -1 & 0 & 1 & 0 & 0 & 0 \\ 1 & -1 & -1 & 0 & 1 & 0 & 0 \\ -1 & 0 & 1 & 0 & 0 & 1 & 0 \\ 0 & -1 & 0 & 0 & 0 & 0 & 1 \end{pmatrix} \quad \text{and} \quad \mathbf{b} = \begin{pmatrix} -1 \\ -2 \\ -1 \\ -3 \end{pmatrix}$$

From inspection, we see that the basis  $B = \{4, 5, 6, 7\}$  is dual feasible, so we can start the dual simplex method from there:

$B = \{4, 5, 6, 7\}$ ,  $\mathcal{T}(B)$  :

$$x_4 = -1 - x_1 + x_2$$

$$x_5 = -2 - x_1 + x_2 + x_3$$

$$x_6 = -1 + x_1 - x_3$$

$$x_7 = -3 + x_2$$

$$z = -x_1 - x_2 - x_3$$

This basis is not feasible. We see that any of  $x_4, x_5, x_6, x_7$  can leave the basis. We choose  $x_4$  to leave. We then add the highest nonnegative scalar multiple  $\lambda$  of the “ $x_4 =$ ” row to the last row to eliminate some variable. We compute

$$\lambda x_4 + z = -\lambda + (-\lambda - 1)x_1 + (\lambda - 1)x_2 - x_3.$$

For the coefficients of  $x_1, x_2, x_3$  to be  $\leq 0$ , the tightest constraint comes from the coefficient of  $x_2$ , so  $x_2$  enters the basis, giving

$B = \{2, 5, 6, 7\}$ ,  $\mathcal{T}(B)$  :

$$x_2 = 1 + x_1 + x_4$$

$$x_5 = -1 + x_3 + x_4$$

$$x_6 = -1 + x_1 - x_3$$

$$x_7 = -2 + x_1 + x_4$$

$$z = -1 - 2x_1 - x_3 - x_4$$

This basis is also not feasible. Any of  $x_5, x_6, x_7$  can leave. Let's choose  $x_5$  to leave. We add the highest nonnegative scalar multiple  $\lambda$  of the “ $x_5 =$ ” row to the last row to eliminate some variable. We compute

$$\lambda x_5 + z = (-\lambda - 1) - 2x_1 + (\lambda - 1)x_3 + (\lambda - 1)x_4.$$

For the coefficients of  $x_1, x_3, x_4$  to be  $\leq 0$ , the tightest constraint of  $\lambda \leq 1$  comes from the coefficient of  $x_3$  and  $x_4$ . We can choose either to enter. Let's choose  $x_3$  to enter, giving

$B = \{2, 3, 6, 7\}$ ,  $\mathcal{T}(B)$  :

$$\begin{aligned}x_2 &= 1 + x_1 + x_4 \\x_3 &= 1 - x_4 + x_5 \\x_6 &= -2 + x_1 + x_4 - x_5 \\x_7 &= -2 + x_1 + x_4 \\z &= -2 - 2x_1 - x_5\end{aligned}$$

This basis is also not feasible. Either of  $x_6$  or  $x_7$  can leave. Let's choose  $x_6$  to leave. We add the highest nonnegative scalar multiple  $\lambda$  of the " $x_6 =$ " row to the last row to eliminate some variable. We compute

$$\lambda x_6 + z = (-2\lambda - 2) + (\lambda - 2)x_1 + \lambda x_4 + (-\lambda - 1)x_5$$

For the coefficients of  $x_1, x_4, x_5$  to be  $\leq 0$ , the tightest constraint of  $\lambda \leq 0$  comes from the coefficient of  $x_4$ , so  $x_4$  to enters, giving

$B = \{2, 3, 4, 7\}$ ,  $\mathcal{T}(B)$  :

$$\begin{aligned}x_2 &= 3 + x_5 + x_6 \\x_3 &= -1 + x_1 - x_6 \\x_4 &= 2 - x_1 + x_5 + x_6 \\x_7 &= x_5 + x_6 \\z &= -2 - 2x_1 - x_5\end{aligned}$$

This basis is still not feasible. Only  $x_3$  can leave. We add the highest nonnegative scalar multiple  $\lambda$  of the " $x_3 =$ " row to the last row to eliminate some variable. We compute

$$\lambda x_3 + z = (-\lambda - 2) + (\lambda - 2)x_1 - x_5 - \lambda x_6$$

For the coefficients of  $x_1, x_5, x_6$  to be  $\leq 0$ , the tightest constraint of  $\lambda \leq 2$  comes from the coefficient of  $x_1$ , so  $x_1$  to enters, giving

$B = \{1, 2, 4, 7\}$ ,  $\mathcal{T}(B)$  :

$$\begin{aligned}x_1 &= 1 + x_3 + x_6 \\x_2 &= 3 + x_5 + x_6 \\x_4 &= 1 - x_3 + x_5 \\x_7 &= x_5 + x_6 \\z &= -4 - 2x_3 - x_5 - 2x_6\end{aligned}$$

This basis  $B = \{1, 2, 4, 7\}$  is now both feasible and dual feasible, so it corresponds to an optimal solution, namely  $\mathbf{x}^* = (1, 3, 0, 1, 0, 0, 0)$ .

SIDE-BY-SIDE VIEW OF PRIMAL AND DUAL

The dual linear program is

$$\min -y_1 - 2y_2 - y_3 - 3y_4 \quad \text{s.t.} \quad \begin{pmatrix} 1 & 1 & -1 & 0 \\ -1 & -1 & 0 & -1 \\ 0 & -1 & 1 & 0 \\ 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \end{pmatrix} \begin{pmatrix} y_1 \\ y_2 \\ y_3 \\ y_4 \end{pmatrix} - \begin{pmatrix} s_1 \\ s_2 \\ s_3 \\ s_4 \\ s_5 \\ s_6 \\ s_7 \end{pmatrix} = \begin{pmatrix} -1 \\ -1 \\ -1 \\ 0 \\ 0 \\ 0 \\ 0 \end{pmatrix}, \quad \begin{pmatrix} s_1 \\ s_2 \\ s_3 \\ s_4 \\ s_5 \\ s_6 \\ s_7 \end{pmatrix} \geq 0.$$

As in HW6, Problem 2, we can eliminate the variables  $\mathbf{y}$  and rewrite this as

$$\min -s_4 - 2s_5 - s_6 - 3s_7 \quad \text{s.t.} \quad \begin{pmatrix} 1 & 0 & 0 & -1 & -1 & 1 & 0 \\ 0 & 1 & 0 & 1 & 1 & 0 & 1 \\ 0 & 0 & 1 & 0 & 1 & -1 & 0 \end{pmatrix} \mathbf{s} = \begin{pmatrix} 1 \\ 1 \\ 1 \end{pmatrix} \quad \text{and } \mathbf{s} \geq 0.$$

So that we can translate to maximization, let's rewrite this as

$$- \max s_4 + 2s_5 + s_6 + 3s_7 \quad \text{s.t.} \quad \begin{pmatrix} 1 & 0 & 0 & -1 & -1 & 1 & 0 \\ 0 & 1 & 0 & 1 & 1 & 0 & 1 \\ 0 & 0 & 1 & 0 & 1 & -1 & 0 \end{pmatrix} \mathbf{s} = \begin{pmatrix} 1 \\ 1 \\ 1 \end{pmatrix} \quad \text{and } \mathbf{s} \geq 0.$$

The basis  $\{1, 2, 3\}$  is clearly feasible. Let us look at a side-by-side progression of the dual simplex method on the original problem and the usual simplex method on this linear program in s:

$\mathcal{T}(\{4, 5, 6, 7\}) :$ $x_4 = -1 - x_1 + x_2$ $x_5 = -2 - x_1 + x_2 + x_3$ $x_6 = -1 + x_1 - x_3$ $x_7 = -3 + x_2$ $z = -x_1 - x_2 - x_3$	$\mathcal{T}(\{1, 2, 3\}) :$ $s_1 = 1 + s_4 + s_5 - s_6$ $s_2 = 1 - s_4 - s_5 - s_7$ $s_3 = 1 - s_5 + s_6$ $z = s_4 + 2s_5 + s_6 + 3s_7$
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$x_4$  leaves,  $x_2$  enters

$s_4$  enters,  $s_2$  leaves

$\mathcal{T}(\{2, 5, 6, 7\}) :$ $x_2 = 1 + x_1 + x_4$ $x_5 = -1 + x_3 + x_4$ $x_6 = -1 + x_1 - x_3$ $x_7 = -2 + x_1 + x_4$ $z = -1 - 2x_1 - x_3 - x_4$	$\mathcal{T}(\{1, 3, 4\}) :$ $s_1 = 2 - s_2 - s_6 - s_7$ $s_3 = 1 - s_5 + s_6$ $s_4 = 1 - s_2 - s_5 - s_7$ $z = 1 - s_2 + s_5 + s_6 + 2s_7$
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$x_5$  leaves,  $x_3$  enters

$s_5$  enters,  $s_3$  leaves

$\mathcal{T}(\{2, 3, 6, 7\}) :$

$$\begin{aligned}x_2 &= 1 + x_1 + x_4 \\x_3 &= 1 - x_4 + x_5 \\x_6 &= -2 + x_1 + x_4 - x_5 \\x_7 &= -2 + x_1 + x_4 \\z &= -2 - 2x_1 - x_5\end{aligned}$$

$x_6$  leaves,  $x_4$  enters

$\mathcal{T}(\{2, 3, 4, 7\}) :$

$$\begin{aligned}x_2 &= 3 + x_5 + x_6 \\x_3 &= -1 + x_1 - x_6 \\x_4 &= 2 - x_1 + x_5 + x_6 \\x_7 &= x_5 + x_6 \\z &= -2 - 2x_1 - x_5\end{aligned}$$

$x_3$  leaves,  $x_1$  enters

$\mathcal{T}(\{1, 2, 4, 7\}) :$

$$\begin{aligned}x_1 &= 1 + x_3 + x_6 \\x_2 &= 3 + x_5 + x_6 \\x_4 &= 1 - x_3 + x_5 \\x_7 &= x_5 + x_6 \\z &= -4 - 2x_3 - x_5 - 2x_6\end{aligned}$$

Done!

$\mathcal{T}(\{1, 4, 5\}) :$

$$\begin{aligned}s_1 &= 2 - s_2 - s_6 - s_7 \\s_4 &= -s_2 + s_3 - s_6 - s_7 \\s_5 &= 1 - s_3 + s_6 \\z &= 2 - s_2 - s_3 + 2s_6 + 2s_7\end{aligned}$$

$s_6$  enters,  $s_4$  leaves

$\mathcal{T}(\{1, 5, 6\}) :$

$$\begin{aligned}s_1 &= 2 - s_3 + s_4 \\s_5 &= 1 - s_2 - s_4 - s_7 \\s_6 &= -s_2 + s_3 - s_4 - s_7 \\z &= 2 - 3s_2 + s_3 - 2s_4\end{aligned}$$

$s_3$  enters,  $s_1$  leaves

$\mathcal{T}(\{3, 5, 6\}) :$

$$\begin{aligned}s_3 &= 2 - s_1 + s_4 \\s_5 &= 1 - s_2 - s_4 - s_7 \\s_6 &= 2 - s_1 - s_2 - s_7 \\z &= 4 - s_1 - 3s_2 - s_4\end{aligned}$$

Done!