# Commutation Classes of Permutations 

Herman Chau

University of Washington

## Commutation Classes of [4321]



What is a Commutation Class?

## Permutations

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The longest permutation $w_{0}$ sends $i \mapsto n-i+1$.
Example ( $n=4$ )

$$
\begin{aligned}
W_{0} & =\left[\begin{array}{llll}
1 & 2 & 3 & 4 \\
4 & 3 & 2 & 1
\end{array}\right] \\
& =\left[\begin{array}{llll}
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& =\sigma_{1} \sigma_{2} \sigma_{3} \sigma_{1} \sigma_{2} \sigma_{1}
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## Wiring Diagrams



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## Reduced Expressions

## Definition

A reduced expression for $w \in S_{n}$ is a product of adjacent transpositions $w=\sigma_{i_{1}} \cdots \sigma_{i_{j}}$ of minimum length. The set of all reduced expressions for $w$ is denoted $\mathcal{R}(w)$.

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The length of a permutation $w$, denoted $\ell(w)$, is the length of any reduced expression for w.

## Fact

The longest permutation $w_{0} \in S_{n}$ has length $\binom{n}{2}$.

## Commutation and Braid Relations

Given a reduced expression of $w \in S_{n}$, we can obtain new reduced expressions via commutation and braid relations.

- Commutation: Let $1 \leq i<i+1<j \leq n-1$. Then $\sigma_{i} \sigma_{j}=\sigma_{j} \sigma_{i}$.
- Braid: Let $1 \leq i \leq n-2$. Then $\sigma_{i} \sigma_{i+1} \sigma_{i}=\sigma_{i+1} \sigma_{i} \sigma_{i+1}$.


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Example ( $\mathrm{n}=4$ )

$$
\begin{gathered}
\sigma_{1} \sigma_{2} \underline{\boldsymbol{\sigma}_{3} \boldsymbol{\sigma}_{1}} \sigma_{2} \sigma_{1} \stackrel{\text { commute }}{=} \sigma_{1} \sigma_{2} \underline{\boldsymbol{\sigma}_{1} \boldsymbol{\sigma}_{3}} \sigma_{2} \sigma_{1} \\
\sigma_{1} \sigma_{2} \sigma_{3} \boldsymbol{\sigma}_{1} \boldsymbol{\sigma}_{2} \boldsymbol{\sigma}_{1} \stackrel{\text { braid }}{=} \sigma_{1} \sigma_{2} \sigma_{3} \underline{\boldsymbol{\sigma}_{2} \boldsymbol{\sigma}_{1} \boldsymbol{\sigma}_{2}}
\end{gathered}
$$

## Commutation and Braid Relations


$\downarrow$ commute $\sigma_{3} \sigma_{1}$


## Commutation and Braid Relations


$\downarrow$ braid $\sigma_{1} \sigma_{2} \sigma_{1}$


## Word Property

Theorem (Matsumoto, Tits)
Any two reduced expressions in $\mathcal{R}(w)$ are connected by a sequence of commutation and braid relations.

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We can visualize this as a graph $G$ whose vertex set is $\mathcal{R}(w)$ and ( $\rho_{1}, \rho_{2}$ ) is an edge if and only if $\rho_{1}, \rho_{2}$ differ by a commutation or braid relation.

## Reduced Expression Graph for [4, 3, 2, 1]



## Commutation Classes

Definition
Two reduced expressions $\rho_{1}, \rho_{2}$ are commutation equivalent, denoted $\rho_{1} \sim \rho_{2}$ if they differ by a series of commutation relations.

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## Definition

$\mathcal{C}(w):=\mathcal{R}(w) / \sim$ is the set of commutation classes of $w$.

## Reduced Expression Graph for [4, 3, 2, 1]



## Commutation Class Graph for [4, 3, 2, 1]



## Wiring Diagrams of Commutation Classes



## Wiring Diagrams of Commutation Classes



## Wiring Diagrams of Commutation Classes



## Elnitsky's Bijection

## Theorem (Elnitsky, 1993)

Let $n \geq 2$. There is a bijection between $\mathcal{C}\left(w_{0}\right)$ and rhombic tilings of the regular $2 n$-gon.

## Bijection in $S_{4}$



## Proof Without Words



## Proof Without Words



## Proof Without Words



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## Bijection for General w

For general $w \in S_{n}, \mathcal{C}(w)$ is in bijection with rhombic tilings of a specific shape, not necessarily the regular $2 n$-gon.

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## Example $(w=[3,4,2,1])$



## Counting Commutation Classes

## Counting Reduced Expressions of wo

## Counting Reduced Expressions of $w_{0}$

## Theorem (Stanley, 1984) (Edelman and Greene, 1987)

The number of reduced expressions $\left|\mathcal{R}\left(w_{0}\right)\right|$ in $S_{n}$ is equal to the number of standard Young tableaux of shape ( $n-1, n-2, \ldots, 1$ ). By the hook-length formula, this is equal to $\binom{n}{2}!/ 1^{n-1} 3^{n-2} \cdots(2 n-3)^{1}$.

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| $n$ | $\left\|\mathcal{R}\left(w_{0}\right)\right\|$ |
| :--- | :--- |
| 1 | 1 |
| 2 | 1 |
| 3 | 2 |
| 4 | 16 |
| 5 | 768 |
| 6 | 292864 |
| $\ldots$ | $\ldots$ |

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Asymptotically, $\left|\mathcal{R}\left(w_{0}\right)\right|=2^{\Theta\left(n^{2} \log n\right)}$.

## Counting Commutation Classes of $w_{0}$

No exact formula is known for the number of commutation classes $\left|\mathcal{C}\left(w_{0}\right)\right|$.

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| $n$ | $\left\|\mathcal{C}\left(w_{0}\right)\right\|$ |
| :--- | :--- |
| 1 | 1 |
| 2 | 1 |
| 3 | 2 |
| 4 | 8 |
| 5 | 62 |
| 6 | 908 |
| $\ldots$ | $\ldots$ |

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| :--- | :--- |
| 1 | 1 |
| 2 | 1 |
| 3 | 2 |
| 4 | 8 |
| 5 | $62=2 \times 31$ |
| 6 | $908=2^{2} \times 227$ |
| $\ldots$ | $\ldots$ |

## Asymptotics of $\left|\mathcal{C}\left(w_{0}\right)\right|$

Theorem (Knuth, 1992)
The number of commutation classes of the longest word $w_{0} \in S_{n}$ grows as $2^{\Theta\left(n^{2}\right)}$.

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What is the constant c in the exponent?

- (Knuth, 1992) $\sim 0.1667 \leq c \leq \sim 0.7924$
- (Felsner, 1997) $c \leq \sim 0.6974$
- (Felsner and Valtr, 2011) $\sim 0.1887 \leq c \leq \sim 0.6571$
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## Upper Bound on $\left|\mathcal{C}\left(w_{0}\right)\right|$

## Definition

Let $\bar{\rho} \in \mathcal{C}\left(w_{0}\right)$. The weaving pattern of $\rho$ is a function $P_{\bar{\rho}}:[n] \rightarrow\{+,-\}^{*}$ that records the sequence of up (+) and down $(-)$ crossings in the wiring diagram of $\bar{\rho}$.

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## Example

## Upper Bound on $\left|\mathcal{C}\left(w_{0}\right)\right|$

## Lemma

The map $\bar{\rho} \mapsto P_{\bar{\rho}}$ is an injection.

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## Proof.

It suffices to find an inverse map that takes a weaving pattern $P_{\bar{\rho}}$ back to a commutation class $\bar{\rho}$.

1. Find $i$ and $i+1$ such that $P_{\bar{\rho}}(i)$ begins with $a+$ and $P_{\bar{\rho}}(i+1)$ begins with a - .
2. Record $\sigma_{i}$, delete the leading + and - , and swap rows $i$ and $i+1$ in the weaving pattern.
3. Repeat until the weaving pattern is empty.

## Upper Bound on $\left|\mathcal{C}\left(w_{0}\right)\right|$

## Lemma

The map $\bar{\rho} \mapsto P_{\bar{\rho}}$ is an injection.
Example

| 4 | --- |
| :--- | :--- |
| 3 | --+ |
| 2 | -++ |
| 1 | +++ |

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| 4 | --- |
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| 3 | --+ |
| 2 | -++ |
| 1 | +++ |

Rows 1 and 2 have a leading + and - respectively.

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## Example

$$
\begin{array}{l|ll}
4 & --- & \\
3 & --+ & \\
2 & ++ & \sigma_{1} \\
1 & ++ &
\end{array}
$$

Delete leading signs, swap, and record $\sigma_{1}$.

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## Lemma

The map $\bar{\rho} \mapsto P_{\bar{\rho}}$ is an injection.

## Example

$$
\begin{array}{l|ll}
4 & --- & \\
3 & --+ & \\
2 & ++ & \sigma_{1} \\
1 & ++ &
\end{array}
$$

Rows 2 and 3 have a leading + and - respectively.

## Upper Bound on $\left|\mathcal{C}\left(w_{0}\right)\right|$

## Lemma

The map $\bar{\rho} \mapsto P_{\bar{\rho}}$ is an injection.

## Example

$$
\begin{array}{l|ll}
4 & --- & \\
3 & + & \sigma_{1} \sigma_{2} \\
2 & -+ & \\
1 & ++ &
\end{array}
$$

Delete leading signs, swap, and record $\sigma_{2}$.

## Upper Bound on $\left|\mathcal{C}\left(w_{0}\right)\right|$

## Lemma

The map $\bar{\rho} \mapsto P_{\bar{\rho}}$ is an injection.

## Example

$$
\begin{array}{l|ll}
4 & --- & \\
3 & + & \sigma_{1} \sigma_{2} \\
2 & -+ & \\
1 & ++ &
\end{array}
$$

Choice of either rows 1 and 2, or rows 3 and 4 .

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Example

| 4 |  |  |
| :--- | :--- | :--- |
| 3 | -- |  |
| 2 | -+ | $\sigma_{1} \sigma_{2} \sigma_{3}$ |
| 1 | ++ |  |

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4 & & \\
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Example

| 4 |  |
| :--- | :--- |
| 3 |  |
| 2 | - |
| 1 | $-\quad \sigma_{1} \sigma_{2} \sigma_{3} \sigma_{1} \sigma_{2}$ |
|  |  |

## Upper Bound on $\left|\mathcal{C}\left(w_{0}\right)\right|$

## Lemma

The map $\bar{\rho} \mapsto P_{\bar{\rho}}$ is an injection.

## Example


$\overline{\sigma_{1} \sigma_{2} \sigma_{3} \sigma_{1} \sigma_{2} \sigma_{1}}$ is the desired commutation class.

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Theorem
The number of commutation classes of $w_{0}$ is at most $\prod_{i=1}^{n}\binom{n-1}{n-i}$.

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Taking logarithms,

$$
\log \prod_{i=1}^{n}\binom{n-1}{n-i} \leq n \log (n-1)!-2 \sum_{i=0}^{n-1} \log j!
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So $\left|\mathcal{C}\left(w_{0}\right)\right| \leq e^{n^{2} / 2}=2^{n^{2} \log _{2}(e) / 2}$ and $\frac{1}{2} \log _{2}(e) \approx 0.7213$.

## Upper Bound on $\left|\mathcal{C}\left(w_{0}\right)\right|$

What is the constant c in the exponent?

- (Knuth, 1992) $\sim 0.1667 \leq c \leq \sim 0.7924$
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## Lower Bound on $\left|\mathcal{C}\left(w_{0}\right)\right|$

## Idea

Suppose $n=3 k$. Count commutation classes of the form $\overline{\rho_{0} \rho_{1} \rho_{2} \rho_{3}}$, where

- $\rho_{0}$ yields $[\underline{2 k+1, \ldots, 3 k}, \underline{k+1, \ldots, 2 k}, \underline{1, \ldots, k]}$,
- $\rho_{1}$ yields $[1, \ldots, k, k+1, \ldots, 2 k, 3 k, \ldots, 2 k+1]$,
- $\rho_{2}$ yields $[1, \ldots, k, 2 k, \ldots, k+1,2 k+1, \ldots, 3 k]$,
- $\rho_{3}$ yields $[k, \ldots, 1, k+1, \ldots, 2 k, 2 k+1, \ldots, 3 k]$.


## Lower Bound on $\left|\mathcal{C}\left(w_{0}\right)\right|$



## Lower Bound on $\left|\mathcal{C}\left(w_{0}\right)\right|$

Some notation:

- $B_{k}:=\left|\mathcal{C}\left(w_{0}\right)\right|$ where $w_{0} \in S_{k}$
- $w_{3 k}:=[2 k+1, \ldots 3 n, k+1, \ldots, 2 k, 1, \ldots k] \in S_{3 k}$


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## Observation

A lower bound for $B_{3 k}$ is $B_{3 k} \geq\left|\mathcal{C}\left(w_{3 k}\right)\right| \cdot B_{k}^{3}$.

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How do we estimate $\left|\mathcal{C}\left(w_{3 k}\right)\right|$ ?

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## Observation

A lower bound for $B_{3 k}$ is $B_{3 k} \geq\left|\mathcal{C}\left(w_{3 k}\right)\right| \cdot B_{k}^{3}$.
How do we estimate $\left|\mathcal{C}\left(w_{3 k}\right)\right|$ ?
Count rhombic tilings via Elnitsky's bijection!

## Lower Bound on $\left|\mathcal{C}\left(w_{0}\right)\right|$

Example ( $\mathrm{k}=2$ )


## Lower Bound on $\left|\mathcal{C}\left(w_{0}\right)\right|$

Example ( $\mathrm{n}=2$ )


## Lower Bound on $\left|\mathcal{C}\left(w_{0}\right)\right|$

## Definition

A plane partition is a matrix of nonnegative integers $\pi_{i, j}$ such that $\pi_{i, j} \geq \pi_{i, j+1}$ and $\pi_{i, j} \geq \pi_{i+1, j}$.

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The plane partition $\begin{array}{ll}2 & 1 \\ 1 & 0\end{array}$

## Lower Bound on $\left|\mathcal{C}\left(w_{0}\right)\right|$

## Lemma

The commutation classes of $w_{3 k}$ are in bijection with $k \times k$ plane partitions with entries at most $k$.

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## Lemma

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## Theorem (MacMahon, 1916)

The number of $i \times j$ plane partitions with entries at most $k$ is equal to

$$
\prod_{a=0}^{i-1} \prod_{b=0}^{j-1} \prod_{c=0}^{k-1} \frac{a+b+c+2}{a+b+c+1}
$$

## Lower Bound on $\left|\mathcal{C}\left(w_{0}\right)\right|$

To approximate the number of such $k \times k$ plane partitions with entries at most $k$, we take logarithms.

$$
\log \prod_{0 \leq a, b, c \leq k-1} \frac{a+b+c+2}{a+b+c+1}
$$

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This yields $\left|\mathcal{C}\left(w_{3 k}\right)\right| \approx e^{0.7848 k^{2}}$.

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Thus, $\log _{2}\left(B_{3 k}\right) \approx \log _{2}(e) \cdot 0.7848 \cdot \frac{3}{2} \cdot \frac{1}{3^{2}} k^{2} \approx 0.1887 k^{2}$.

## Lower Bound on $\left|\mathcal{C}\left(w_{0}\right)\right|$

What is the constant c in the exponent?

- (Knuth, 1992) $\sim 0.1667 \leq c \leq \sim 0.7924$
- (Felsner, 1997) $c \leq \sim 0.6974$
- (Felsner and Valtr, 2011) $\sim 0.1887 \leq c \leq \sim 0.6571$
- (Dumitrescu and Mandal, 2020) $\sim 0.2083 \leq c$


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- How are the sizes of commutation classes distributed?

Thank You!

