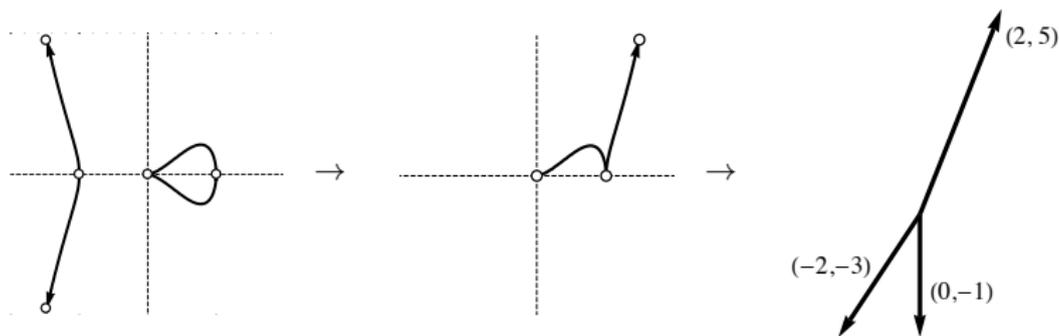


# Numerical methods for computing real and complex tropical curves

Cynthia Vinzant

North Carolina State University



joint with Daniel Brake and Jonathan Hauenstein

# Puiseux series, valuations, and tropical varieties

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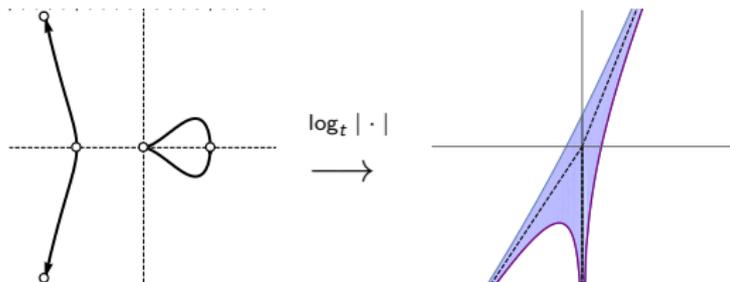
We can take the variety of  $I \subset \mathbb{k}[x_1, \dots, x_n]$  over  $\mathbb{k}\{\{t\}\}$ .

The  **$\mathbb{k}$ -tropical variety** of  $I$  is  $\text{Trop}_{\mathbb{k}}(I) = \overline{-\text{val}(\mathcal{V}_{\mathbb{k}\{\{t\}\}}(I))} \subset \mathbb{R}^n$ .

# Logarithmic limit sets

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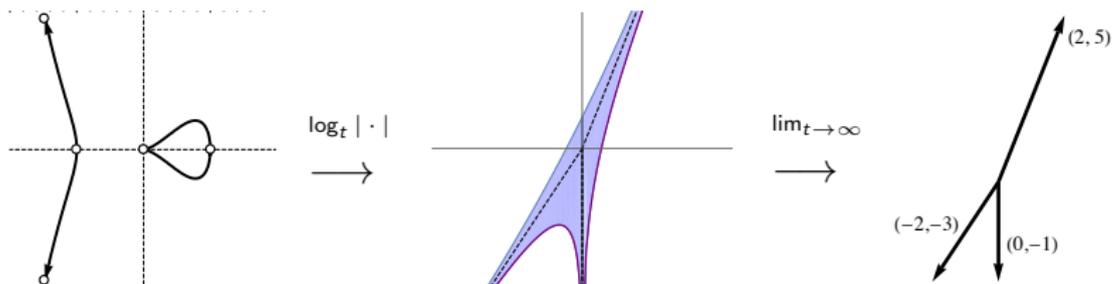
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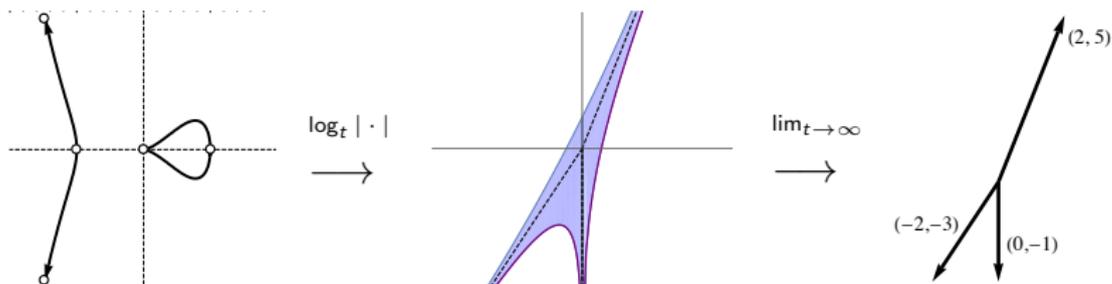
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For both  $\mathbb{k} = \mathbb{R}, \mathbb{C}$ ,  $\mathcal{A}_\infty(\mathcal{V}_{\mathbb{k}}(I))$  equals  $\text{Trop}_{\mathbb{k}}(I)$ .



# Connections with initial ideals

For  $w \in \mathbb{R}^n$ ,  $f = \sum_{\alpha} c_{\alpha} x^{\alpha}$ , define  $\text{in}_w(f) = \sum_{\alpha \in \mathcal{A}} c_{\alpha} x^{\alpha}$ , where  $\mathcal{A}$  is the set of  $\alpha$  maximizing  $w \cdot \alpha$ . Then  $\text{in}_w(I) = \langle \text{in}_w(f) : f \in I \rangle$ .

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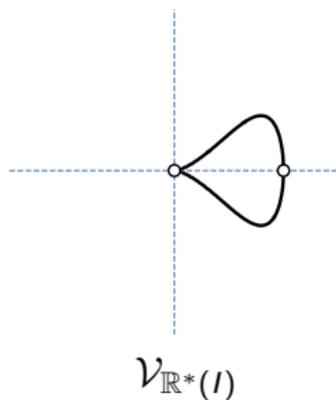
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E.g. For  $I = \langle (x - y)^2 + 1 \rangle$ ,  $\mathcal{V}_{\mathbb{R}}(I) = \emptyset$  and  $\text{Trop}_{\mathbb{R}}(I) = \emptyset$ ,  
but  $\mathcal{V}_{\mathbb{R}^*}(\text{in}_{(1,1)}(I)) = \mathcal{V}_{\mathbb{R}^*}((x - y)^2) \neq \emptyset$ .

# Example: sextic plane curve

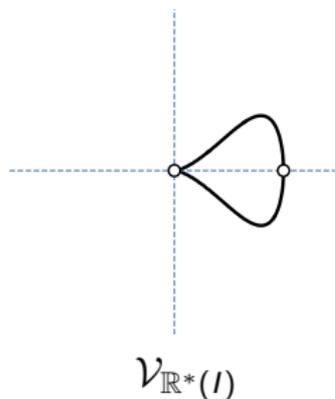
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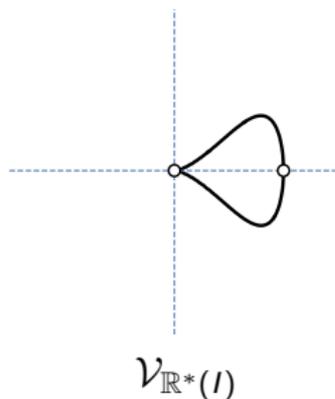


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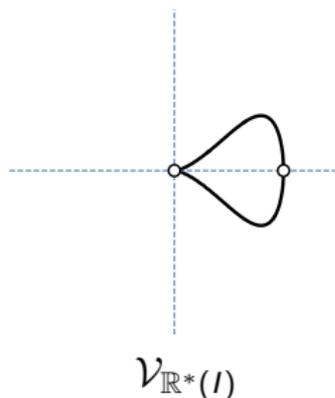
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$\mathcal{V}_{\mathbb{R}^*}(I)$

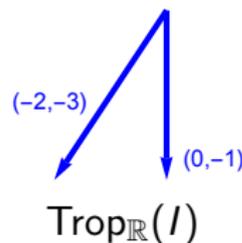
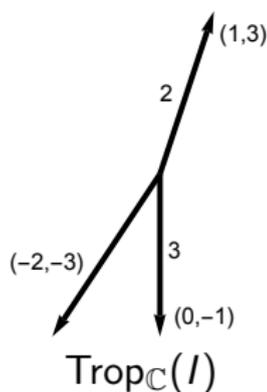
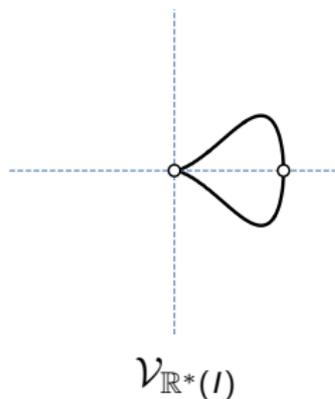
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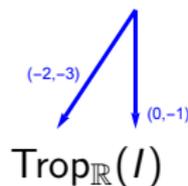
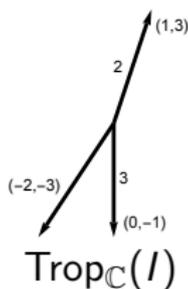
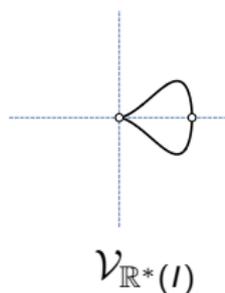
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# Tropical varieties are polyhedral fans

## Theorem (Fundamental Theorem of Tropical Geometry)

For an irreducible ideal  $I \subset \mathbb{C}[x_1, \dots, x_n]$ ,  $\text{Trop}_{\mathbb{C}}(I)$  is a rational pure-dimensional polyhedral fan of dimension  $d = \dim(\mathcal{V}_{\mathbb{C}^*}(I))$ . Counting cones with appropriate multiplicities, it is balanced.



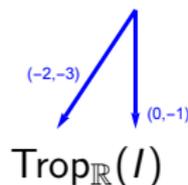
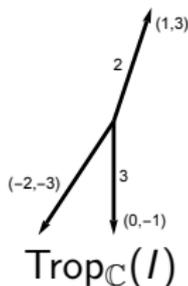
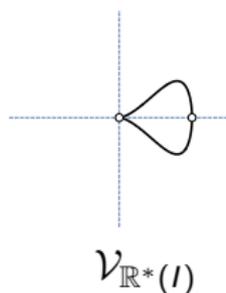
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## Theorem (Alessandrini, 2013)

For an ideal  $I \subset \mathbb{R}[x_1, \dots, x_n]$ ,  $\text{Trop}_{\mathbb{R}}(I)$  is a rational polyhedral fan of dimension  $d \leq \dim(\mathcal{V}_{\mathbb{R}^*}(I))$ .



# Computing Tropical Varieties

**Symbolic Algorithm:** Gfan – developed by Jensen  
uses Gröbner bases to compute  $\text{Trop}_{\mathbb{C}}(I)$  for any  $I \subset \mathbb{Q}[x_1, \dots, x_n]$ .

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## Numerical Algorithms:

- ▶  $\text{Trop}_{\mathbb{C}}$  of hypersurfaces (Hauenstein, Sottile, 2014)
- ▶  $\text{Trop}_{\mathbb{C}}$  of curves (Jensen, Leykin, Yu, 2015)
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Numerical algorithms do not require defining equations.

Curves are tractable and are used in internal computations for  $\text{Trop}_{\mathbb{C}}$  of larger dimensional varieties.

# Strategy for computing $\text{Trop}_{\mathbb{k}}(I)$

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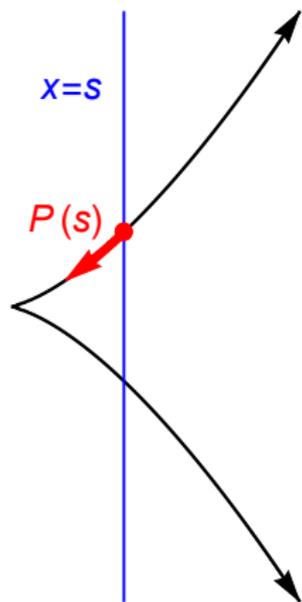
**Cauchy Integrals:** If  $f(z)$  is analytic on  $\{z \in \mathbb{C} : |z| \leq \tau\}$ , then

$$f^{(k)}(0) = \frac{k!}{2\pi} \int_0^{2\pi} \frac{f(\tau e^{i\theta})}{(\tau e^{i\theta})^{k+1}} d\theta,$$

and  $f(z) = \sum_{k=0}^{\infty} f^{(k)}(0) \cdot \frac{1}{k!} \cdot z^k$  for  $|z| \leq \tau$ .

# Finding an analytic parametrization: monodromy

Suppose  $\{P(s) : s \in [0, \tau]\} \subset C$  and  $P_j(s) = s$ .

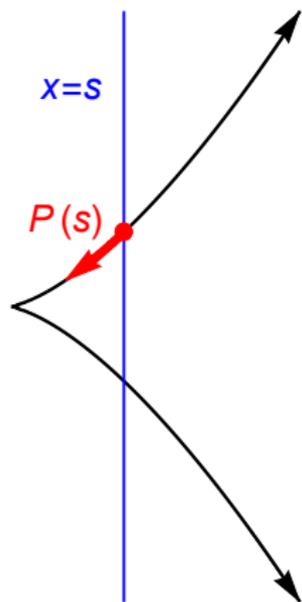


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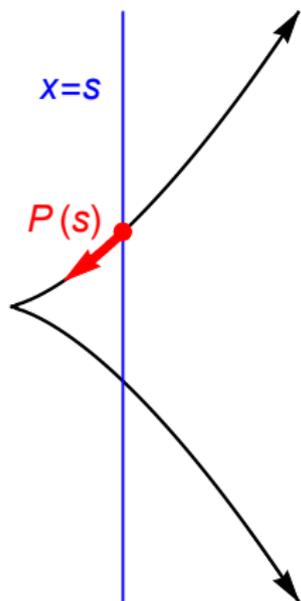
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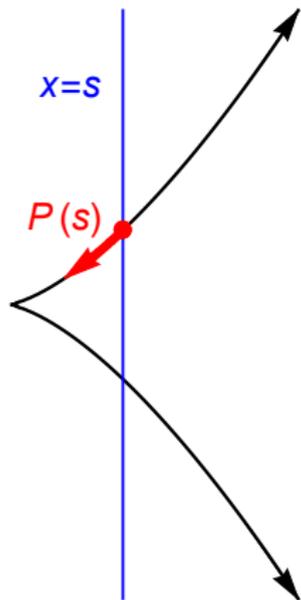
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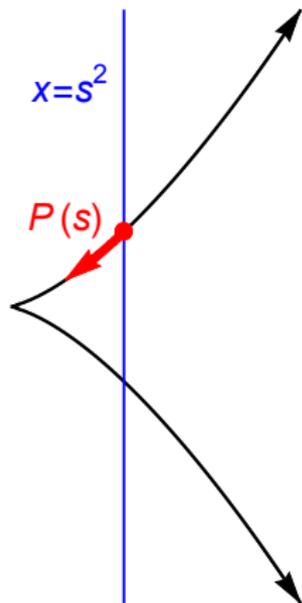
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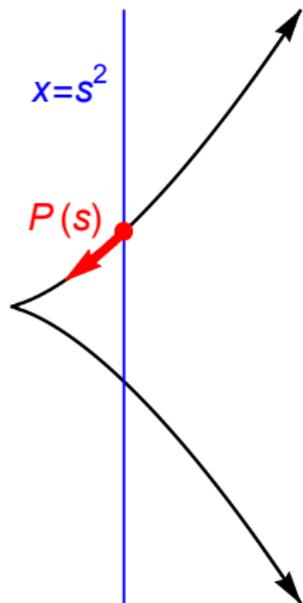
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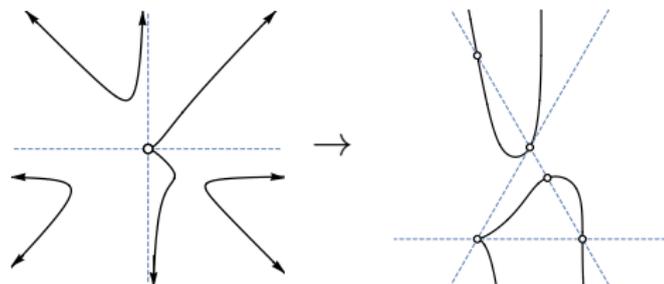
Re-parametrize:  $s \mapsto s^2$ ,  $P(s) = (s^2, s^3)$



# Sketch of Algorithm

Let  $C \subset \mathbb{C}^n$  be an irreducible curve.

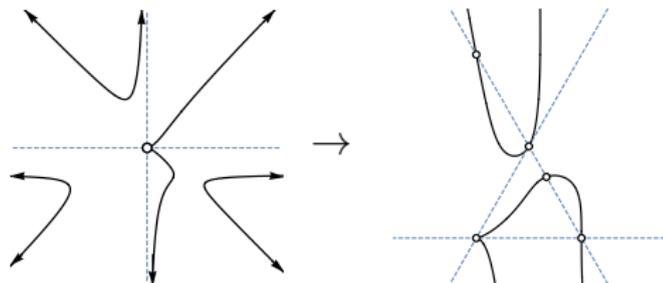
- ▶ Take  $\overline{C} \subset \mathbb{P}^n(\mathbb{C})$  and an affine slice  $\widehat{C} = \{\ell = 1\} \cap \overline{C}$  containing all the points  $\overline{C} \cap \mathcal{V}(x_0 x_1 \cdots x_n)$ .



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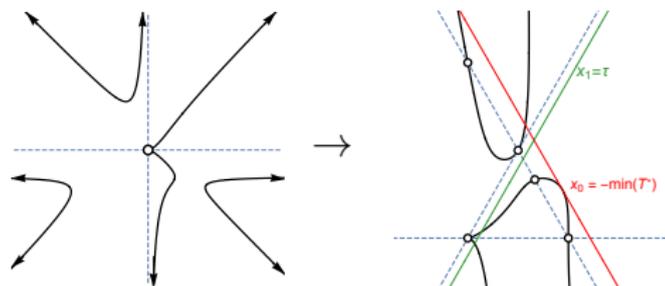
- ▶ Take  $\overline{C} \subset \mathbb{P}^n(\mathbb{C})$  and an affine slice  $\widehat{C} = \{\ell = 1\} \cap \overline{C}$  containing all the points  $\overline{C} \cap \mathcal{V}(x_0 x_1 \cdots x_n)$ .
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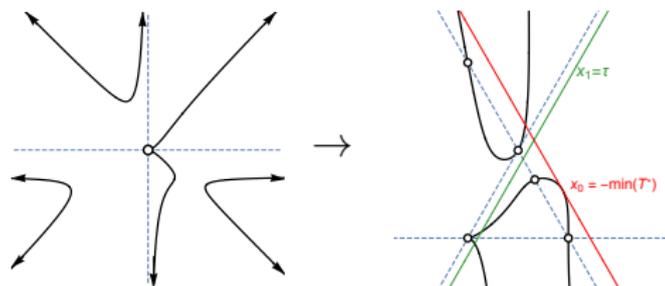
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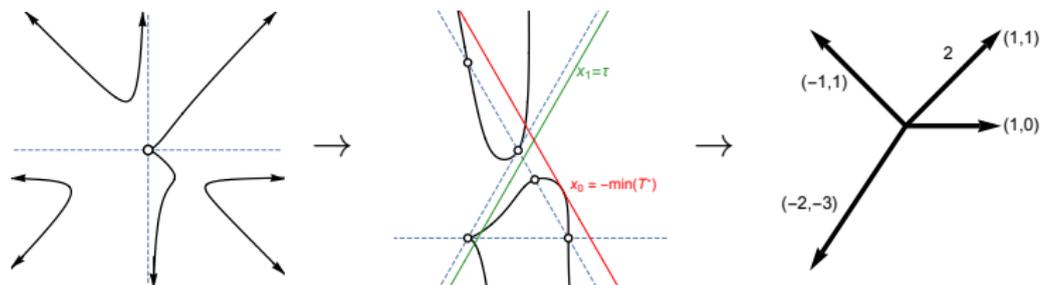
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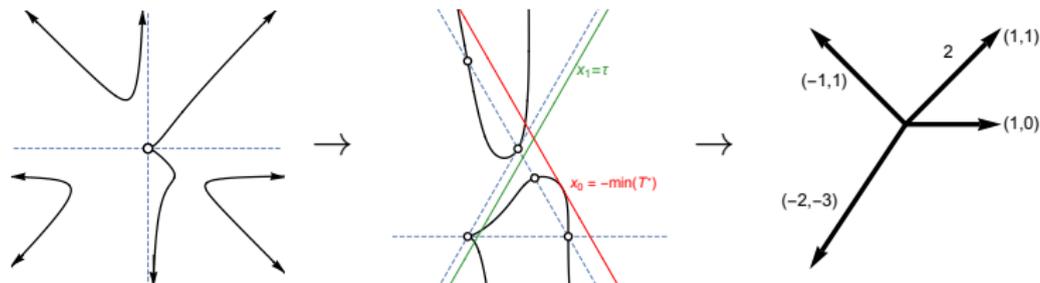
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# Example: quartic plane curve

Replace  $C = \mathcal{V}(x_1^3 x_2 - x_1 x_2^3 + x_1^3 - x_2^2) \subset \mathbb{C}^2$  with

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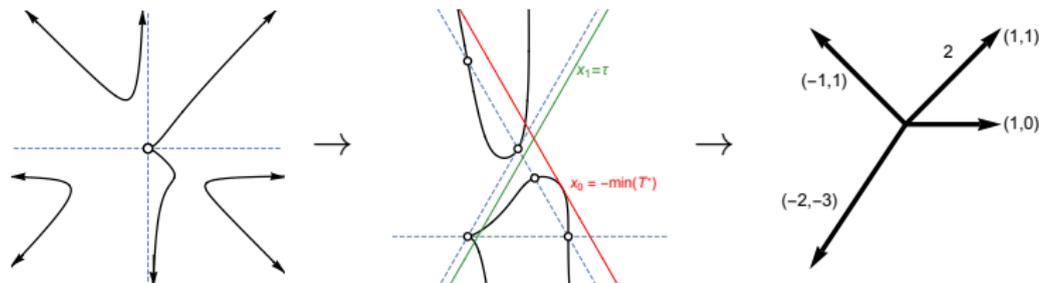


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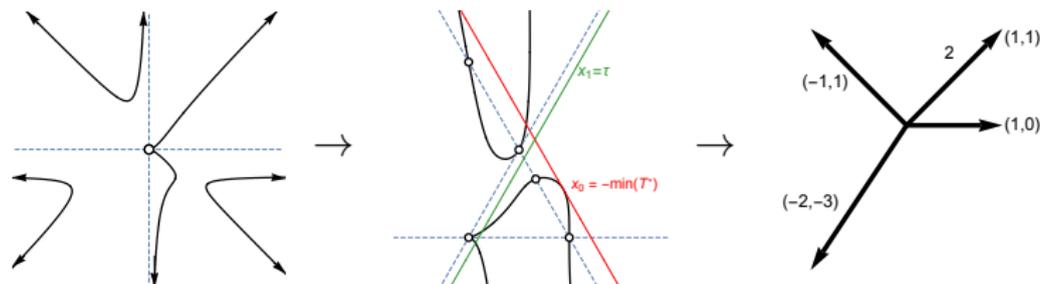
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The point  $p \approx (0.8293, 0.1, 0.0354) \in \widehat{C} \cap \{x_1 = .1\}$  tracks to  $(1, 0, 0)$ .



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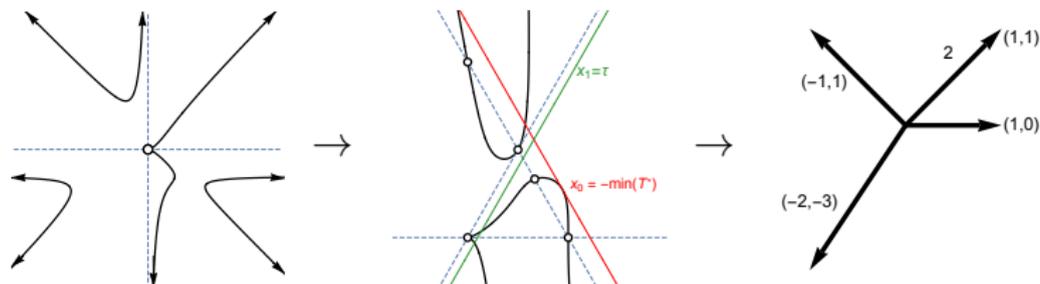
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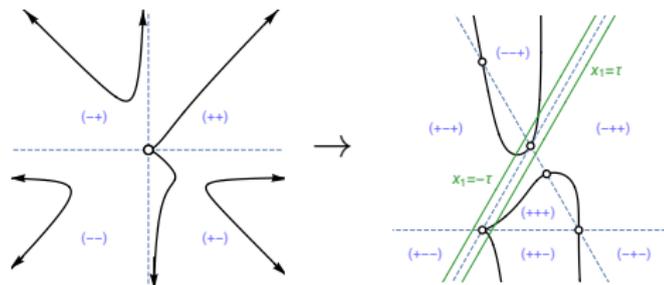
Valuation of the path is  $(0, 2, 3) \rightarrow (-2, -3) \in \text{Trop}_{\mathbb{C}}(C)$ .



# Real Tropical Strategy

We can compute  $\text{Trop}_{\mathbb{R}}(C)$  similarly to  $\text{Trop}_{\mathbb{C}}(C)$ .

This requires checking  $\{x_j = \pm\tau\}$  for real points and only considering real paths converging to  $\widehat{C} \cap \{x_j = 0\}$ .

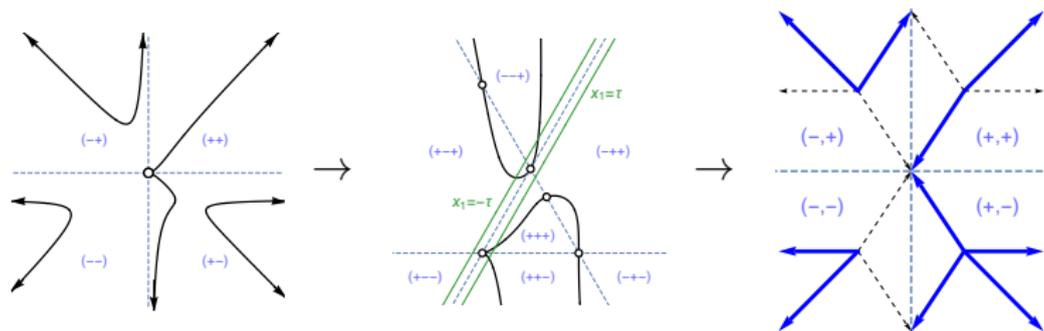


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Keeping track of signs of the parametrized paths gives the **signed real tropical variety**.



# Computing curves in large spaces

A central curve defined by ...

$$\left[ \begin{array}{l} x_1 s_1 - x_j s_j \text{ for } j = 2, \dots, 7 \\ -u_0 + t^2 - x_1 \\ -v_0 + t^4 - x_2 \\ u_1 - x_3 \\ v_1 - x_4 \\ t(u_0 + v_0) - v_1 - x_5 \\ t^2 u_0 - u_1 - x_6 \\ t^2 v_0 - u_1 - x_7 \\ t^2 x_2 + x_3 + x_7 - t^6 \\ t^2 x_1 + x_3 + x_6 - t^4 \\ tx_1 + tx_2 + x_4 + x_5 - t^3 - t^5 \\ s_1 - ts_5 - t^2 s_6 \\ s_2 - ts_5 - t^2 s_7 - 1 \\ s_3 - s_6 - s_7 \\ s_4 - s_5 \end{array} \right]$$

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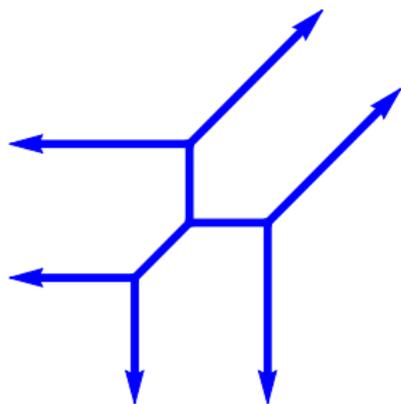
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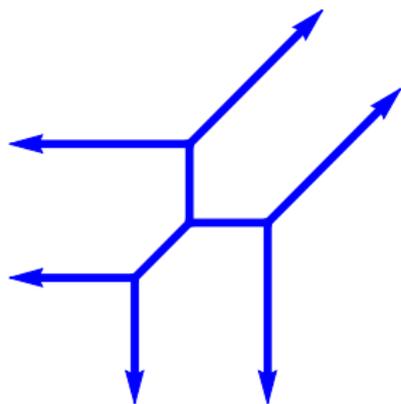
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Our algorithm find that the tropical variety  $\text{Trop}_\mathbb{C}(C) = \text{Trop}_\mathbb{R}(C)$  consists of 10 rays with multiplicites 6, 4, 3, 2, 2, 1, 1, 1, 1, 1.



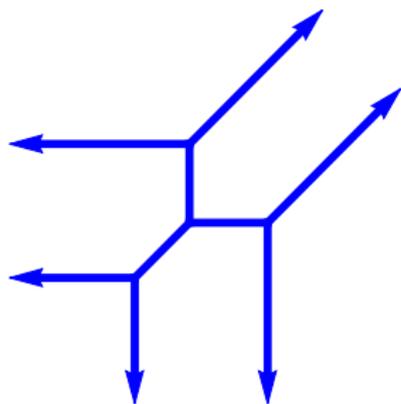
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We would like to develop these methods for varieties of dimension  $\geq 1$ , like *surfaces*, over both  $\mathbb{C}$  and  $\mathbb{R}$ .

Tropicalizations of varieties higher dimension can also be used to compute tropical varieties of ideals in  $\mathbb{k}\{\{t\}\}[x_1, \dots, x_n]$ , which are used in many applications (like central curves).



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